

# ~~Secure applications on top of L4~~

FOSDEM'14

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2014

# Security Gap

Red FOSDEM '14

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2014

# Syuzhet

- Intro
- Myth about uKernel: Security vs Performance
- Attacks on stack
- W xor X memory support in L4Re
- Conclusion

# About me

- Sartakov A. Vasily
- Ksys labs – Small RnD company
  - Mobile and network Hardware-software systems
  - Not only uKernels
  - Open Source and Research projects
    - Evaluate, apply, implement
    - Industry point of view

# About us

- Joined to community 3 years ago
- Fiasco.OC + L4Re
- Genode
- Respect Open Source – we publish too.

# What has changed since..?

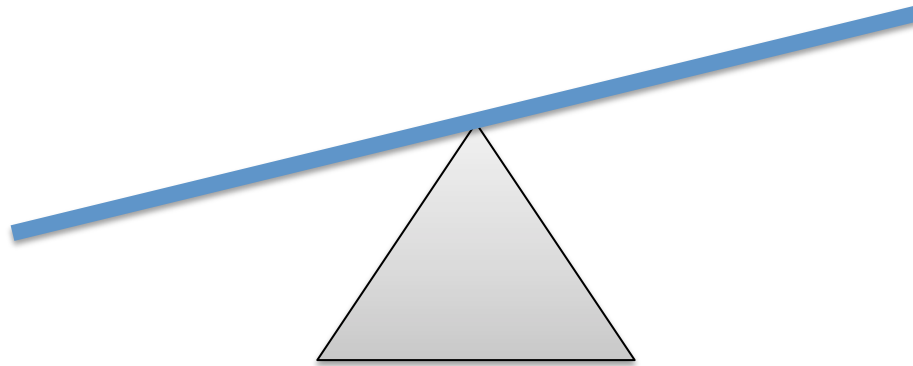
- Transformation from university to commercial projects
- New step of maturity

# Syuzhet

- Intro
- **Myth about uKernel: Security vs Performance**
- Attacks on stack
- W xor X memory support in L4Re
- Conclusion

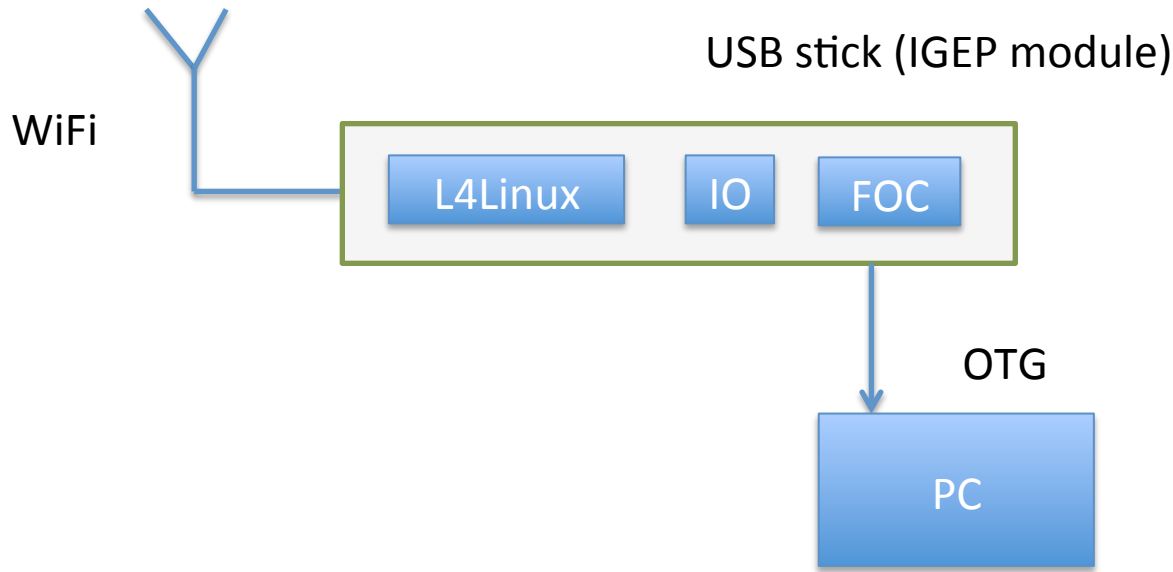
# Myth

- Dialectical pairs: performance vs security
- Security is a most strong part...(??)





