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Quick intro

<u>Linphone</u>

- VOIP SIP client around since 2001
- Is available on GNU/Linux, android, iOS, Windows, Mac
- Uses SIP standards for audio, video and instant messaging
- Secure group messaging using a Signal protocol derivative

Linphone's team also provides

- Flexisip, an open source SIP Proxy
- A free SIP service sip.linphone.org





- I. Principles of Forward Error Correction
- II. Why the Linphone team chose the Flexible FEC scheme
- III. Flexible FEC scheme algorithm described in RFC 8627
- IV. Implementation challenges
- V. Results





I. PRINCIPLES OF FORWARD ERROR CORRECTION

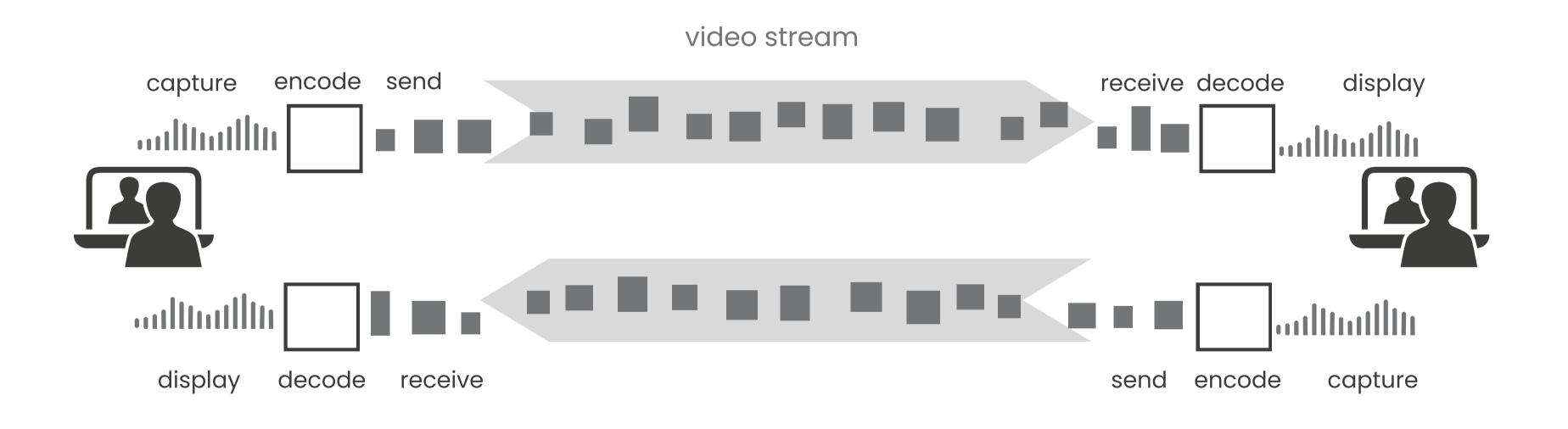




PACKET LOSS IN VIDEO CALL

Real-time Transport Protocol (RTP)

- internet protocol for video/audio transmission
- packets format
- more reliable than UDP

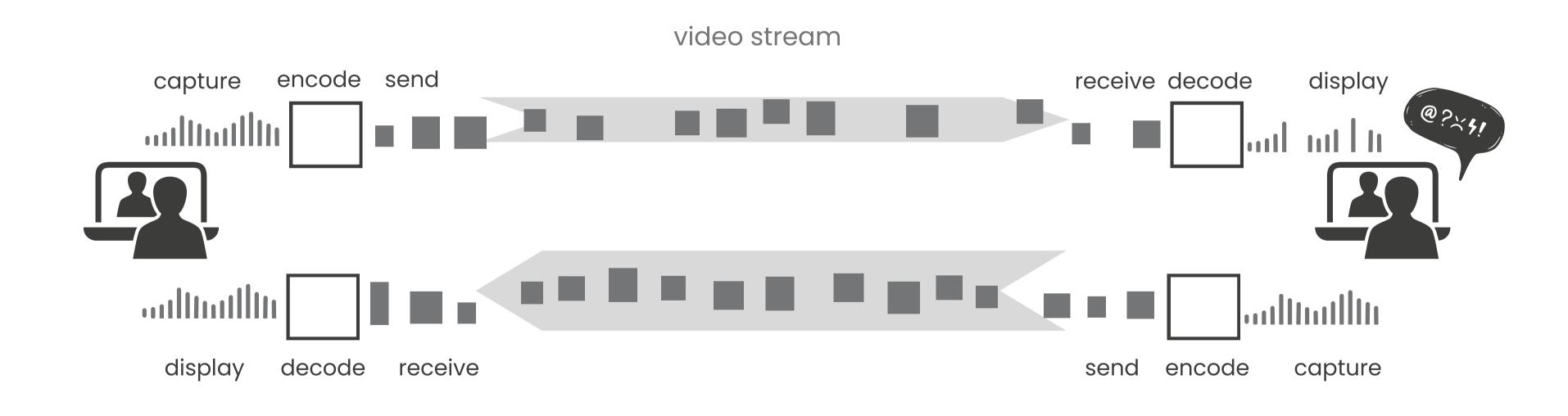




PACKET LOSS IN VIDEO CALL

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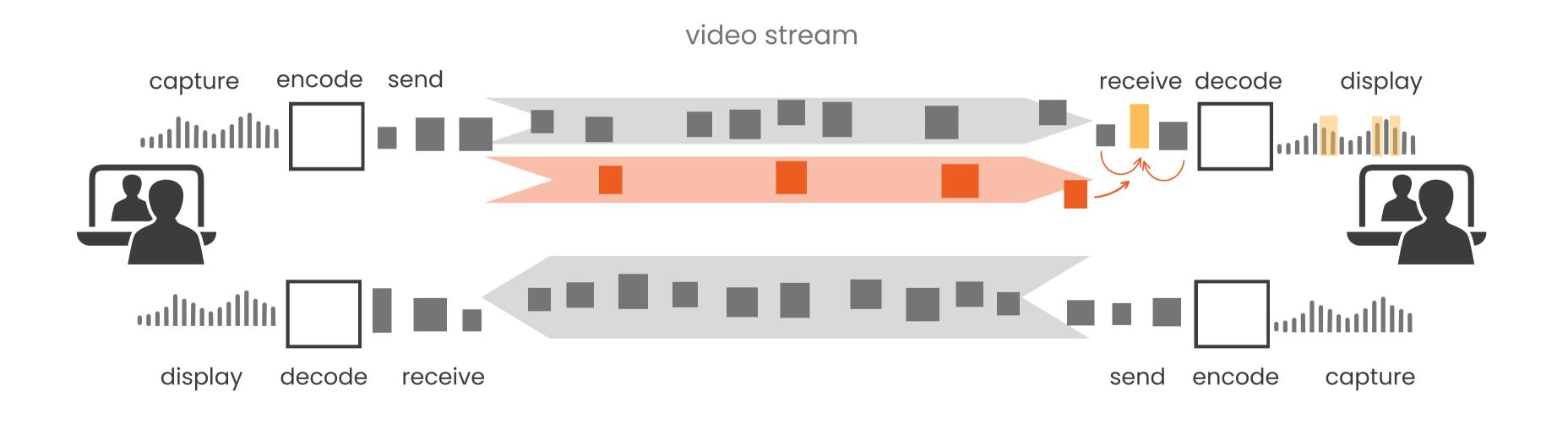




PACKET LOSS IN VIDEO CALL

Real-time Transport Protocol (RTP)

- internet protocol for video/audio transmission
- packets format
- more reliable than UDP





II. WHY THE LINPHONE TEAM CHOSE THE FLEXIBLE FEC SCHEME





STRATEGIES TO RECOVER LOST PACKETS

- Loss detected
 - o ask to send the packet again
 - send preventively the packet twice
 - recover the packet
- Forward Error Correction
 - Low Density Parity Check codes
 - Flexible FEC
 - simplicity: based on packet combinaison with XOR
 - free solution
 - recent standard
 - interoperability with webrtc





RFC 8627

Internet Engineering Task Force (IETF)

Request for Comments: 8627 Category: Standards Track

ISSN: 2070-1721

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Cisco
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callstats.io
A. Begen
Networked Media
G. Mandyam
Qualcomm Inc.
July 2019

RTP Payload Format for Flexible Forward Error Correction (FEC)

Abstract

This document defines new RTP payload formats for the Forward Error Correction (FEC) packets that are generated by the non-interleaved and interleaved parity codes from source media encapsulated in RTP. These parity codes are systematic codes (Flexible FEC, or "FLEX FEC"), where a number of FEC repair packets are generated from a set of source packets from one or more source RTP streams. These FEC repair packets are sent in a redundancy RTP stream separate from the source RTP stream(s) that carries the source packets. RTP source packets that were lost in transmission can be reconstructed using the source and repair packets that were received. The non-interleaved and interleaved parity codes that are defined in this specification offer a good protection against random and bursty packet losses, respectively, at a cost of complexity. The RTP payload formats that are defined in this document address scalability issues experienced with the earlier specifications and offer several improvements. Due to these changes, the new payload formats are not backward compatible with earlier specifications; however, endpoints that do not implement this specification can still work by simply ignoring the FEC repair packets.

