Mathematics and development of fast TLS handshakes

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Web content delivery & protection

- 2013: WAF development by request of Positive Technologies "Visionar" from Gartner magic quadrant'15
 - Web attacks
 - L7 HTTP/HTTPS DDoS attacks
- Nginx, HAProxy, etc. perfect HTTP proxies, not HTTP filters
- Netfilter works in TCP/IP stack (softirg) => HTTP(S)/TCP/IP stack
- Tempesta FW:
 - hybrid of HTTP accelerator & firewall
 - embedded into the Linux TCP/IP stack

	CHALLENGERS
L	
L	04-04
	Crinx
	Akama
	Barracuda Networks 🔵
	NSFOCUS Fortinet
	Penta Security
	Ergon Informatik
	B United Security Provider
	DBAPPSecurity



Gartner



As of July 2015



Tempesta TLS

https://netdevconf.info/0x12/session.html?kernel-tls-handhakes-for-https-ddos-mitigation

- Part of Tempesta FW, an open source Application Delivery Controller
- Open source alternative to F5 BIG-IP or Fortinet ADC
- "TLS CPS/TPS" is a common specification for network security appliances & ADCs





Linux kernel TLS handshaks

- Very fast light-weight Linux kernel implementation
 - …even for session resumption
 - there is modern research in the field
- Resistant against DDoS on TLS handshakes (asymmetric DDoS)
- Privileged address space for sensitive security data
 - Varnish: TLS is processed in separate process Hitch http://varnish-cache.org/docs/trunk/phk/ssl.html
 - Resistance against attacks like CloudBleed https://blog.cloudflare.com/incident-report-on-memory-leak-caused-by-cloudflare-parser-bug/



Why NIST p256?

- ECDSA
 - RSA and NIST curves p256, p384, and p521 are the only allowed for CA certificates https://cabforum.org/baseline-requirements-documents/
 - P256 is the fastest NIST curve
 - P521 isn't recommended by IANA https://www.iana.org/assignments/tls-parameters/tls-parameters.xml#tls-parameters-8
 - RSA is slow and vulnerable to asymmetric DDoS https://vincent.bernat.im/en/blog/2011-ssl-dos-mitigation
- Curve25519
 - Much faster than NIST p256
 - In practice for ECDHE only



TLS libraries performance issues

- Copies, memory initializations/erasing, memory comparisons
- memcpy(), memset(), memcmp() and their constant-time analogs
- Many dynamic allocations
- Large data structures
- Some math is outdated

```
12.79%
        libc-2.24.so
 9.34%
        nqinx(openssl)
 7.40%
        nginx(openssl)
 3.54%
        nqinx(openssl)
 2.87%
        nqinx(openssl)
 2.79%
        libc-2.24.so
 2.49%
        libc-2.24.so
 2.30%
        nginx(openssl)
 1.82%
        libc-2.24.so
                             malloc
 1.57%
        [kernel.kallsyms]
 1.45%
        libc-2.24.so
                             free
 1.32%
        nginx(openssl)
 1.18%
        nqinx(openssl)
 1.12%
        nginx(openssl)
 0.93%
        libc-2.24.so
        nginx(openssl)
 0.81%
 0.75%
        libc-2.24.so
 0.57%
        nqinx(openssl)
 0.54%
        nginx(openssl)
 0.54%
        nginx(openssl)
```

0.50% [kernel.kallsyms]



_int_malloc __ecp_nistz256_mul_montx __ecp_nistz256_sqr_montx sha256_block_data_order_avx2 ecp_nistz256_avx2_gather_w7 int free malloc_consolidate OPENSSL cleanse

do_syscall_64

ecp_nistz256_ord_sqr_montx ecp_nistz256_point_doublex __ecp_nistz256_sub_fromx _memmove_avx_unaligned_erms _ecp_nistz256_mul_by_2x __memset_avx2_unaligned_erms aesni_ecb_encrypt ecp_nistz256_point_addx EVP MD CTX reset entry_SYSCALL_64