# L4Linux, USB-OTG, WiFi, Omap3

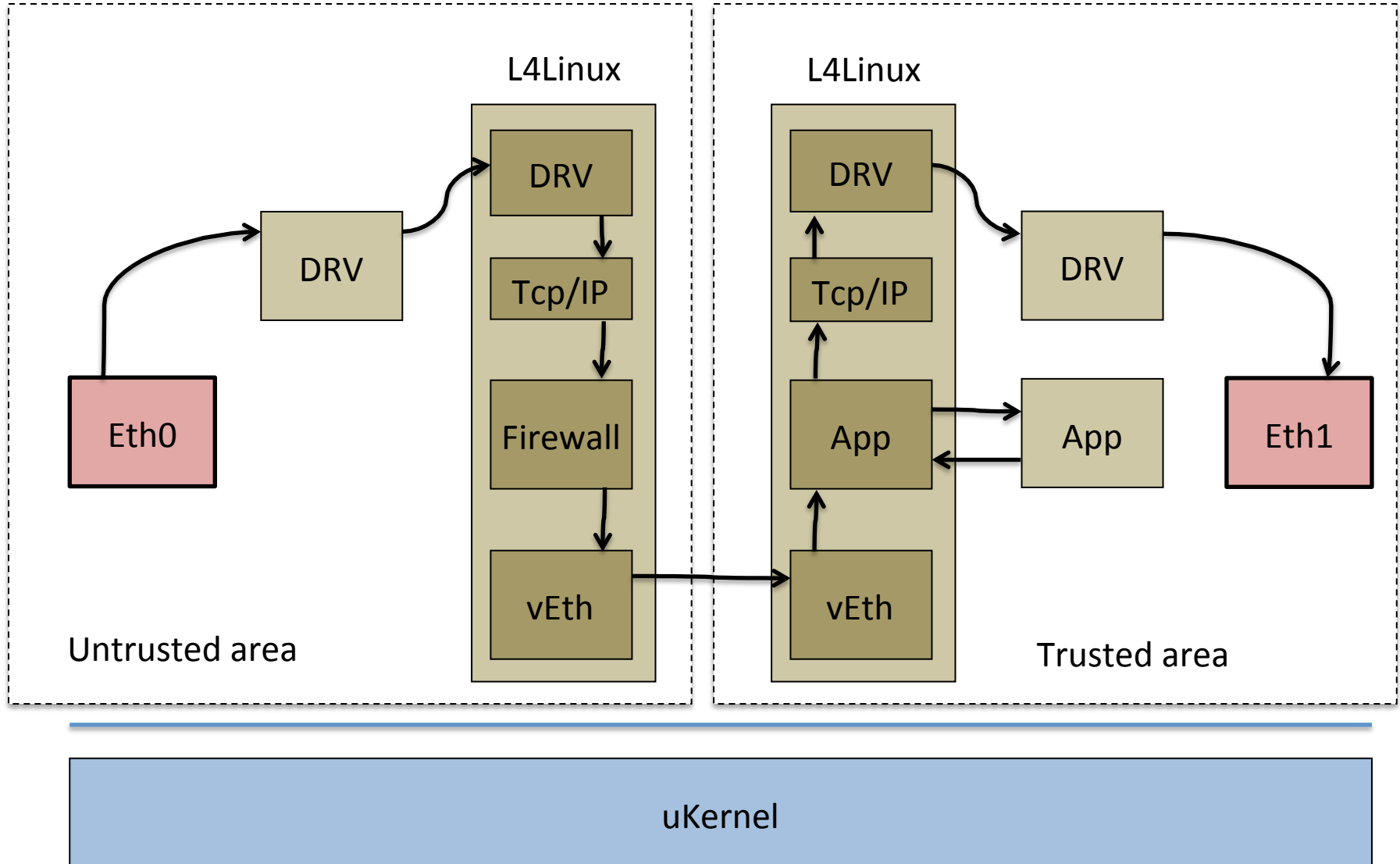


Linux vs uKernel:

L: 5.5 Mbits/sec

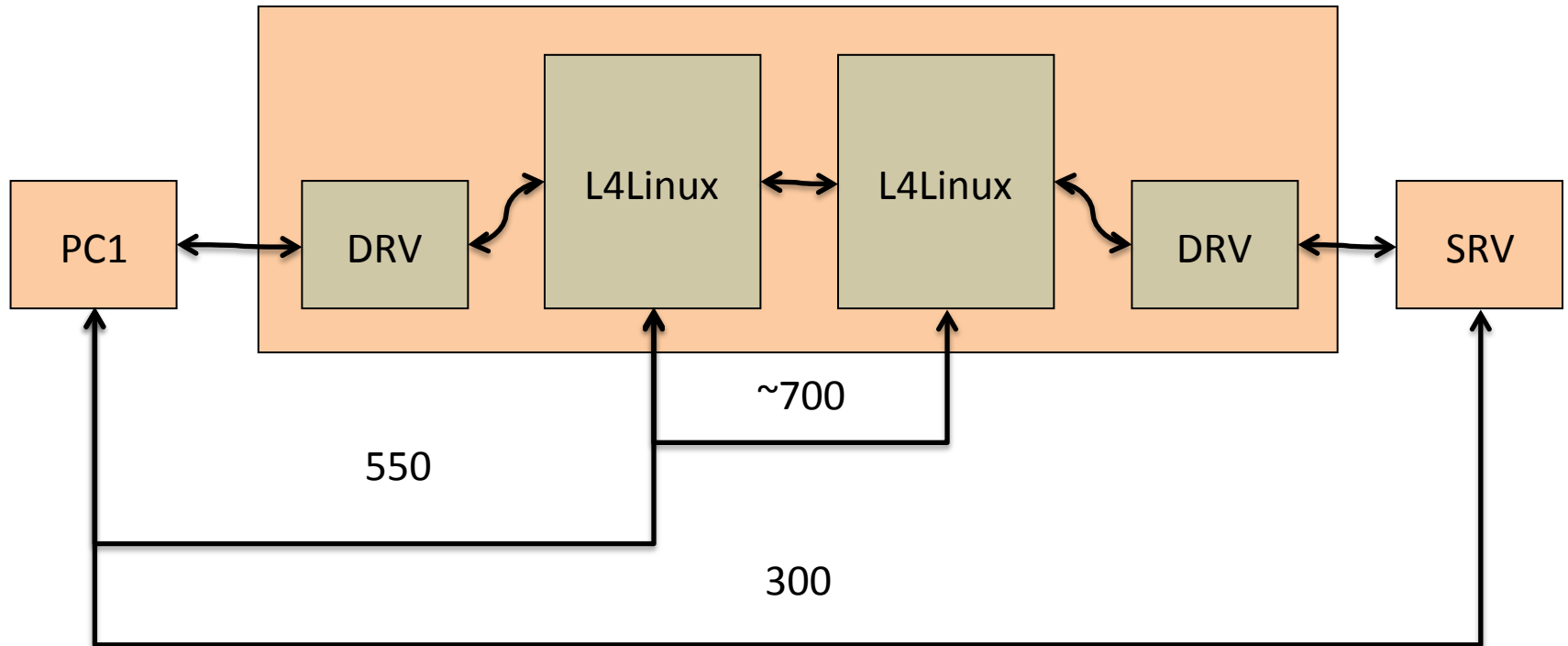
L4Linux: 7.35 Mbits/sec

# Example2: Network