- Full description of how a RTP stream is protected
- RTP payload and header format for FEC packets
- Parity codes for reconstruction
- All media: video, audio, text, application





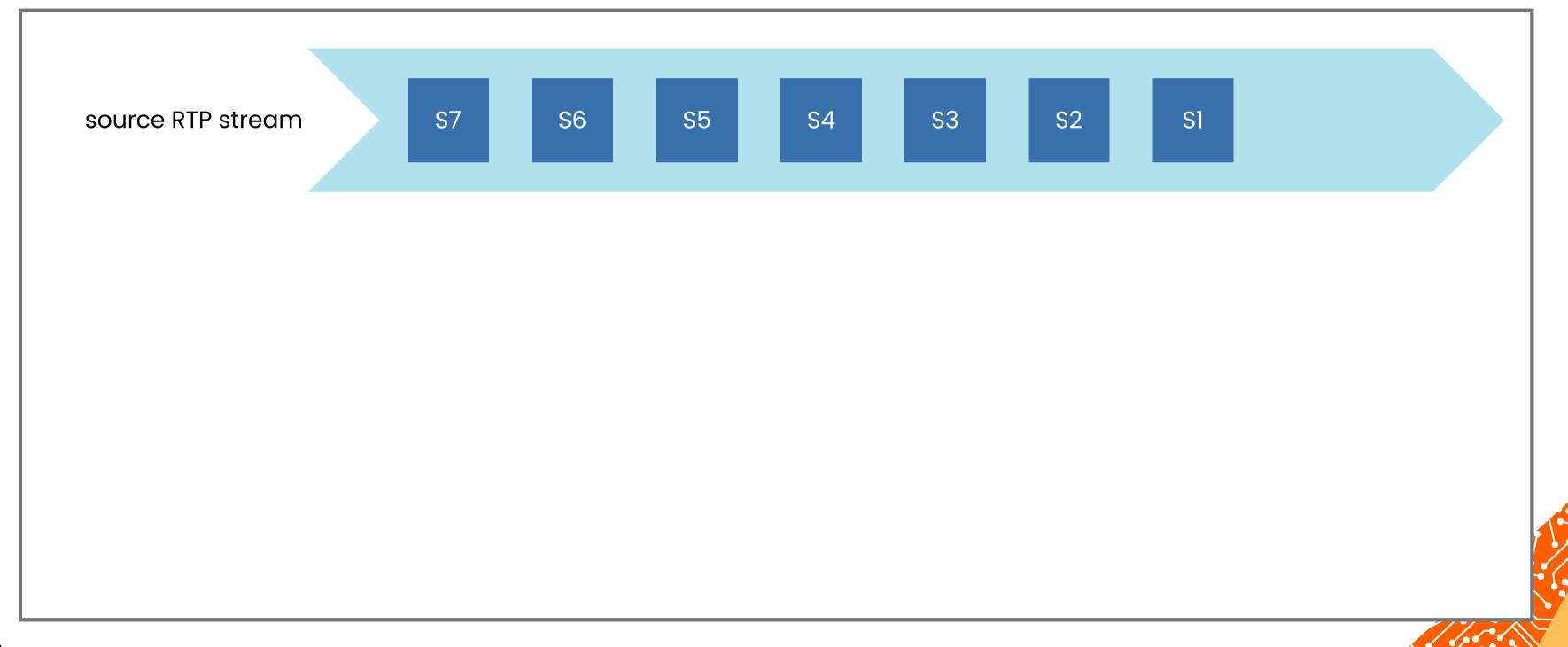
III. FLEXIBLE FEC SCHEME ALGORITHM DESCRIBED IN RFC 8627





