

# Keccak, More Than Just SHA3SUM

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# Outline

- 1 How it all began
- 2 Introducing KECCAK
- 3 More than just SHA3SUM
- 4 Inside KECCAK
- 5 KECCAK and the community

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# Let's talk about hash functions...



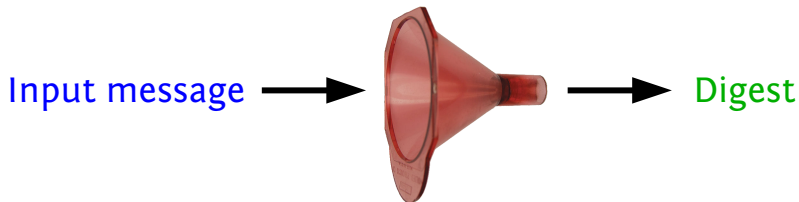
???



These are “hashes” of some sort, but they ain't **hash functions**...

# Cryptographic hash functions

$$h : \{0, 1\}^* \rightarrow \{0, 1\}^n$$



- MD5:  $n = 128$  (Ron Rivest, 1992)
- SHA-1:  $n = 160$  (NSA, NIST, 1995)
- SHA-2:  $n \in \{224, 256, 384, 512\}$  (NSA, NIST, 2001)

# Why should you care?

- You probably use them several times a day:
  - website authentication,
  - digital signature,
  - home banking,
  - secure internet connections,
  - software integrity,
  - version control software,
  - ...

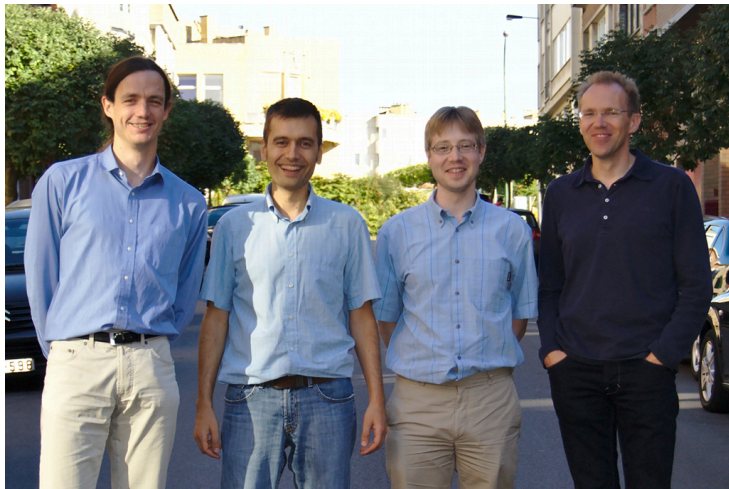
# Breaking news in crypto



- 2004: SHA-0 broken (Joux et al.)
- 2004: MD5 broken (Wang et al.)
- 2005: practical attack on MD5 (Lenstra et al., and Klima)
- 2005: SHA-1 theoretically broken (Wang et al.)
- 2006: SHA-1 broken further (De Cannière and Rechberger)
- **2007: NIST calls for SHA-3**

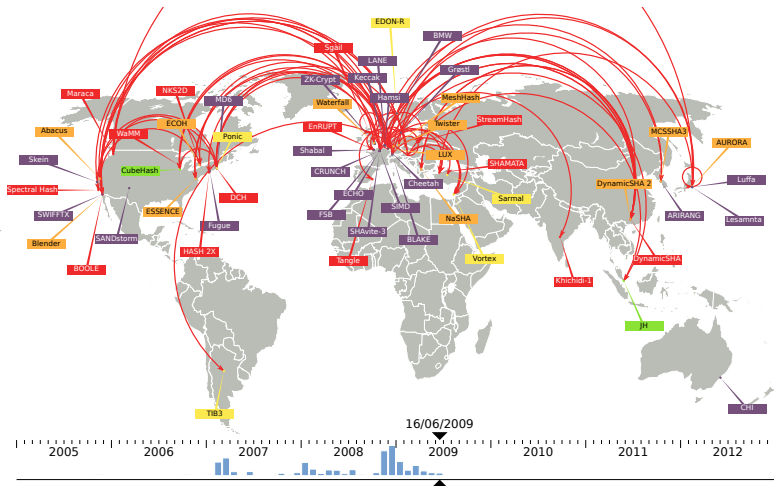
**Who answered NIST's call?**

# KECCA Team to the rescue!





# The battlefield



[courtesy of Christophe De Cannière]