service on top x86



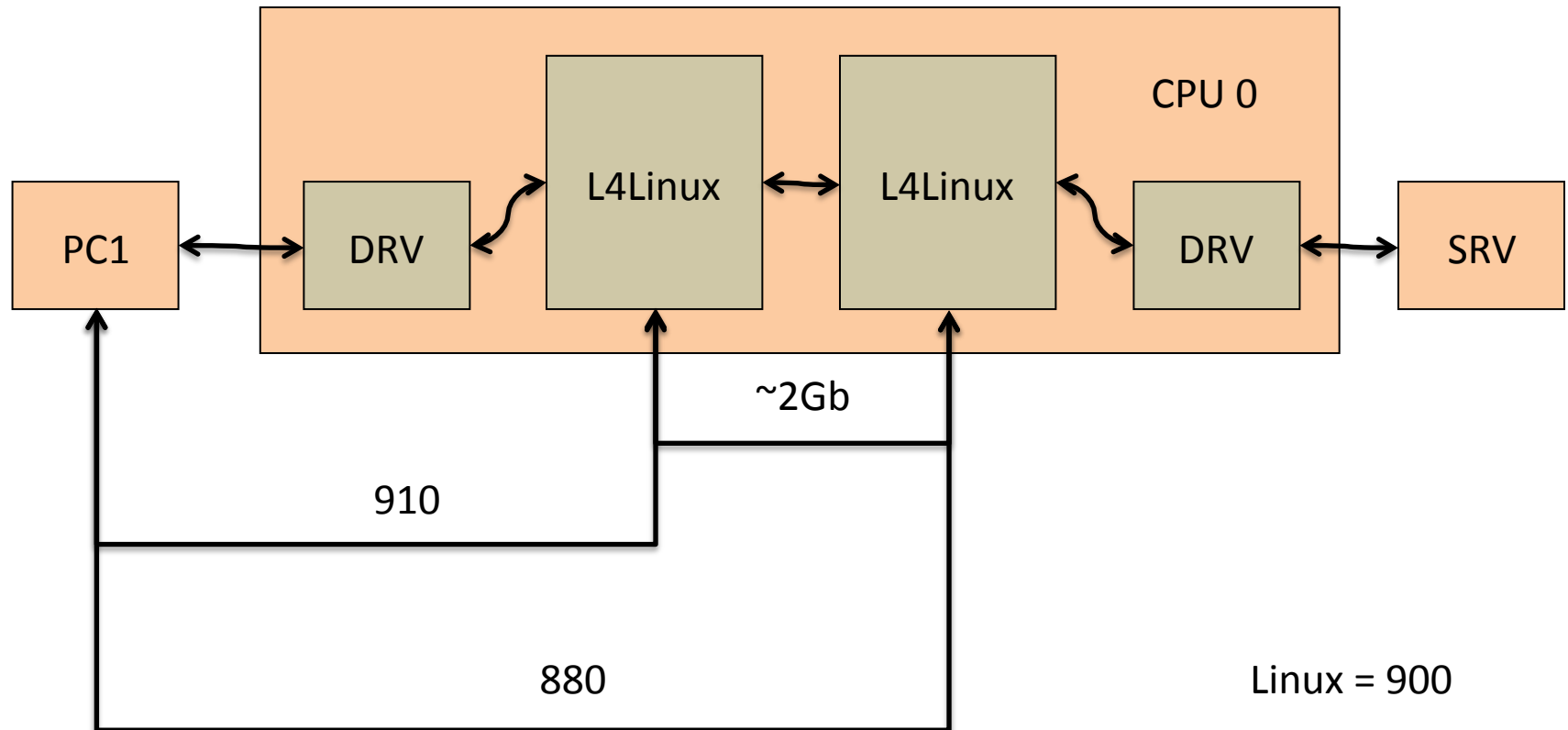
# Example2: Performance tests (Genode)