FEC PROTECTION

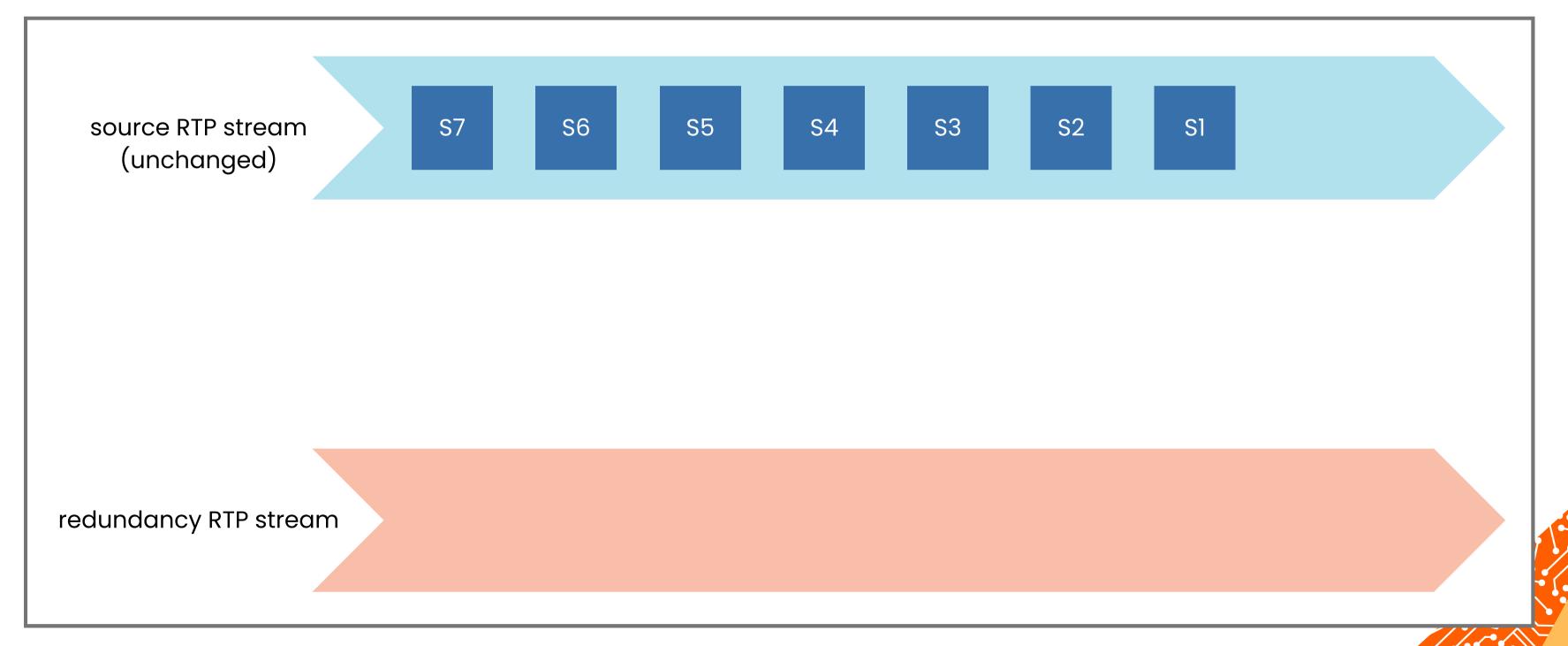






FEC PROTECTION

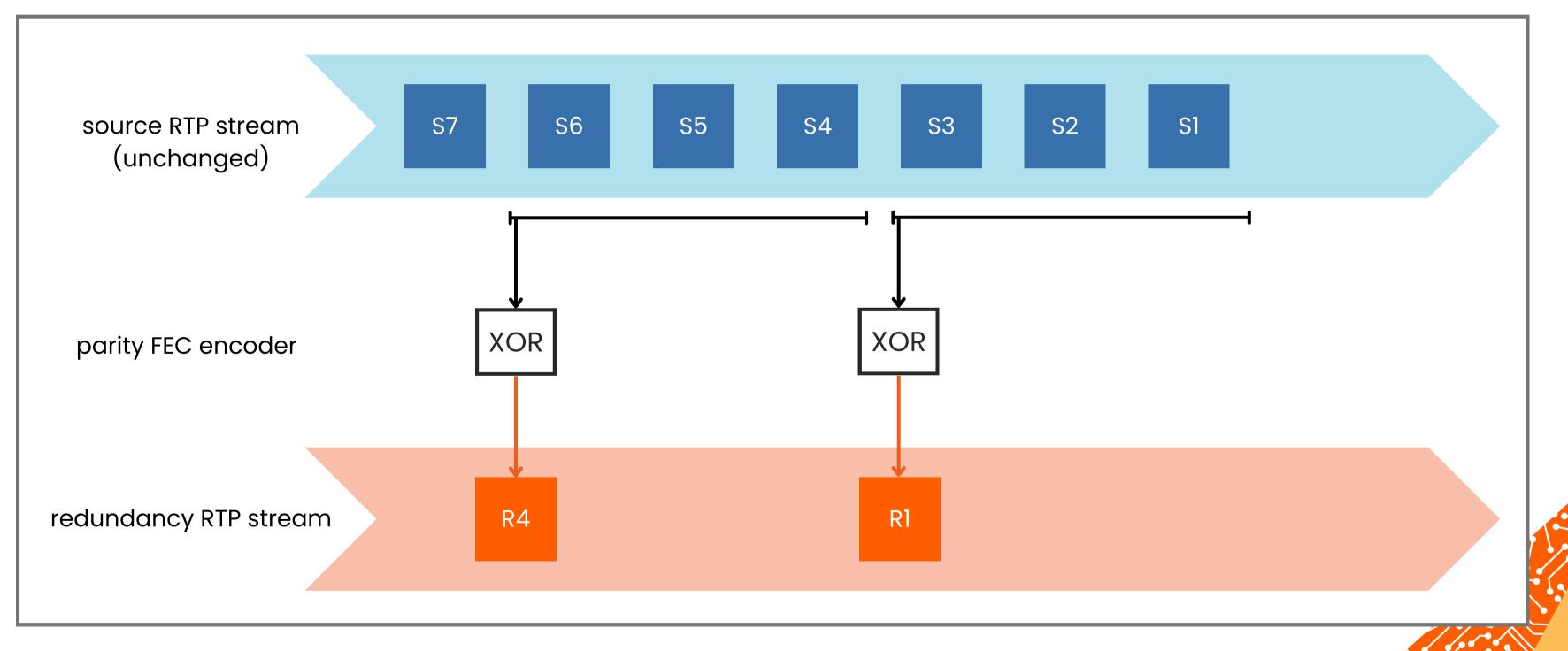






FEC PROTECTION: SENDER

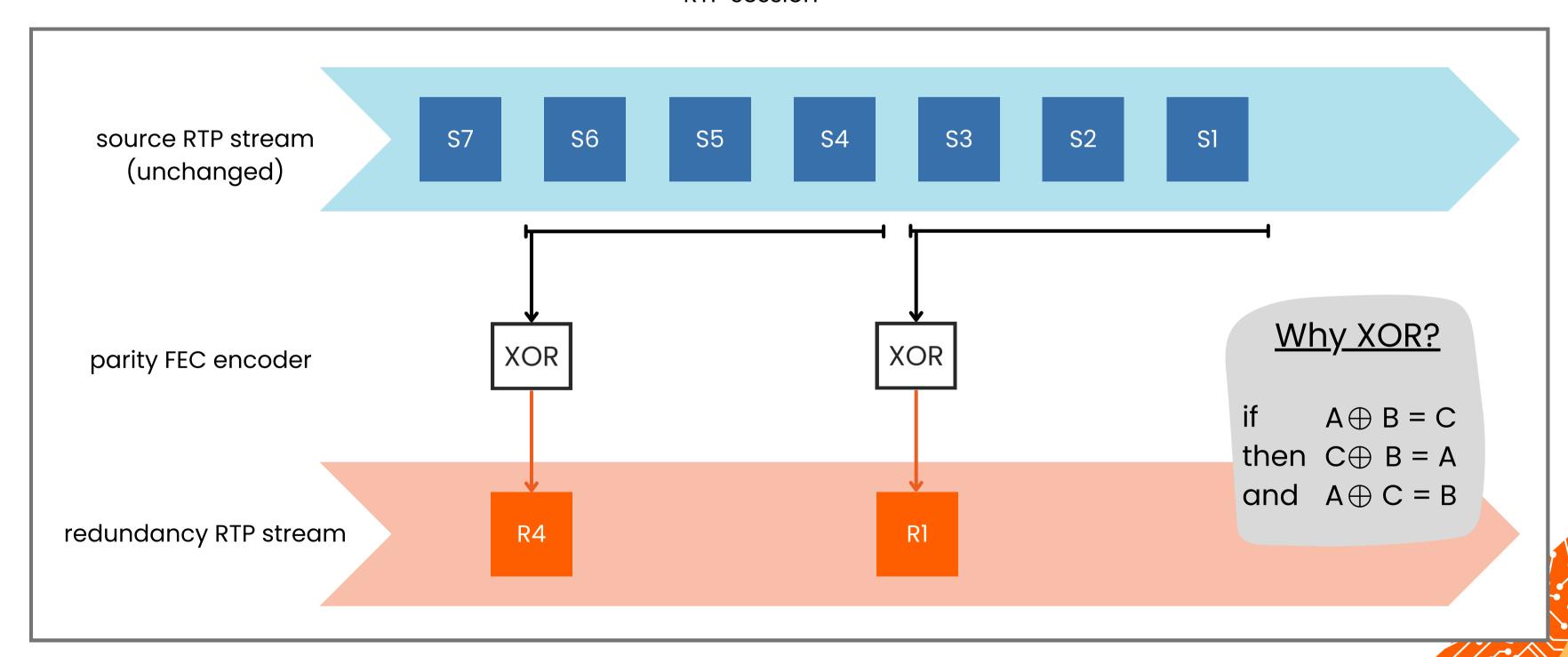
RTP session





FEC PROTECTION: SENDER

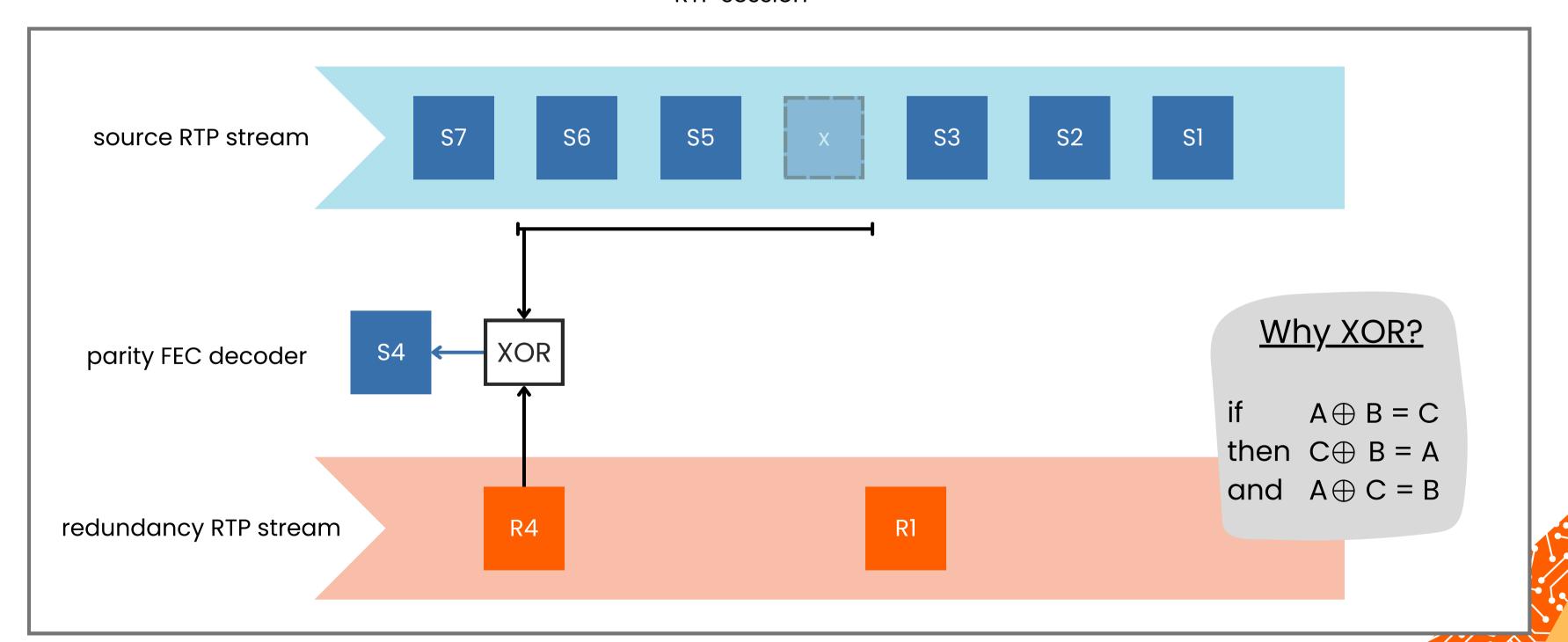
RTP session





FEC PROTECTION: RECEIVER

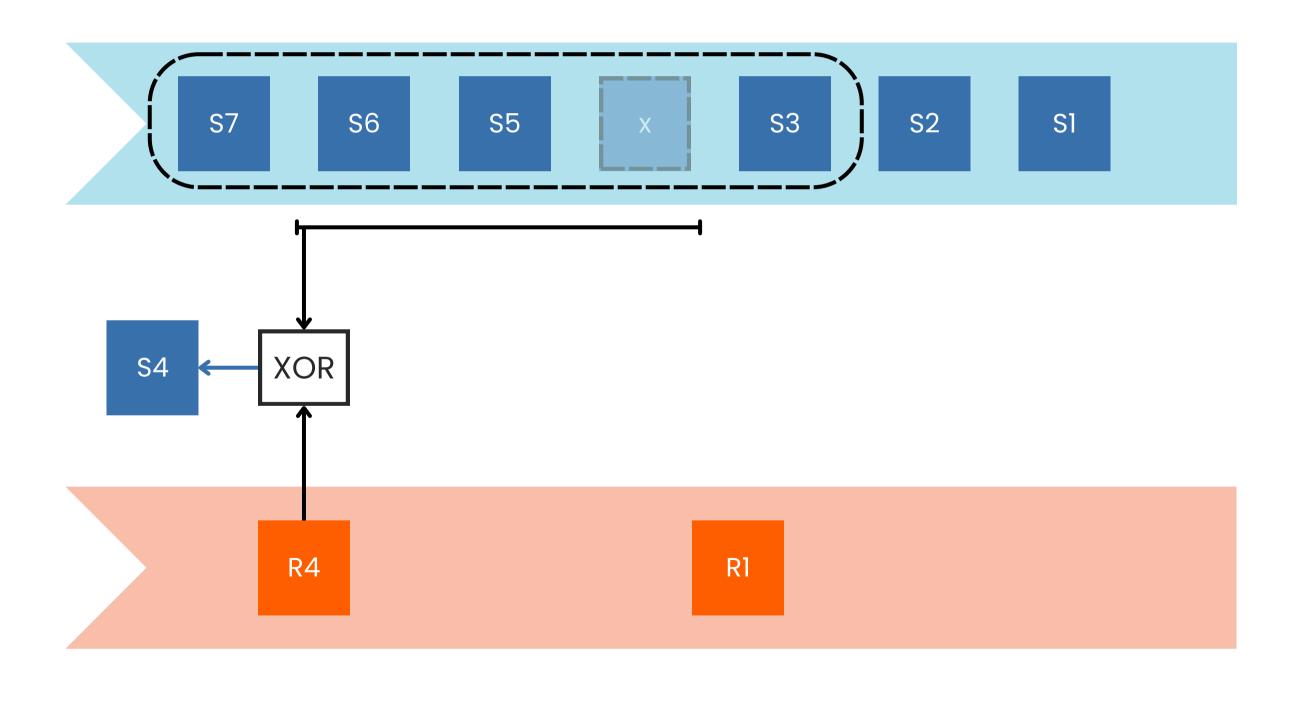
RTP session





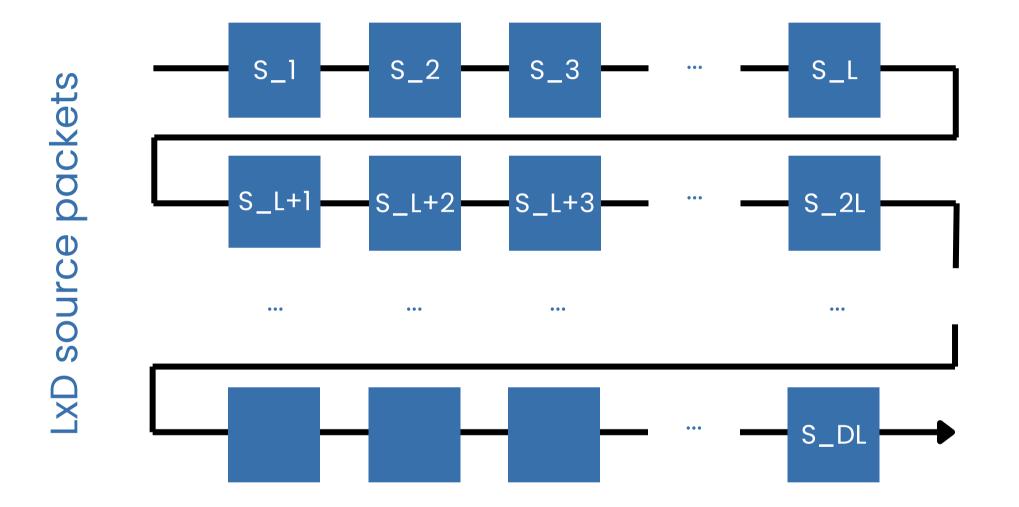