# SHA-3 time schedule



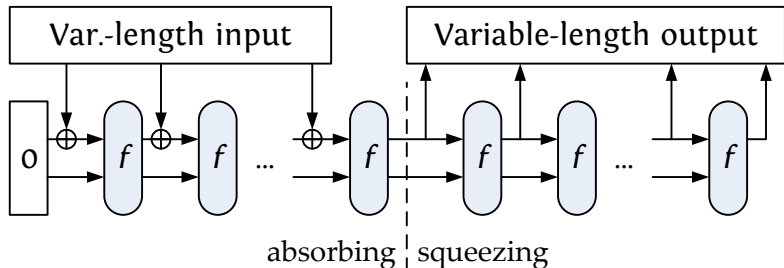
- 2007: SHA-3 initial call
- 2008: submission deadline
- 2009: first SHA-3 conference
- 2010: second SHA-3 conference
- 2010: **finalists** are Blake, Grøstl, JH, KECCAK and Skein
- 2012: final SHA-3 conference
- Oct. 2, 2012: **KECCAK** wins!

Participants: 64 → 51 → 14 → 5 → 1

# Outline

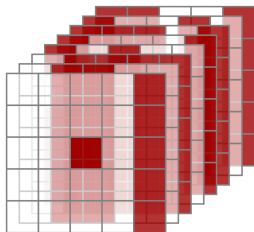
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# KECCAK, a sponge function



- Arbitrary *input* and *output* length
- More flexible than regular hash functions
- Parameters
  - $r$  bits of *rate* (defines the **speed**)
  - $c$  bits of *capacity* (defines the **security** level)
- Use the **permutation** KECCAK- $f$

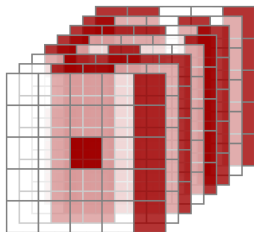
# The seven permutation army



- $(5 \times 5)$  lanes
- up to 64-bit each

- 7 permutations:
  - 25, 50, 100, 200, 400, 800, 1600 bits
  - toy, lightweight, fastest
- repetition of a simple round function
  - operates on a 3D state
- like a block cipher but **without** a key

# The seven permutation army



- $(5 \times 5)$  lanes
- up to 64-bit each

- First, choose your permutation ...
  - e.g. *width* = 1600
- ...then choose the *rate* and *capacity*
  - such that *rate* + *capacity* = 1600
- Security-speed trade-offs using the same permutation:

Rate	Capacity	Strength	Speed
1344	256	128	×1.312
1216	384	192	×1.188
1088	512	256	×1.063
1024	576	288	1.000

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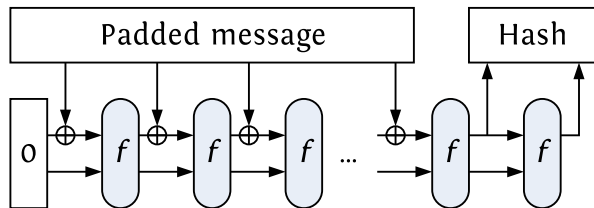
# One primitive to rule them all



- **Full range** of cryptographic functions
  - hashing (regular, salted)
  - key derivation
  - message authentication
  - encryption
  - authenticated encryption
- ...in a **simple way**
  - simple & straightforward usage
  - easy to understand security claim
- ...and increasing **diversity** of standard portfolio
  - very different from SHA-1 and SHA-2
  - very different from AES and block cipher modes