The source code

https://github.com/tempesta-tech/tempesta/tree/master/tls

- Still in-progress: we implement some of the algorithms on our own
- Initially the fork of mbed TLS 2.8.0 (https://tls.mbed.org/) x40 faster!
 - very portable and easy to move into the kernel
 - cutting edge security
 - too many memory allocations (https://github.com/tempesta-tech/tempesta/issues/614)
 - big integer abstractions (https://github.com/tempesta-tech/tempesta/issues/1064)
 - inefficient algorithms, no architecture-specific implementations, ...
- We also take parts from WolfSSL (https://github.com/wolfSSL/wolfssl/)
 - very fast, but not portable
 - **Security** *https://github.com/wolfSSL/wolfssl/issues/3184*





ECDSA & ECDHE mathematics: Tempesta TLS, OpenSSL, WolfSSL

OpenSSL 1.1.1h

256 bits ecdsa (nistp256) 36473 sign/s 256 bits ecdh (nistp256) 16620 op/s

WolfSSL (current master)

ECDSA	256	sign
ECDHE	256	agree

43260 ops/sec (+19%) 40878 ops/sec (+146%)

Tempesta TLS (full TLS handshake operation)

ECDSA sign (nistp256): ops/s=38393 ECDHE srv (nistp256): ops/s=13418

OpenSSL & WolfSSL don't include ephemeral keys generation (one more *m* * *G* operation)

Demo!

- Tempesta TLS, Nginx-1.14.2/OpenSSL-1.1.1d, Nginx-1.17.8/WolfSSL
- ► TLS 1.2
 - full handshakes
 - abbreviated handshakes
- tls-perf

https://github.com/tempesta-tech/tls-perf

- establish & drop many TLS connections in parallel
- like TLS-THC-DOS, but faster, more flexible, more options



Data for proprietary vendors

- BIG-IP is only 30-50% faster than Nginx/OpenSSL/DPDK https://www.youtube.com/watch?v=Plv87h8GtLc
- Avi Vantage (VMware) makes ~2000 handshakes/second per 1CPU https://avinetworks.com/docs/latest/ssl-performance/



Why faster?

- No memory allocations in run time
- No context switches
- No copies on socket I/O
- Less message queues
- Zero-copy handshakes state machine https://netdevconf.info/0x12/session.html?kernel-tls-handhakes-for-https-ddos-mitigation
- State of the art cryptography mathematics



Elliptic curve cryptography

• Secp256r1: $y^2 = x^3 - 3x + b$ defined over the **field** GF(p) $p256 = 2^{256} - 2^{224} + 2^{192} + 2^{96} - 1$



Point multiplication

OpenSSL: "Fast prime field elliptic-curve cryptography with 256-bit primes" by Gueron and Krasnov

 $\mathbf{P} = \mathbf{m} + \mathbf{P}$ - the most expensive elliptic curve operation

```
for i in bits(m):
     Q \leftarrow point_double(Q)
     if m_i == 1:
          Q \leftarrow point\_add(Q, P)
```

- Point multiplications in TLS handshake:
 - known point multiplication: precompute the table for doubled G
 - perfect forward secrecy ECDHE: generate keys G * d (d random)
 - handshake: 2 known & 1 unknown point multiplications



Point representation and coordinate systems

http://www.hyperelliptic.org/EFD/g1p/auto-shortw-jacobian.html

"Analysis and optimization of elliptic-curve single-scalar multiplication", Bernstein & Lange, 2007

- Jacobian coordinates (rough estimations)
 - conversion overhead: 39*M + 4*S + 3*I (for w(-indow) = 4)
 - point addition (mixed) $8 \times M + 3 \times S$, doubling $2 \times M + 4 \times S$
- Affine coordinates (rough estimations)
 - point addition $13 \times M + 4 \times S$, doubling $4 \times M + 5 \times S$
- NIST 256 bits, D = 256 / w = 64 Comb rounds (addition & doubling):

 $64 \times (10 \times M + 7 \times S) < < 64 \times (17 \times M + 9 \times S)$





Point addition

http://www.hyperelliptic.org/EFD/g1p/auto-shortw-jacobian.html#addition-add-2007-bl