FEC PROTECTION: PARAMETERS

- Repair window
- Protection pattern





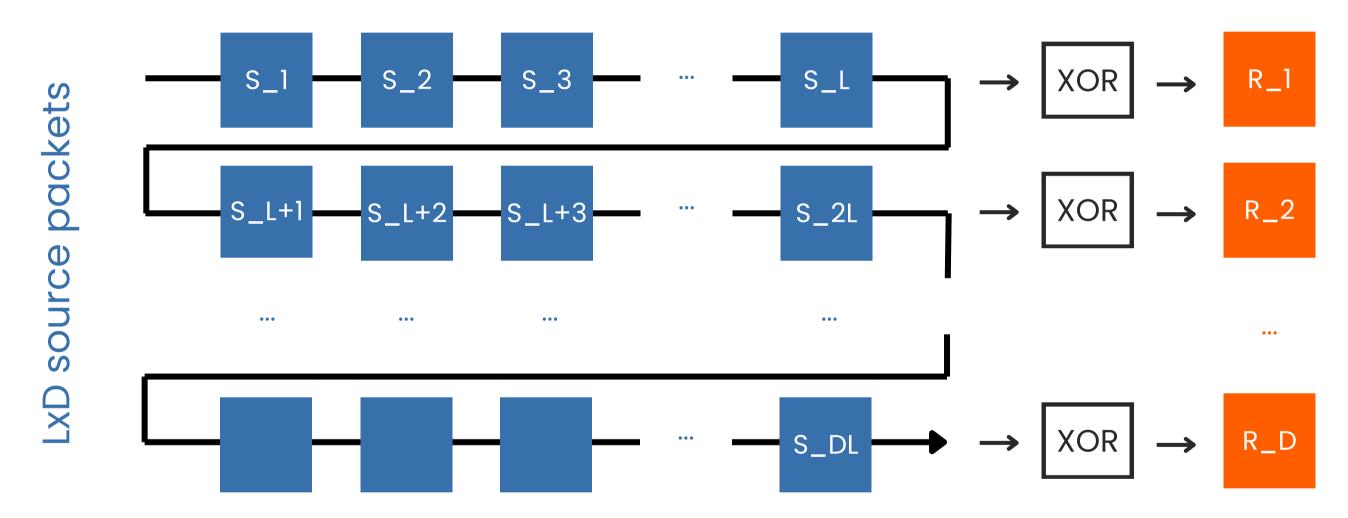
1D NON INTERLEAVED FEC PROTECTION







1D NON INTERLEAVED FEC PROTECTION

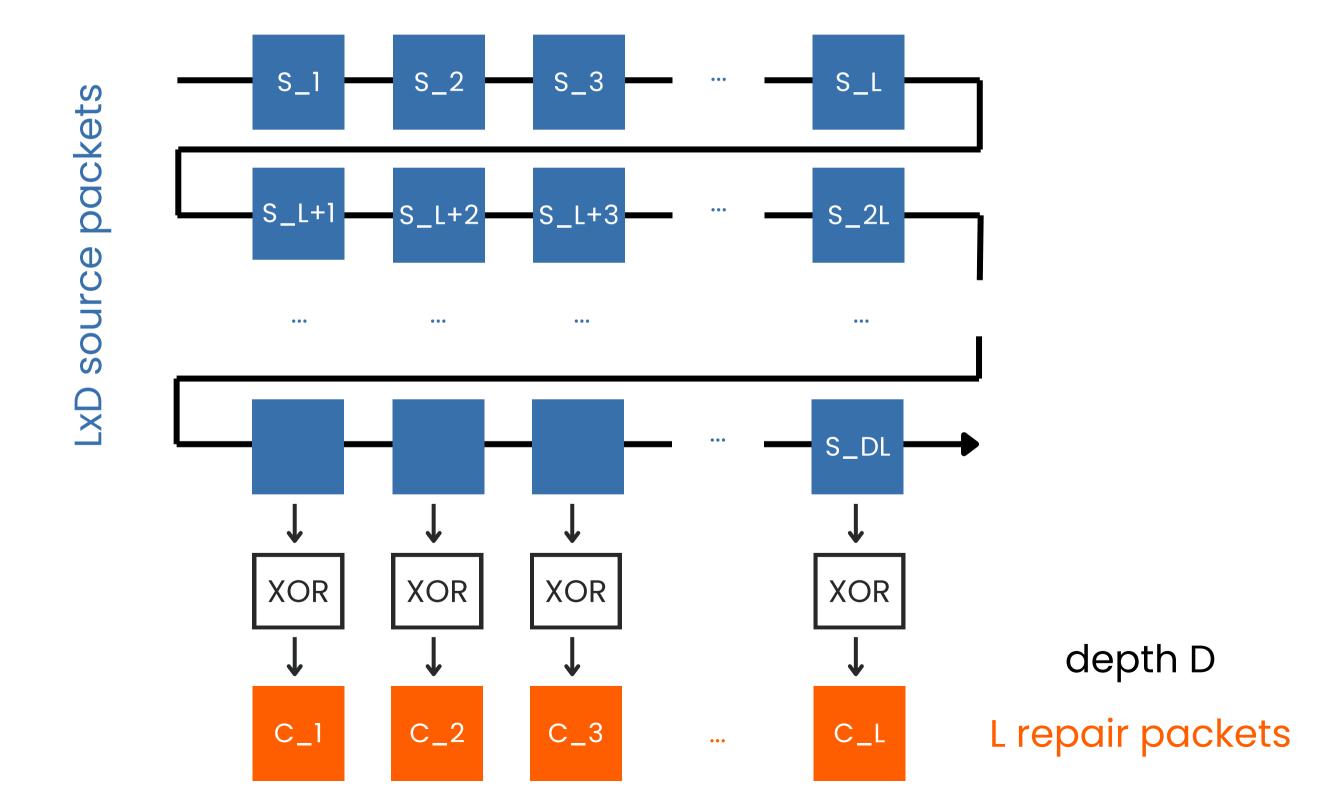


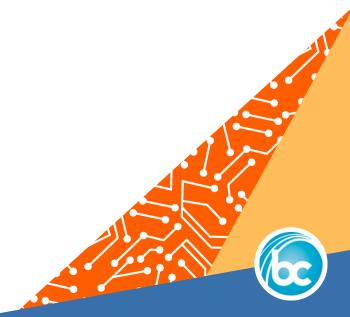
D repair packets length L



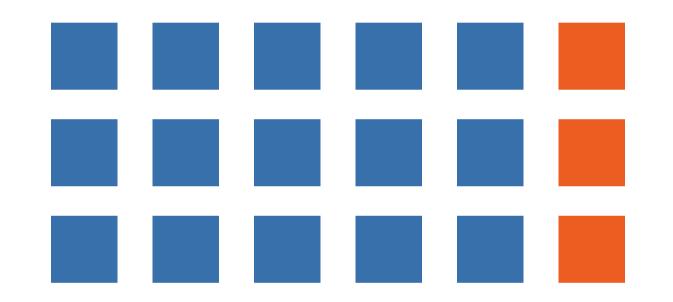


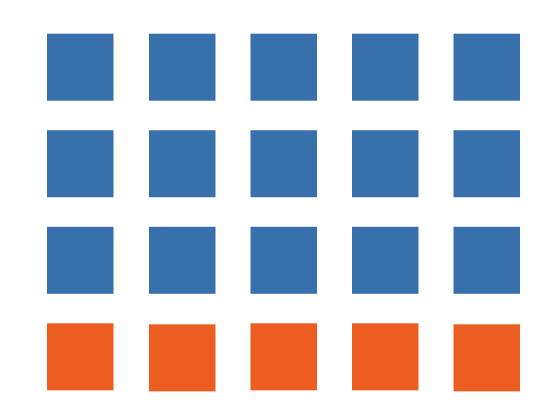