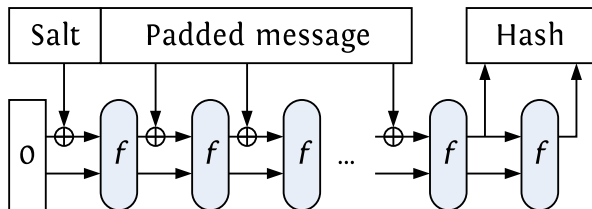


## Use KECCAK for regular hashing



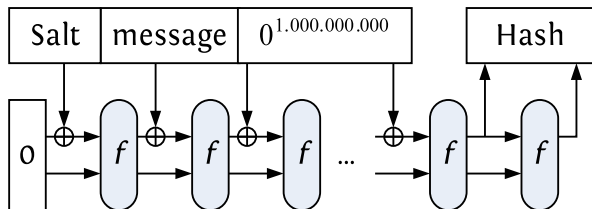
- Electronic signatures, message integrity (*GPG, X.509 ...*)
- Data integrity (*shaxsum ...*)
- Data identifier (*Git, Mercurial, online anti-virus, peer-2-peer ...*)

# Use KECCAK for salted hashing



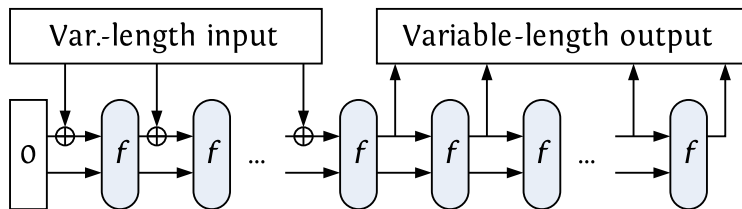
- Goal: defeat rainbow tables
- Web cookie
- Password storage and verification (*Kerberos*, */etc/shadow* ...)

# Use KECCAK for salted hashing



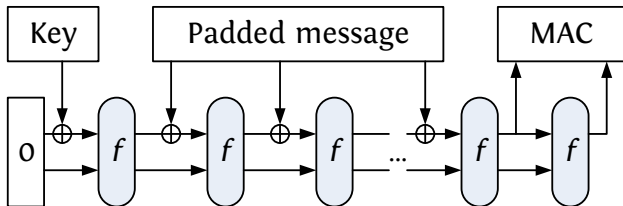
- Goal: defeat rainbow tables
- Web cookie
- Password storage and verification (*Kerberos, /etc/shadow ...*)
- ...Can be as **slow** as you like it!

# Use KECCAK as a mask generation function



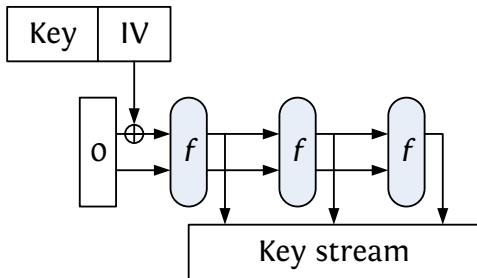
- Key derivation function in SSL, TLS
- Full-domain hashing in public key cryptography
  - electronic signatures RSA PSS [PKCS#1]
  - encryption RSA OAEP [PKCS#1]
  - key establishment RSA KEM [IEEE Std 1363a]

## Use KECCAK for MACing



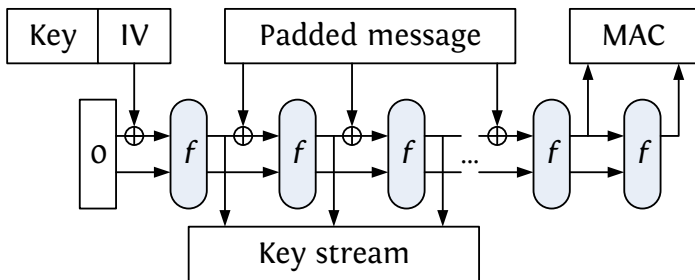
- As a message authentication code
- Simpler than HMAC [FIPS 198]
  - HMAC: special construction for MACing with SHA-1 and SHA-2
  - Required to plug a security hole in SHA-1 and SHA-2
  - No longer needed for KECCAK which is sound

# Use KECCAK for (stream) encryption



- As a stream cipher

# Single pass authenticated encryption



- Authentication and encryption in a **single** pass!
- Secure messaging (*SSL/TLS, SSH, IPSEC ...*)
- **Same** primitive KECCAK- $f$  but in a (slightly) different mode
  - **Duplex** construction
  - Also for random generation with **reseeding** (`/dev/urandom ...`)

# Tuning KECCAK to your own security requirements

Online tool available at <http://keccak.noekeon.org/tune.html>

## Tune KECCAK to your requirements

---

The capacity parameter and chosen output length in KECCAK can be freely chosen. Their combination determines the attainable security and the capacity has an impact on performance. This page gives you the optimal capacity and output length values, given the classical hash function criteria.