• ex. addition in Jacobian coordinates (cost: 11M + 5S) $A = (x_1, y_1, z_1), B = (x_2, y_2, z_2), \text{ then } C = A + B = (x_3, y_3, z_3) \text{ is}$

$$\begin{array}{l} U_1 \,=\, X_1 Z_2^2 \\ U_2 \,=\, X_2 Z_1^2 \\ S_1 \,=\, Y_1 Z_2^3 \\ S_2 \,=\, Y_2 Z_1^3 \\ H \,=\, U_2 - \, U_1 \\ R \,=\, S_2 - \, S_1 \\ z_3 \,=\, H Z_1 Z_2 \\ x_3 \,=\, R^2 - \, H^3 - \, 2 U_1 H^2 \\ y_3 \,=\, (U_1 H^2 - \, X_3) R - \, S_1 H^3 \end{array}$$



The cost

- ▶ Modular multiplication (M) is the most expensive basic scalar operation
- Modular squaring (s) faster than M, usually 0.8M (Montgomery) (0.9 for optimized FIPS due to more expensive modular reduction)
- Modular inversion (I) is very expensive, about 100м

alar operation omery) eduction)



Modular arithmetics

- ex. prime field F(29) 17 + 20 = 8 since $37 \mod 29 = 8$ addition: 17 - 20 = 26 since $-3 \mod 29 = 26$ subtraction: 17 * 20 = 21 since 340 mod 29 = 21 multiplication: since $17 \cdot 12 \mod 29 = 1$ $17^{-1} = 12$ inversion:
- Montgomery reduction (the most used)
 - there is some overhead, but each modular operation is cheaper

FIPS reduction

- Can be faster if small number of modular operations is used
- There are optimization techniques, e.g. "Low-Latency Elliptic Curve Scalar Multiplication" Bos, 2012
- But still about 65% slower than Montgomery reduction



Montgomery multiplication in P256

"Montgomery Multiplication", Henry S. Warren, Jr.

- Fast 256-bit integer multiplication with modular reduction on P256
- ▶ a, b < m (m modulus P256)</p>
- ► Set n = 2²⁵⁶
- Transform multipliers to Montgomery domain (overhead): $a' = an \mod m$ $b' = bn \mod m$
- Fast multiplication with reduction: $u = a'b'/n \mod m$
 - compute only 256 bits of $(a'b' + (-m^{-1}a'b' \mod n)m)/n$
 - if u > m, then $u \leftarrow u m$ (unconditionally, carry as a mask)
- Convert to ordinary number: $v = un^{-1} \mod m$



The math layers

- Different multiplication algorithms for fixed and unknown point
 - "Efficient Fixed Base Exponentiation and Scalar Multiplication based on a Multiplicative Splitting Exponent Recoding", Robert et al 2019
- Point doubling and addition everyone seems use the same algorithms
- Jacobian coordinates: different modular inversion algorithms
 - "Fast constant-time gcd computation and modular inversion", Bernstein et al, 2019
- Modular reduction for scalar multiplications:
 - Montgomery has overhead vs FIPS speed: if we use less multiplcations it could make • sense to use different reduction method FIPS (seems deadend)
 - "Low-Latency Elliptic Curve Scalar Multiplication" Bos, 2012
- => Balance between all the layers



Example of math layers balancing

- For w=5 we need 52 point additions for an unknown point multiplication
- Jacobian coordinates addition takes 11M + 5S
- Affine-Jacobian coordinates addition takes 8M + 3S
 - about 4.4M cheaper if S = 0.8M
 - requires 2 coordinates normalizations for Comba precomputation
 - coordinates normalization: $2^{w-1} * (6M + 1S) + 1I$
- Almost the same for S = 0.9M and fast inversion I < 100M
- Montgomery arithmetics (S = 0.8M):
 - ECDHE +28% and ECDSA +6% performance





Side channel attacks (SCA) resistance

- Timing attacks, simple power analysis, differential power analysis etc.
- Protections against SCAs:
 - Constant time algorithms
 - Dummy operations
 - Point randomization
 - e.g. modular inversion **741 vs 266** iterations
- RDRAND allows to write faster non-constant time algorithms
 - SRBDS mitigation costs about 97% performance https://software.intel.com/security-software-guidance/insights/processors-affected-special-register-buffe r-data-sampling?fbclid=IwAR1ifj3ZuAtNOabKkj3vFItBLSvOnMqlxH2I-QeN5KB-aji54J1BCJa9ILk https://www.phoronix.com/scan.php?page=news_item&px=RdRand-3-Percent&fbclid=IwAR2vmmR_Lir oekUuw7KMRaHB7KThpqz0tIr1fX2GCW3HAvwt5Kb1p9xpLKo