1D INTERLEAVED FEC PROTECTION









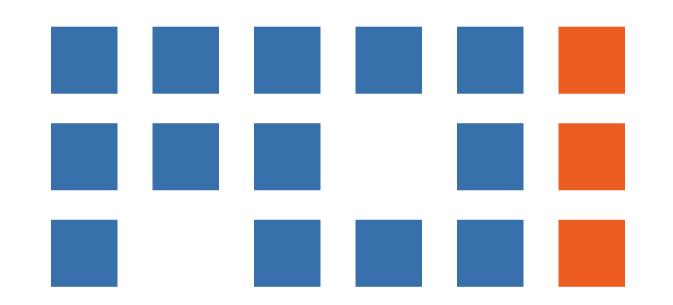


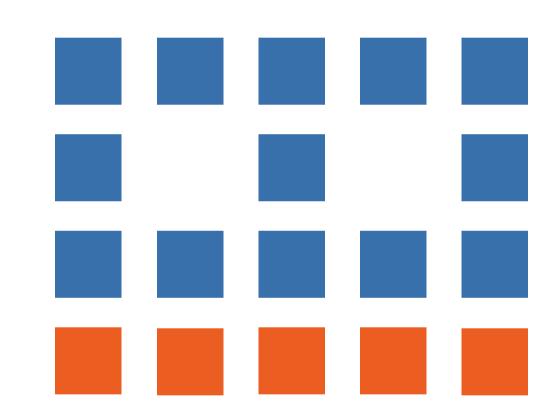
1D non interleaved L = 5 1D interleaved D = 3





random packet loss



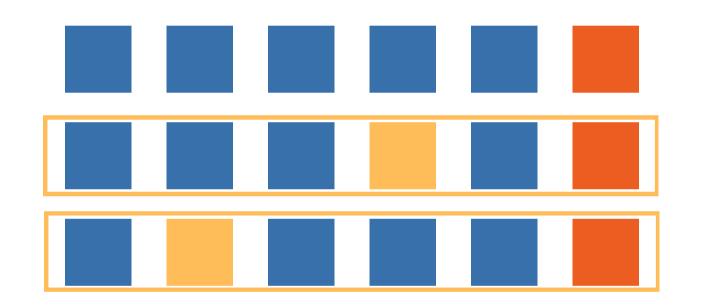


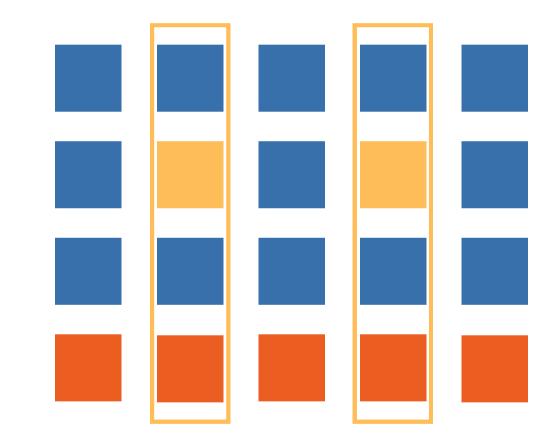
1D non interleaved L = 5 1D interleaved

$$D = 3$$









1D non interleaved

L = 5

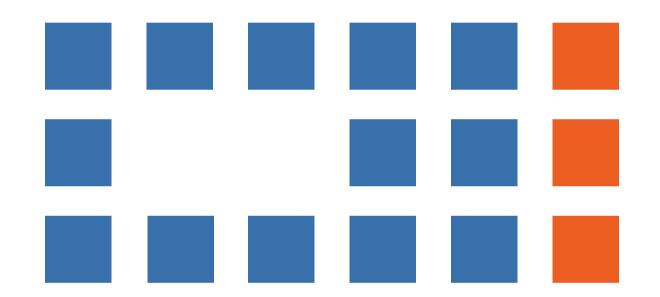