**Required collision resistance:  $2^x$**      $x =$

**Required (second) preimage resistance:  $2^y$**      $y =$

Please specify your requirements...



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Required collision resistance:  $2^x$   $x =$

Required (second) preimage resistance:  $2^y$   $y =$

The optimal choice of parameters is:

**KECCAK[r=1344,c=256]** with a least **256 bits** of output.

## Speed

For long messages, this function is **31.3% faster than** KECCAK[] (KECCAK with the default parameters). On the reference processor proposed by NIST, long messages should take about **9.6 cycles/byte**.

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Required collision resistance: $2^x$	x =	<input type="text" value="128"/>	⬇
Required (second) preimage resistance: $2^y$	y =	<input type="text" value="256"/>	⬇

The optimal choice of parameters is:

**KECCAK[r=1088,c=512]** with a least **256 bits** of output.

## Speed

For long messages, this function is **6.25% faster than** KECCAK[] (KECCAK with the default parameters). On the reference processor proposed by NIST, long messages should take about **11.9 cycles/byte**.

# Tuning KECCAK to your own security requirements

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## Security claim

In line with our **hermetic sponge strategy**, we make a **flat sponge claim** with  $c=256$  bits of capacity: for any output length, we claim this KECCAK sponge function resists any attack up to  $2^{128}$  operations (each of complexity equivalent to one call to KECCAK- $f$ ), unless easier on a random oracle. For 256 bits of output specifically, this translates into the following claimed security level:

	Claimed security level
<b>Collision resistance</b>	$2^{128}$
<b>(Second) preimage resistance</b>	$2^{128}$

## Addendum: how big is $2^{128}$ ?

If an attacker has access to one billion computers, each performing one billion evaluations of KECCAK- $f$  per second, it would take about  $1.1 \times 10^{13}$  years (**770 times the estimated age of the universe**) to evaluate the permutation  $2^{128}$  times.

# Tuning KECCAK to your own security requirements

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## Security claim

In line with our **hermetic sponge strategy**, we make a **flat sponge claim** with  $c=512$  bits of capacity: for any output length, we claim this KECCAK sponge function resists any attack up to  $2^{256}$  operations (each of complexity equivalent to one call to KECCAK- $f$ ), unless easier on a random oracle. For 256 bits of output specifically, this translates into the following claimed security level:

	Claimed security level
<b>Collision resistance</b>	$2^{128}$
<b>(Second) preimage resistance</b>	$2^{256}$

## Addendum: how big is $2^{256}$ ?

If an attacker has access to one billion computers, each performing one billion evaluations of KECCAK- $f$  per second, it would take about  $3.7 \times 10^{51}$  years ( $2.6 \times 10^{41}$  times the estimated age of the universe) to evaluate the permutation  $2^{256}$  times.

Considering an irreversible computer working at 2.735°K (the average temperature of the universe), Landauer's principle implies that it cannot consume less than  $2.62 \times 10^{-23}$  joule every time a bit is changed. (Computers actually consume much more than that.) Just counting from 1 to  $2^{256}$  would take at least  $3 \times 10^{54}$  joules (the total energy output of the Sun during  $2.5 \times 10^{20}$  years).