Linux = 900

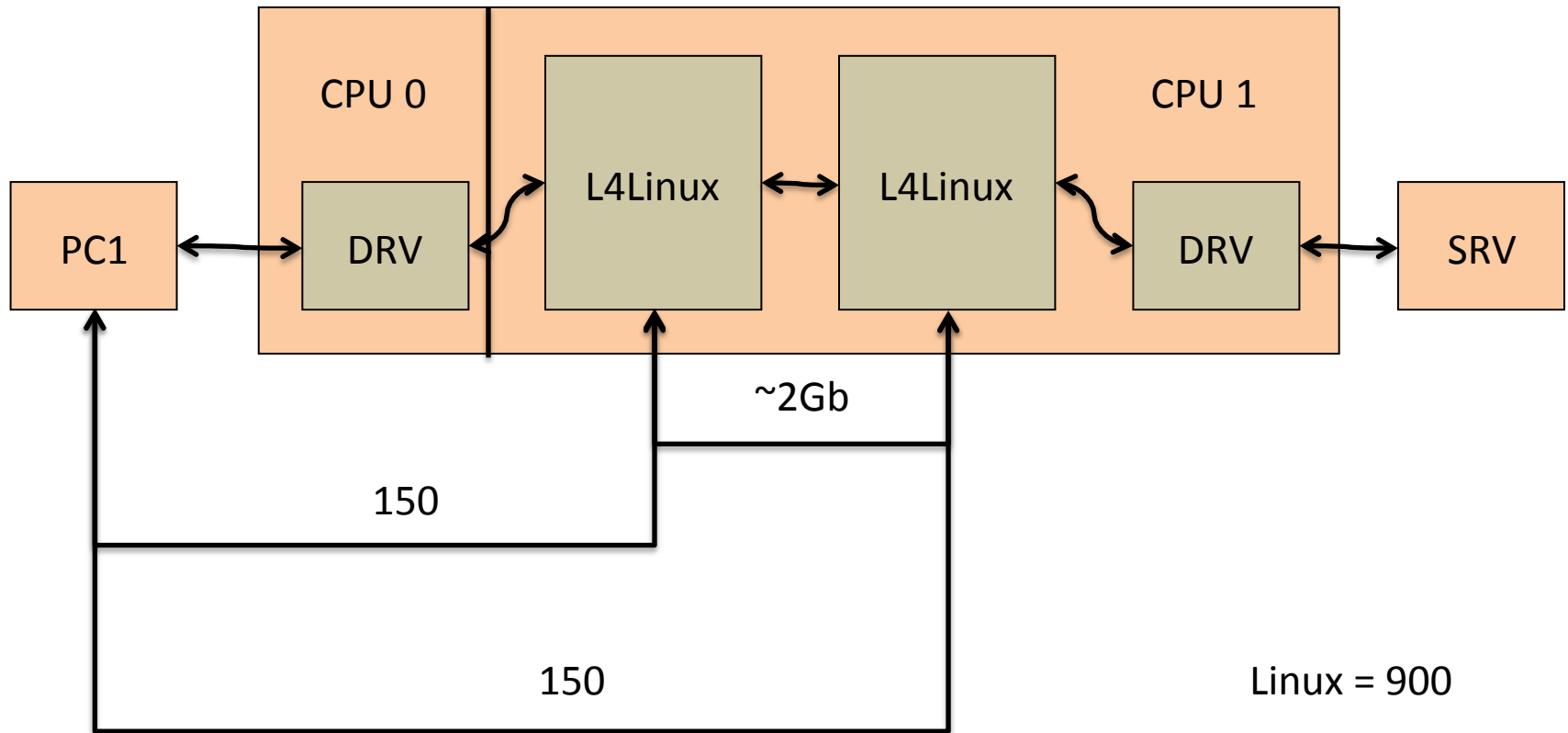


Disclaimer: This is a old experiment result, since that time Genodes rework event mechanism and add some new software

# Example2: Performance tests (L4Re)



# Example3: Performance tests (L4Re) (SMP)

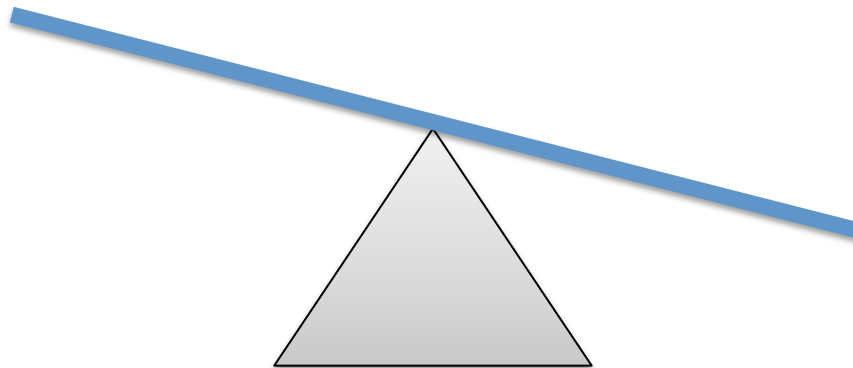


# Conclusion

- Performance is not an unsolvable problem
- Performance is not an out-of-box feature

# Come back again to Myth

- uKernels are secure
  - Small size of TCB (less errors, verification)
  - Stable api/abi
  - Drivers in userspace
  - Isolation/separation of components
- This is true, but...



# Stack protection

- Linux: Exec shield, since 2003
- Linux: PaX, since 2000
- Windows (sic!) : DEP, since XP SP2 (!!!)
  