1D interleaved

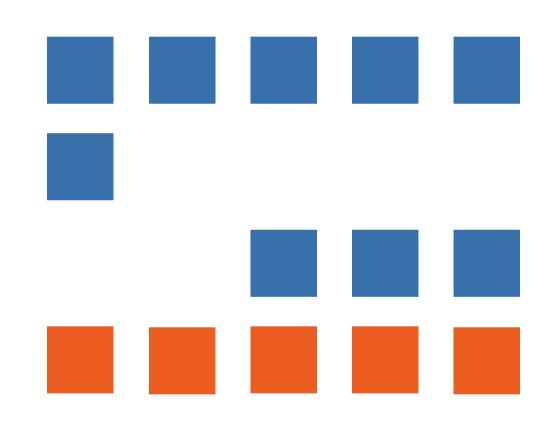
$$D = 3$$





burst: consecutive packet loss



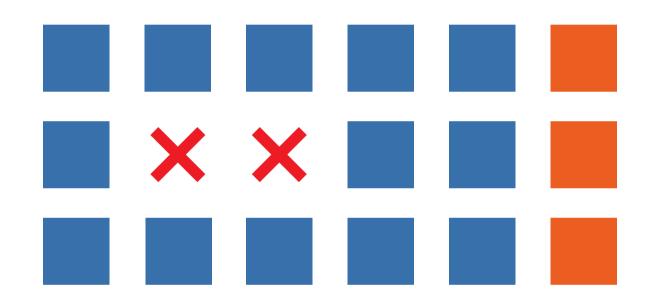


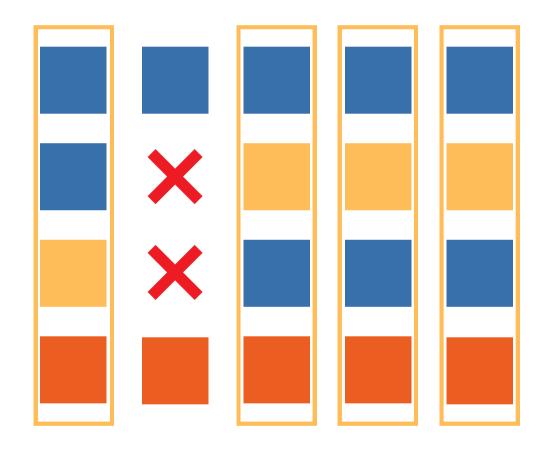
1D non interleaved L = 5 1D interleaved D = 3





burst: consecutive packet loss -> recovery fail





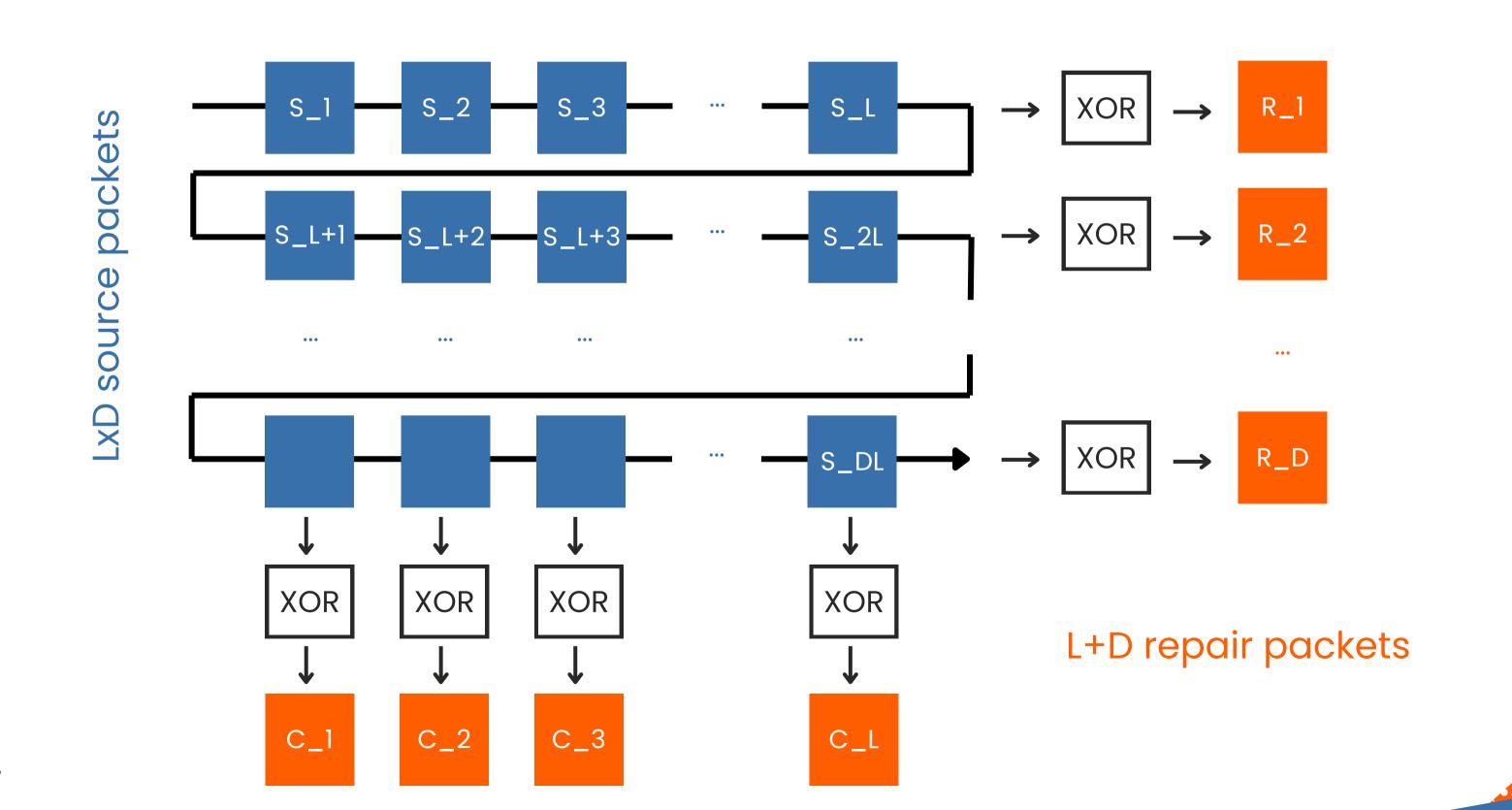
1D non interleaved L = 5 1D interleaved D = 3