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# KECCAK-f in pseudo-code

```

KECCAK-F[b](A) {
  forall i in 0...nr-1
    A = Round[b](A, RC[i])
  return A
}

Round[b](A, RC) {
  θ step
  C[x] = A[x,0] xor A[x,1] xor A[x,2] xor A[x,3] xor A[x,4], forall x in 0...4
  D[x] = C[x-1] xor rot(C[x+1],1), forall x in 0...4
  A[x,y] = A[x,y] xor D[x], forall (x,y) in (0...4,0...4)

  ρ and π steps
  B[y,2*x+3*y] = rot(A[x,y], r[x,y]), forall (x,y) in (0...4,0...4)

  χ step
  A[x,y] = B[x,y] xor ((not B[x+1,y]) and B[x+2,y]), forall (x,y) in (0...4,0...4)

  ι step
  A[0,0] = A[0,0] xor RC

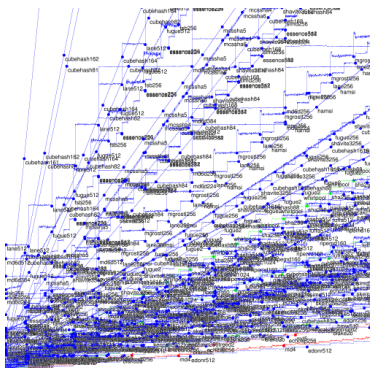
  return A
}

```

[http://keccak.noekeon.org/specs\\_summary.html](http://keccak.noekeon.org/specs_summary.html)

# Performance in software

- Faster than SHA-2 on all modern PC
- KECCAKTREE faster than MD5 on some platforms



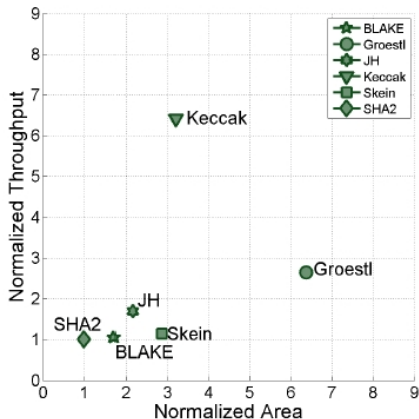
C/b	Algo	Strength
4.79	keccakc256treed2	128
4.98	md5 <b>broken!</b>	64
5.89	keccakc512treed2	256
6.09	sha1 <b>broken!</b>	80
8.25	keccakc256	128
10.02	keccakc512	256
13.73	sha512	256
21.66	sha256	128

[eBASH, hydra-6, <http://bench.cr.yp.to/>]

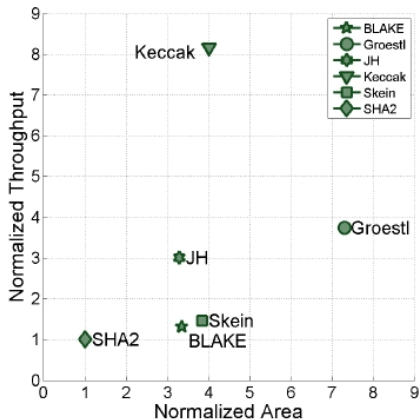
# Efficient and flexible in hardware

From Kris Gaj's presentation at SHA-3, Washington 2012:

## ASIC

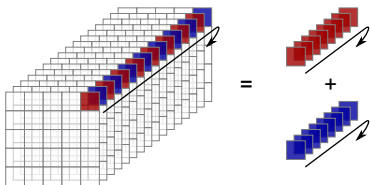


## Stratix III FPGA





# Bit interleaving



- Ex.: map 64-bit lane to 32-bit words
  - $\rho$  seems the critical step
  - **Even** bits in one word
  - **Odd** bits in a second word
  - $\text{ROT}_{64} \leftrightarrow 2 \times \text{ROT}_{32}$
- Can be generalized
  - to 16- and 8-bit words
- Can be combined
  - with lane/slice-wise architectures
  - with most other techniques
- No mismatch CPU words vs. security level

[KECCAK impl. overview, Section 2.1]

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## SHA-3, an open contest

- Open submissions, as required by NIST:
  - Public algorithm details
  - Open-source reference and optimized implementations
  - No patents
- Open cryptanalysis
- Open benchmarks [eBASH] [XBX]

### KECCAKTOOLS

A set of **documented** C++ classes to help analyze KECCAK-*f*

- To encourage cryptanalysis (*we use it too!*)
- To help verify our claims [KECCAK Team, FSE 2012]
- And also to generate optimized code