Memory usage & SCA

• ex. ECDSA precomputed table for fixed point multiplication

- mbed TLS: ~8KB dynamically precomputed table, point randomization, constant-time algorithm, full table scan
- OpenSSL: ~150KB static table, full scan
- WolfSSL: ~150KB, direct index access (fixed in the new version) https://github.com/wolfSSL/wolfssl/issues/3184

=> 150KB is far larger than L1d cache size, so many cache misses:



Big Integers (aka MPIs)

"BigNum Math: Implementing Cryptographic Multiple Precision Arithmetic", by Tom St Denis

- All the libraries use them (not in hot paths), mbed TLS overuses them
- > linux/lib/mpi/, linux/include/linux/mpi.h

```
typedef unsigned long int mpi_limb_t;
struct gcry_mpi {
 int nlimbs; /* number of valid limbs */
 unsigned flags;
 mpi_limb_t *d; /* array with the limbs */
};
```

Need to manage variable-size integers => size-specific assembly implemetations





Easy assembly

// a := a + b // x[0] is the less significant limb, // x[1] is the most significant limb. void s_mp_add(unsigned long *a, unsigned long *b) { unsigned long carry; a[0] += b[0]; carry = (a[0] < b[0]);a[1] += b[1] + carry; }

// Pointer to **a** is in %RDI, pointer to **b** is in %RSI (%rdi), %r8 movq <mark>8(</mark>%rdi), %r9 movq addq (%rsi), %r8 // add with carry 8(%rsi), %r9 // use the carry in the next addition addc (%r8), (%rdi) movq (%r9), <mark>8</mark>(%rdi) movq







Open questions and further research

- Ice Lake CPUs have negligeable downclocking on AVX-512 https://travisdowns.github.io/blog/2020/08/19/icl-avx512-freq.html
- Parallel Montgomery computations J.W.Bos, "Montgomery Arithmetic from a Software Perspective", 2017
 - SIMD multiplications & squarings of two and more products
 - Interleaved Montgomery multiplications
- Better methods for point multiplications





Going to the Linux kernel upstream

- Details and the discussion
 - https://netdevconf.info/0x14/session.html?talk-performance-study-of-kernel-TLS-handshakes •
 - https://github.com/tempesta-tech/tempesta/issues/1433 ullet
- Server-side only
- The full TLS handshake is in softirg (just like TCP)
- Fallback to a user-space TLS library on ClientHello
- Batches of handshakes in 1 FPU context





TODO

More cryptography mathematics performance optimizations https://github.com/tempesta-tech/tempesta/issues/1064 https://github.com/tempesta-tech/tempesta/issues/1335

TLS 1.3

https://github.com/tempesta-tech/tempesta/issues/1031

- Moving to the kernel asymmetric keys API https://github.com/tempesta-tech/tempesta/issues/1332
- The Linux kernel /crypto API performance issues
 - SHA-256 (crucial for TLS handshake) 30-100% slower than OpenSSL https://github.com/tempesta-tech/tempesta/issues/1483
 - Extra copying and memory allocations in kTLS https://github.com/tempesta-tech/tempesta/issues/1064



Netdev papers about Tempesta TLS

- "Kernel HTTP/TCP/IP stack for HTTP DDoS mitigation", Netdev 2.1, https://netdevconf.info/2.1/session.html?krizhanovsky
- "Kernel TLS handshakes for HTTPS DDoS mitigation", Netdev 0x12, https://netdevconf.info/0x12/session.html?kernel-tls-handhakes-for-https-ddos-mitigation
- Performance study of kernel TLS handshakes", Netdev 014, https://netdevconf.info/0x14/session.html?talk-performance-study-of-kernel-TLS-handsh akes







Contact us if you're interested in fast Linux kernel TLS!

https://github.com/tempesta-tech/tempesta

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