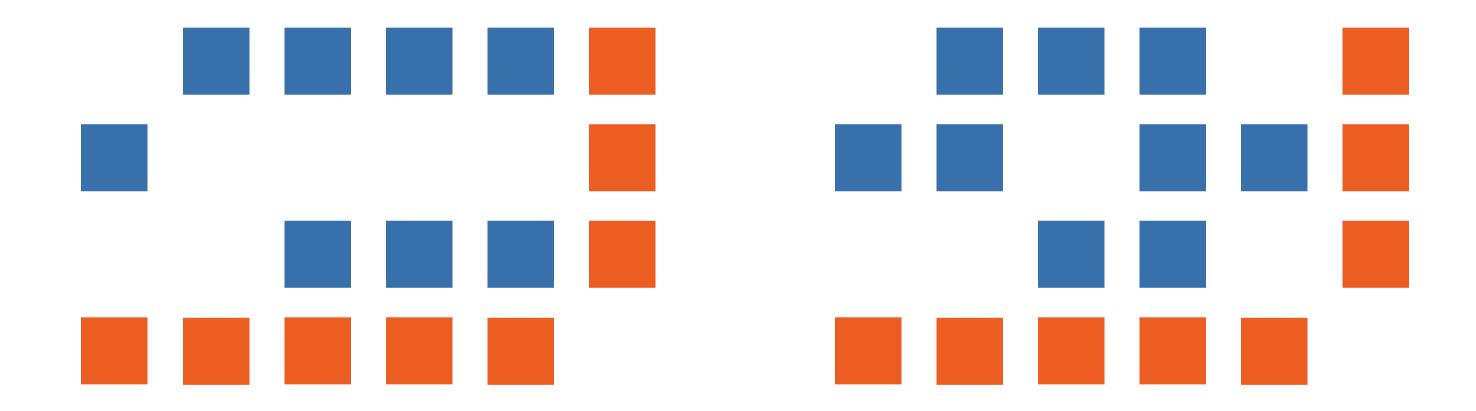
2D PARITY FEC PROTECTION

2D interleaved







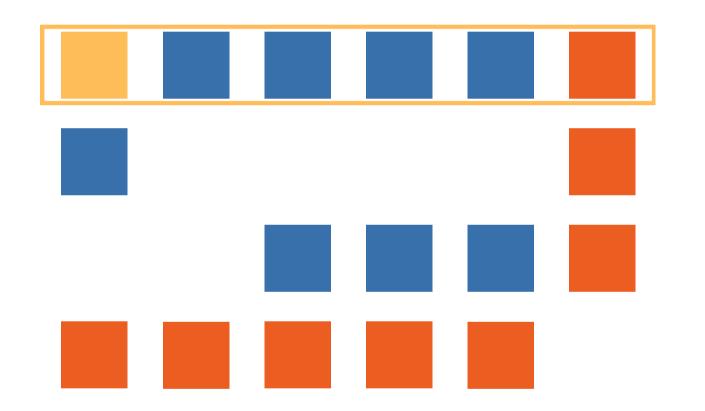


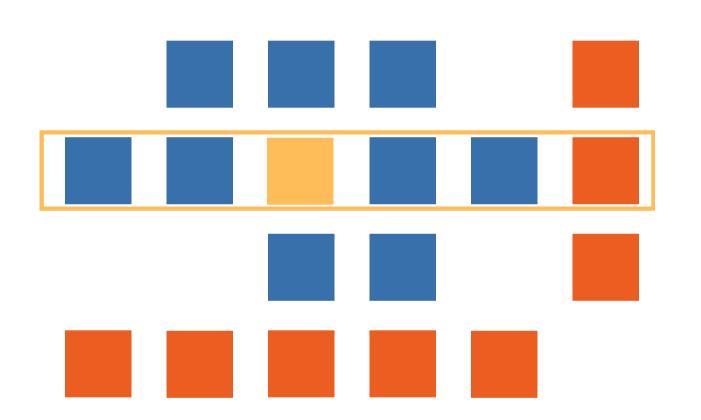
$$L = 5$$

$$D = 3$$







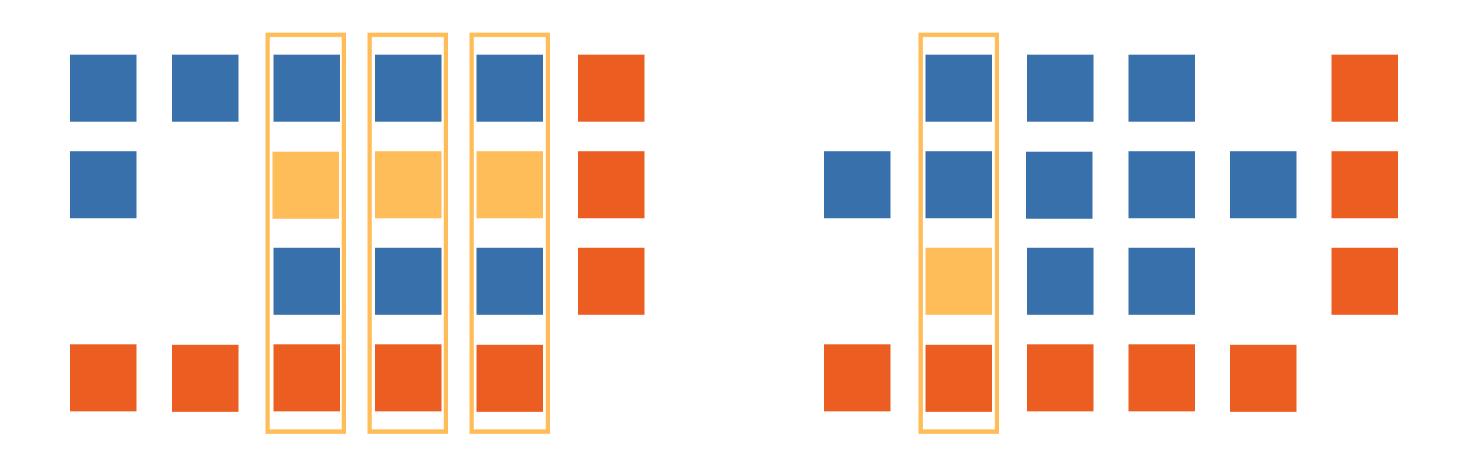


repair rows

$$D = 3$$







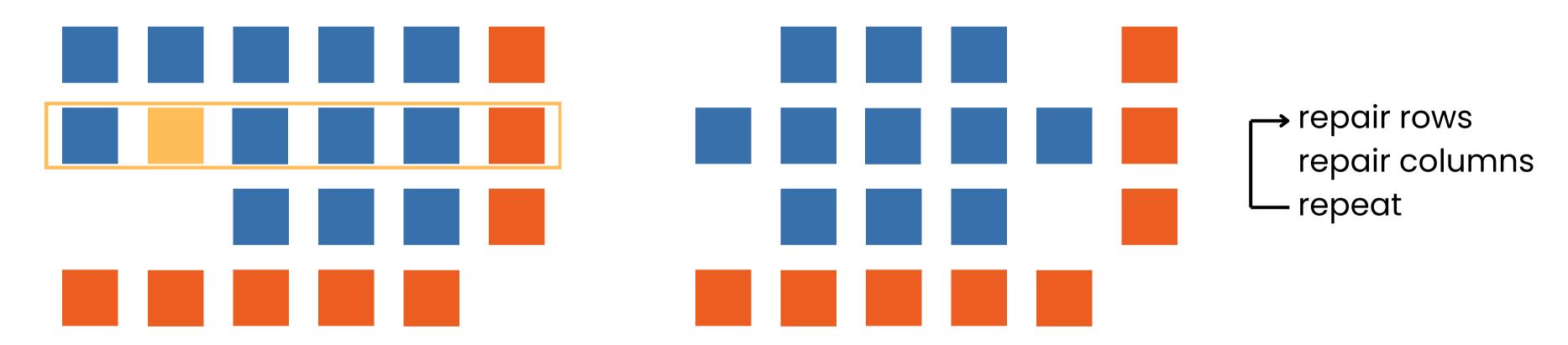
repair rows repair columns

$$L = 5$$

$$D = 3$$





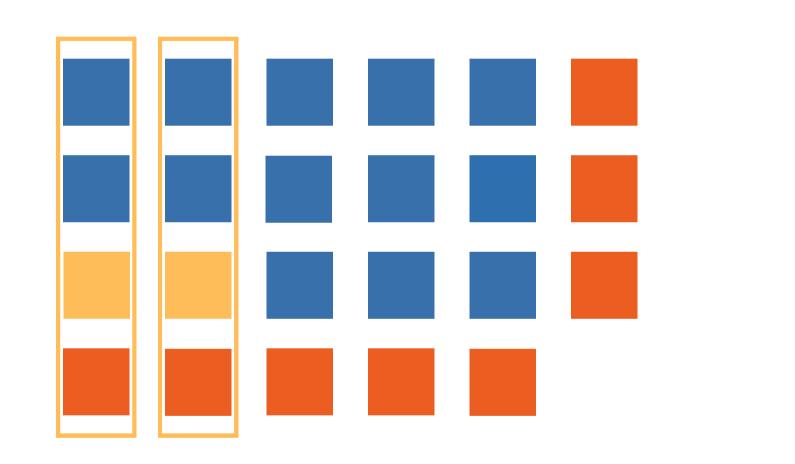


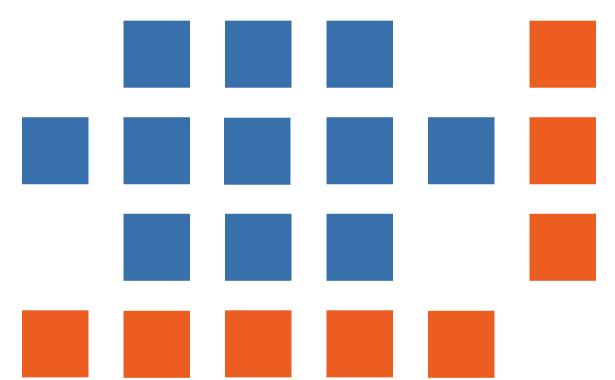
$$L = 5$$

$$D = 3$$









repair rows

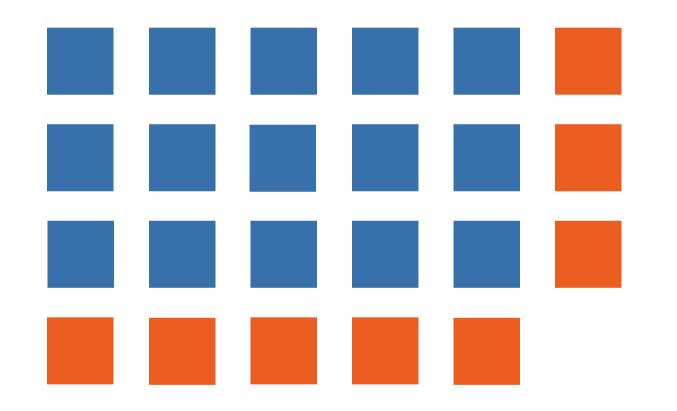
→ repair columns
repeat

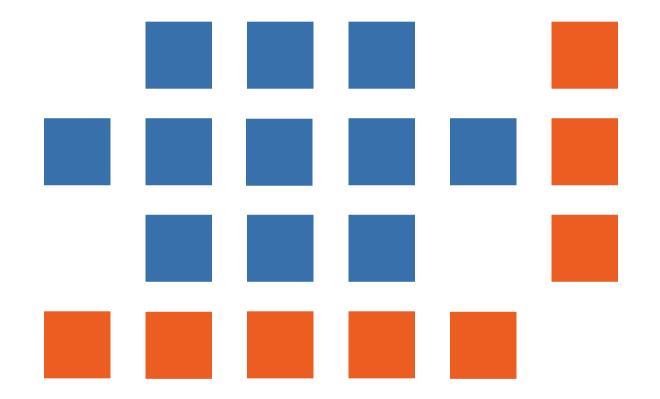
$$L = 5$$

$$D = 3$$





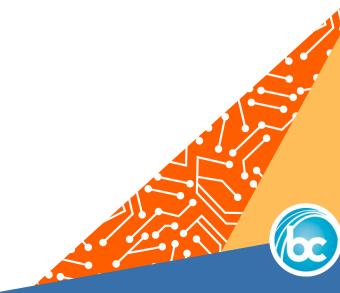




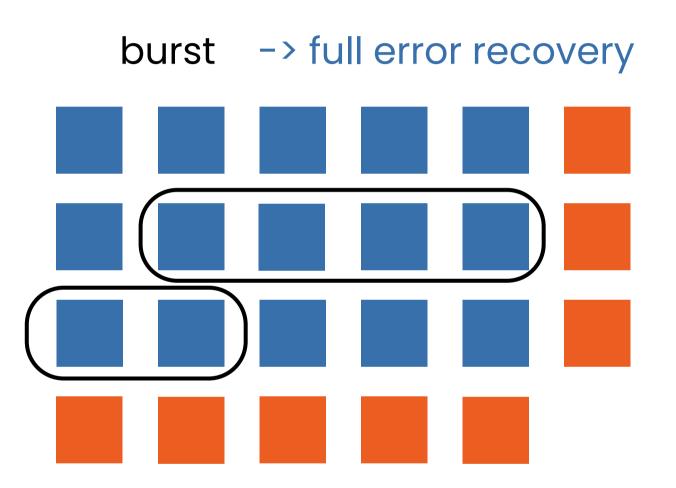
repair rows
repair columns
repeat
stop

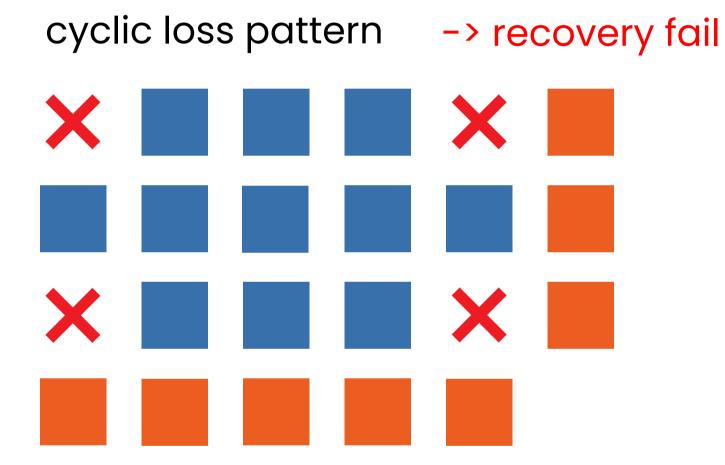
$$L = 5$$

$$D = 3$$









repair rows
repair columns
repeat
stop

- some loss pattern fail
- more robust to bursts

$$L = 5$$
$$D = 3$$

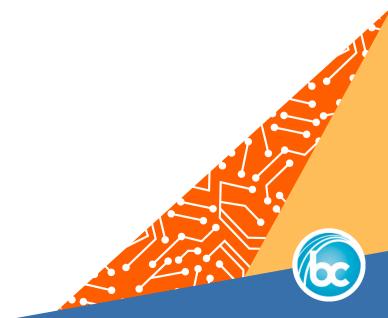






overhead = number of byte of the repair packets

number of byte of the protected source packets







• size(repair packet) > size (source packet)





OVERHEAD

- size(repair packet) > size (source packet)
- if all source packets have the same size :

overhead =
$$\frac{1}{I}$$

1D interleaved

overhead =
$$\frac{1}{D}$$

∘ 2D

overhead =
$$\frac{1}{L} + \frac{1}{D}$$





OVERHEAD

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$$\frac{1}{I}$$

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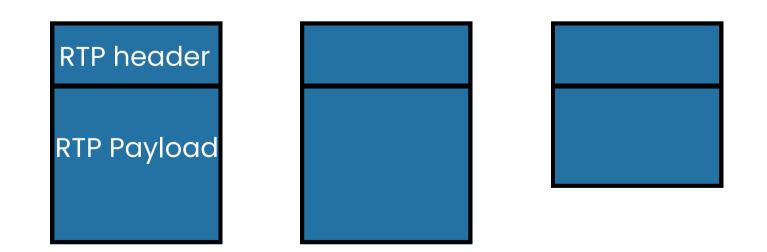
overhead =
$$\frac{1}{L} + \frac{1}{D}$$