- Genode – “canaries” is disabled in toolchains (StackGuard)
- L4Re – “canaries” is disabled at compilation time by gcc flag



# Why I am care

- Third part software (Linux, BSD)
- A lot of wrappers
- -> potentially vulnerable points
- -> malicious software and intrusion

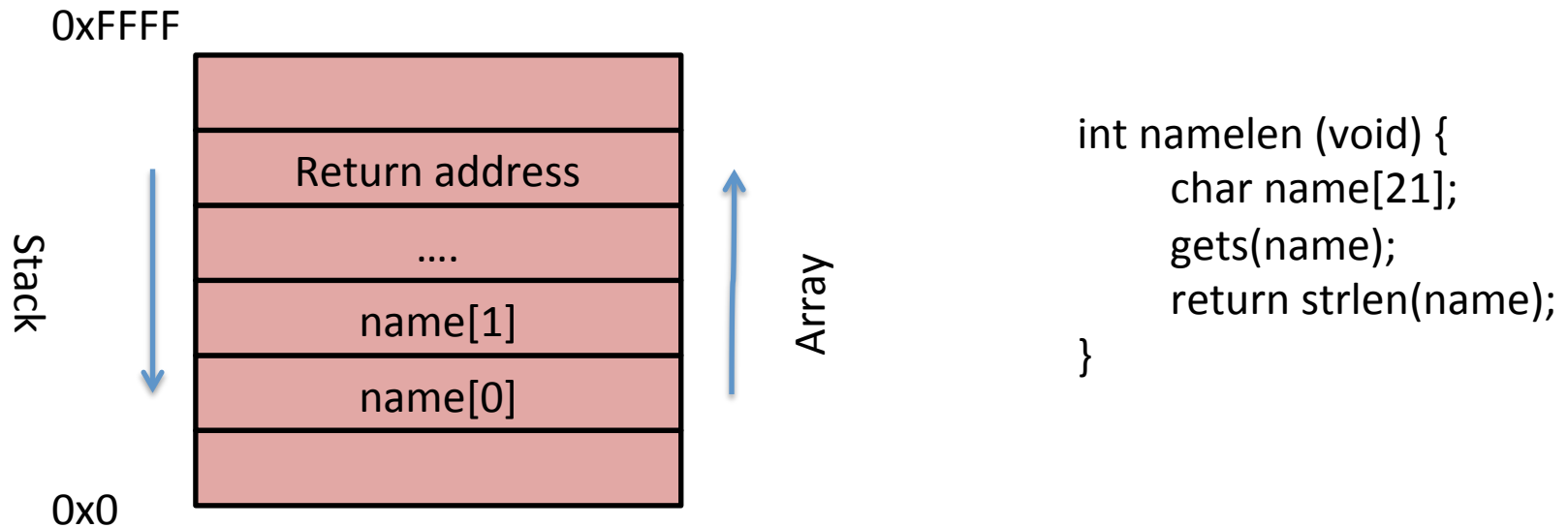
# Syuzhet

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# Smashing the stack for fun and profit ©

- Von Neumann architecture:
  - Data and instructions are in the same place
  - There is not difference
  - Type of memory defined by operations on it
  - Data can be used as instructions

# Smashing the stack



- Best case – Segfault
- Worst case - malicious execution

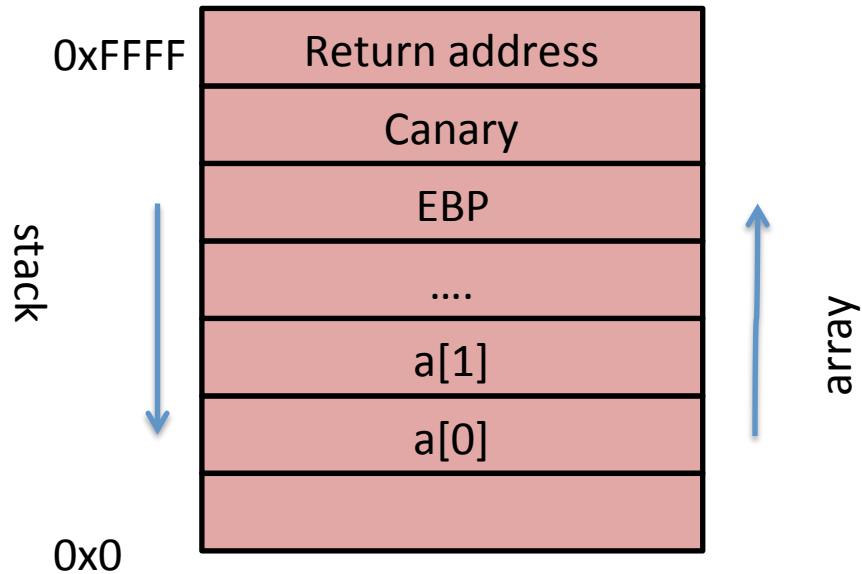
# Intrusion

- Payload
- Execution flow
- Memory that devoted for data becomes set of instructions