# Prizes for best cryptanalysis results



- **Four** cryptanalysis prizes awarded!
- 25 bottles of Belgian trappist beer  
[CICO problem & cube testers, Aumasson and Khovratovich]
- Bialetti coffee machine  
[zero-sum, Aumasson and Meier]
- Lambic-based beers and book  
[zero-sum, Boura and Canteaut]
- Belgian finest chocolates  
[second preimage, Bernstein]

# Crunchy Crypto Collision and Preimage Contest

- Goal:
  - Motivate 3rd-party cryptanalysis
  - Give an instant view on current state-of-the-art
- Scope: 1 to 12 rounds, including smaller instances
  - KECCAK[ $r = 40, c = 160$ ], ← *no challenge broken yet!*
  - KECCAK[ $r = 240, c = 160$ ],
  - KECCAK[ $r = 640, c = 160$ ], and
  - KECCAK[ $r = 1440, c = 160$ ]
- Results so far:
  - Preimages found for 1-2 rounds
  - Collisions found for 1-4 rounds

[http://keccak.noekeon.org/crunchy\\_contest.html](http://keccak.noekeon.org/crunchy_contest.html)

# Hex-Hot-Ticks



- Contest for stimulating developers in using "exotic" platforms
- Winners:
  - KECCAK on a NVIDIA GPU using CUDA [Gerhard Hoffmann]
  - KECCAKTREE on a NVIDIA GPU also using CUDA [Guillaume Sevestre]

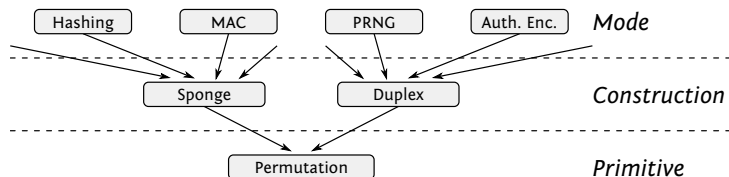
# Implementations

- Reference implementations
  - Focused on readability
  - In C, C++ and Python
- Optimized implementations
  - For 8-bit, 32-bit (bit interleaving), 64-bit platforms + 128-bit SIMD
  - In C or in assembly (x86, ARM, AVR)
  - In-place for reduced memory footprint
  - KECCAKTOOLS to generate optimized code

Available at <http://keccak.noekeon.org/files.html>

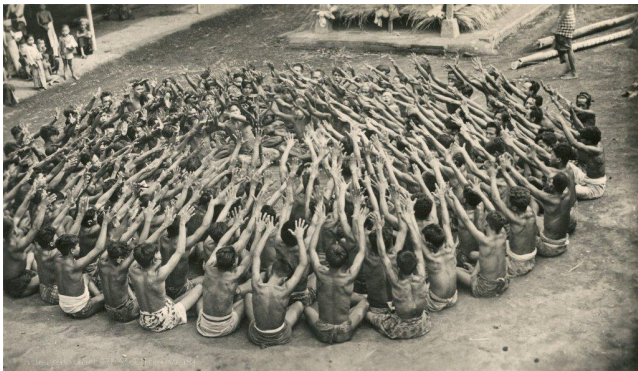
# Do you want to help?

- You can
  - make static / dynamic libraries,
  - optimize current implementations,
  - write a new implementation in your favorite language.
- Implementation-oriented doc. [KECCAK implementation overview]
- Please respect the SPONGE / DUPLEX interfaces
  - API guideline to be published soon
- SHA-3 not standardized yet!





# Questions?



More information on

<http://sponge.noekon.org/>

<http://keccak.noekon.org/>

# Credits

- Creative Commons Attribution
  - [http://en.wikipedia.org/wiki/File:Japanese\\_Secret\\_Puzzle\\_Box.jpg](http://en.wikipedia.org/wiki/File:Japanese_Secret_Puzzle_Box.jpg), from Nipaylah.
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  - <http://www.flickr.com/photos/marcelgermain/2078076913/>, from MarcelGermain (via <http://photopin.com>).
- SHA-3 battlefield picture courtesy of Christophe De Cannière