# Counteraction

- Canaries
- $W$  xor  $X$  memory
- Address space layout randomization (ASLR)

# Canaries (stack guard)



Prolog and epilog of function

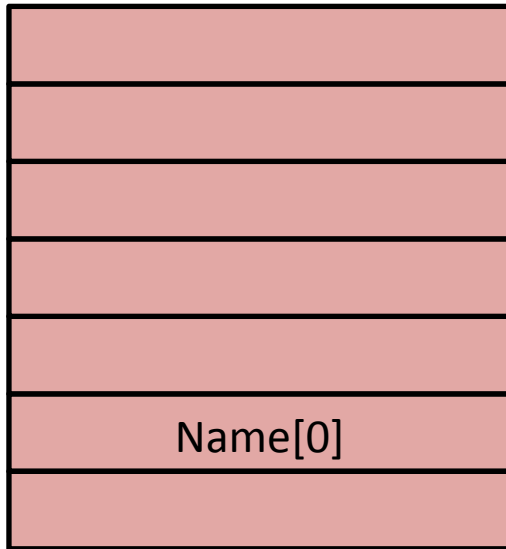
# ASLR

First

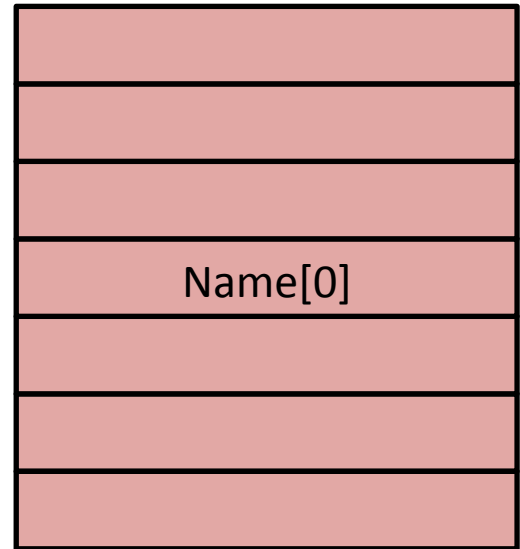
Second

0x7f..ff

7cb7ba740



ef5415a90



```
char name[64];  
printf("%p\n", name);  
puts("What's your name?");  
gets(name);  
printf("Hello, %s!\n", name);
```



# W xor X memory

- Hardware or software implementation
- Memory protected from execution
- Prevents payload uploading

# Counteraction

- Canaries
- $W \text{ xor } X$  memory
- Address space layout randomization (ASLR)

# Syuzhet

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# W xor X memory

- Hardware support: AMD64, ARM,
- NX bit
- Disable execution
- Requires support by kernel and environment.

# KE 1: memory objects

- Entities:
  - Dataspace
  - Region mapper
- Semantic:

```
1 L4Re::Env::env()->mem_alloc()->alloc(size, ds, L4Re::Mem_alloc::Executable)
```

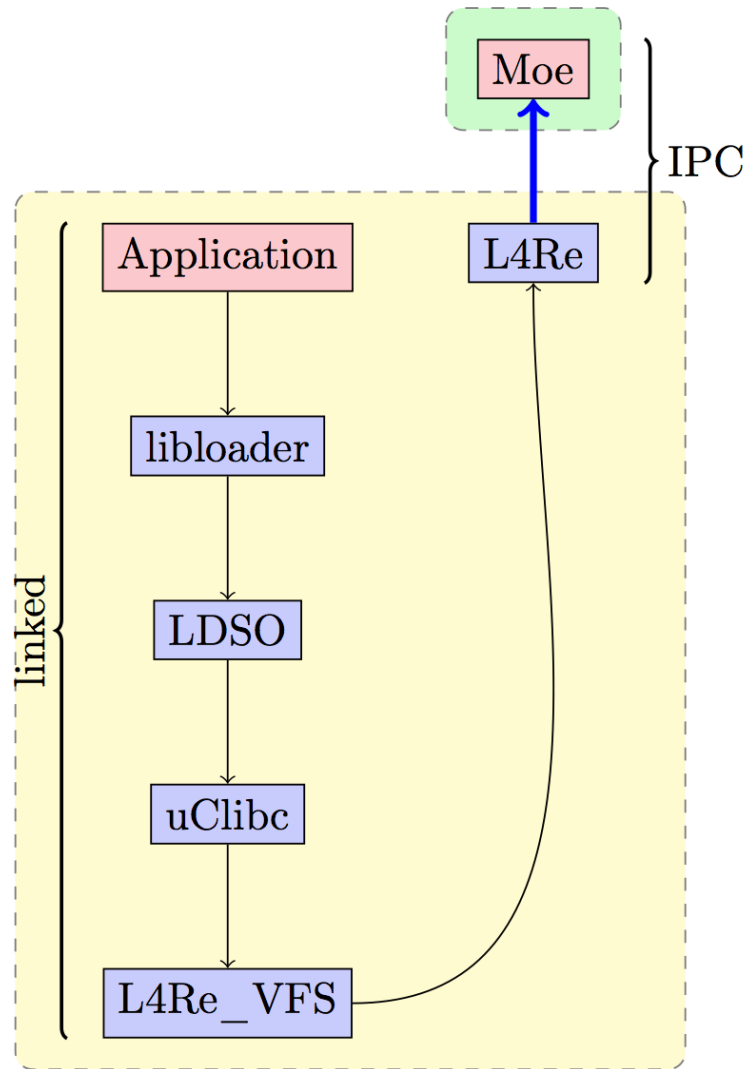
```
1 L4Re::Env::env()->rm()->attach(&ptr, size,  
2     L4Re::Rm::Search_addr | L4Re::Rm::Eager_map  
3     | L4Re::Rm::Executable, ds)
```

# KE 2: Starting, ELF

- ELF file contains sections with access flags
- Elf-loader creates region according to ELF

1	PHDR	off	0x0000000000000040	vaddr	0x0000000001000040	paddr			
			0x0000000001000040	align	2**3				
2		filesz	0x0000000000000118	memsz	0x0000000000000118	flags	r--		
3	LOAD	off	0x0000000000000000	vaddr	0x0000000001000000	paddr			
			0x0000000001000000	align	2**12				
4		filesz	0x000000000001cbe8	memsz	0x000000000001cbe8	flags	r-x		
5	LOAD	off	0x000000000001d000	vaddr	0x000000000101d000	paddr			
			0x000000000101d000	align	2**12				
6		filesz	0x000000000000370	memsz	0x00000000000143e0	flags	rw-		
7	TLS	off	0x000000000001d000	vaddr	0x000000000101d000	paddr			
			0x000000000101d000	align	2**3				
8		filesz	0x0000000000000000	memsz	0x0000000000000018	flags	rw-		
9	0x60000014	off	0x00000000000147a0	vaddr	0x00000000010147a0	paddr			
			0x00000000010147a0	align	2**4				
10		filesz	0x0000000000000018	memsz	0x0000000000000018	flags	r--		

# KE 2: Starting



# KE 3: The Gentleman's Set of Tests

Tect	L4Re	L4Re + NX dataspace
anonmap	Vulnerable	Killed
execbss	Vulnerable	Killed
execdata	Vulnerable	Killed
execstack	Vulnerable	Killed
mprotanon	Vulnerable	Killed
mprotbss	Vulnerable	Killed
mprotdata	Vulnerable	Killed
mprotheap	Vulnerable	Killed
mprotstack	Vulnerable	Killed



# Restrictions

- For well protection all techniques should be used
  - W xor X
  - ASLR
  - Canaries
  - Other...
- I386 does not have a hardware NX

# Restrictions: L4Linux

- L4Linux uses low level Fiasco.OC calls
- L4Linux starts program self
- Obtain one big dataspace from kernel.
- L4Linux has to manage W xor X allocation self
- Does not support AMD64
- A big hole in security

# Thank you for attention

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\* Please do not fork me on github