L	D	overhead
10	0	0.10
5	0	0.20
5	5	0.40
3	3	0.67



Source packets

- unchanged
 - RTP header
 - RTP payload





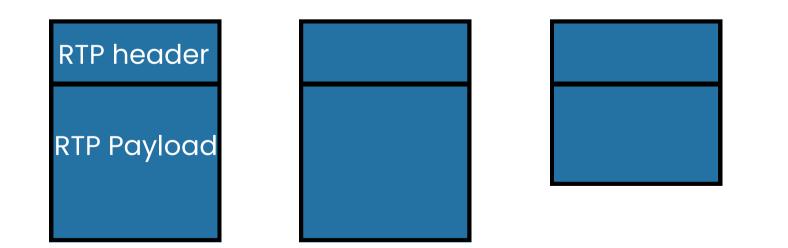


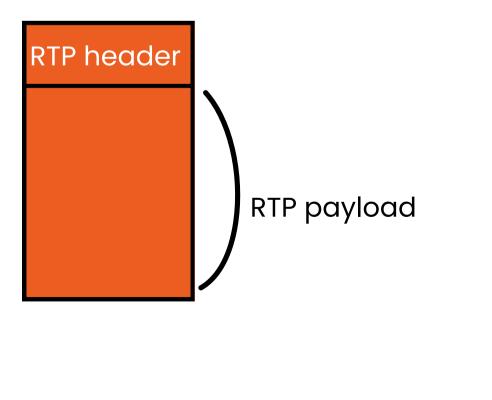
Source packets

- unchanged
 - RTP header
 - RTP payload

Repair packet

- RTP header
- RTP payload





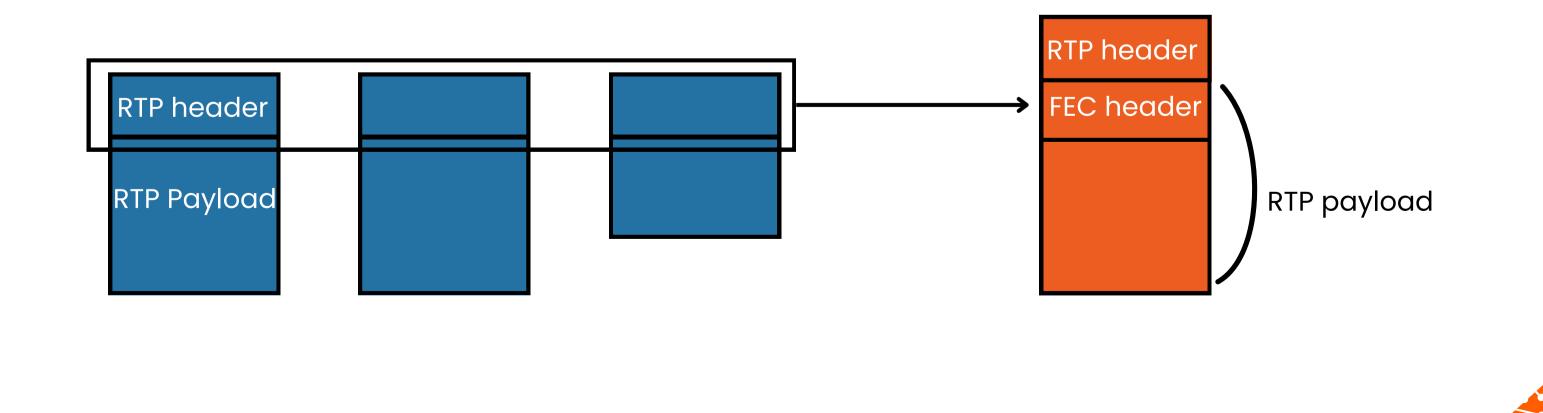


Source packets

- unchanged
 - RTP header
 - RTP payload
- unique sequence number in RTP header

Repair packet

- RTP header
- payload
 - FEC header: identification of the source packets protected



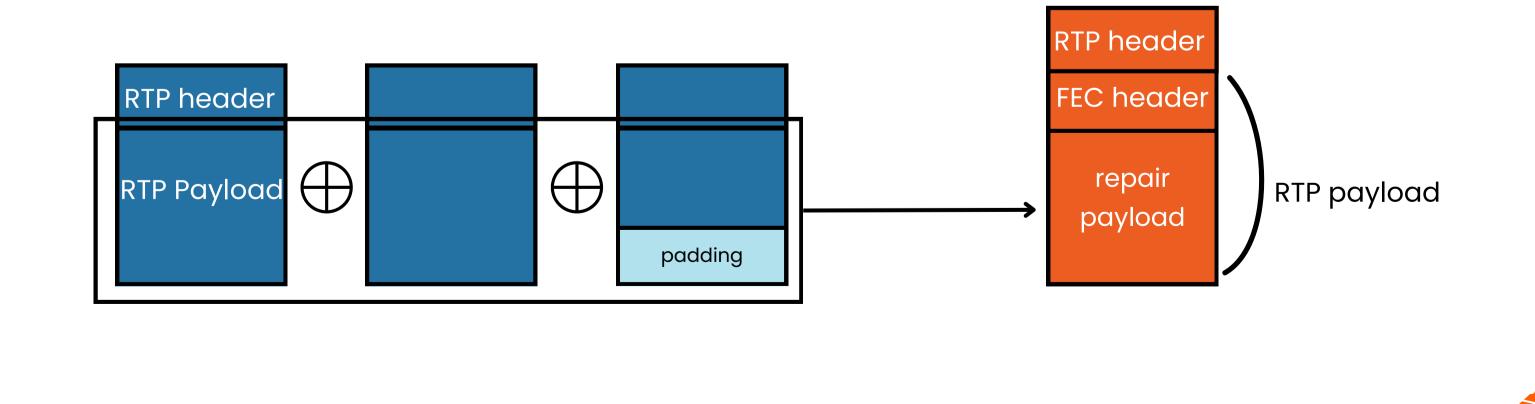


Source packets

- unchanged
 - RTP header
 - RTP payload
- unique sequence number in RTP header

Repair packet

- RTP header
- payload
 - FEC header: identification of the source packets protected
 - o repair payload: XOR result





- A single repair packet carries information about
 - the set of source packets protected
 - the protection configuration with L and D





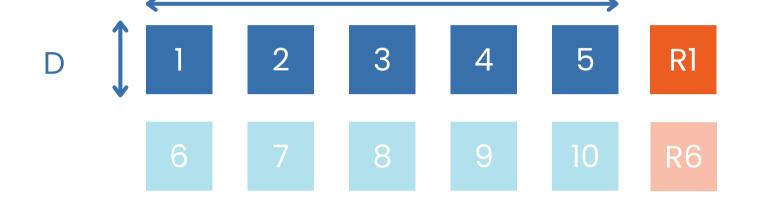


- A single repair packet carries information about
 - the set of source packets protected
 - the protection configuration with L and D
- For fixed L columns and D rows
 - L > 0, D = 0: row FEC, source packets: SN, SN+1, SN+(L-1)





- A single repair packet carries information about
 - the set of source packets protected
 - the protection configuration with L and D



For fixed L columns and D rows



L > 0, D = 1: row FEC followed by column, all columns following all rows repair packets



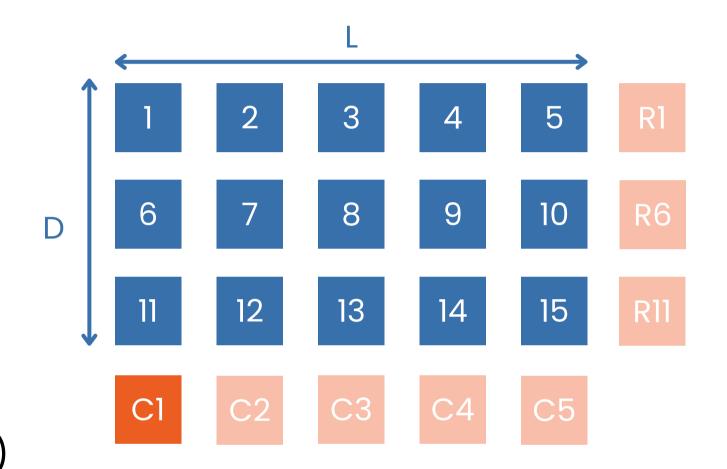
IR PACKET

- A single repair packet carries information about
 - the set of source packets protected
 - the protection configuration with L and D
- For fixed L columns and D rows
 - \circ L > 0, D = 0: row FEC, source packets: SN, SN+1, SN+(L-1)



L > 0, D = 1: row FEC followed by column, all columns following all rows repair packets

□ L > 0, D > 1: column FEC, source packets: SN, SN+L, ..., SN+(D-1)*L



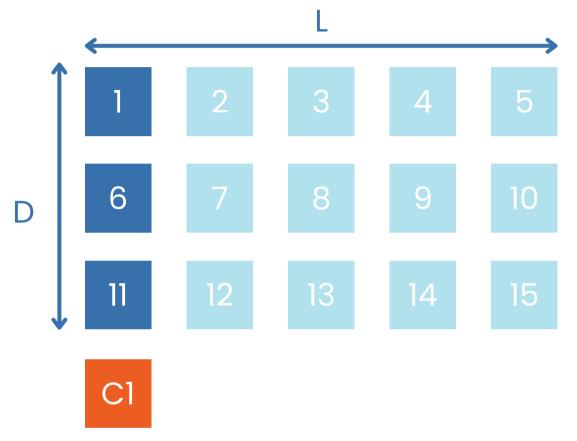


- A single repair packet carries information about
 - the set of source packets protected
 - the protection configuration with L and D
- For fixed L columns and D rows





□ L > 0, D > 1: column FEC, source packets: SN, SN+L, ..., SN+(D-1)*L





IV. IMPLEMENTATION CHALLENGES





IN LINPHONE

- Parameters
 - (L, D): (10, 0), (5, 0), (5, 5), (3, 3)
 - ideally 2D parity protection
 - adaptation to the loss rate and the network capabilities
 - o repair window: 200 ms
 - contain full source block, given any value of L, D
 - do not cause any delay in the video
- Implementation in Linphone SDK in C and C++
 - oRTP: library implementing the RTP protocol
 - Mediastreamer2: streaming engine for voice/video





OUR STRATEGY FOR VIDEO QUALITY

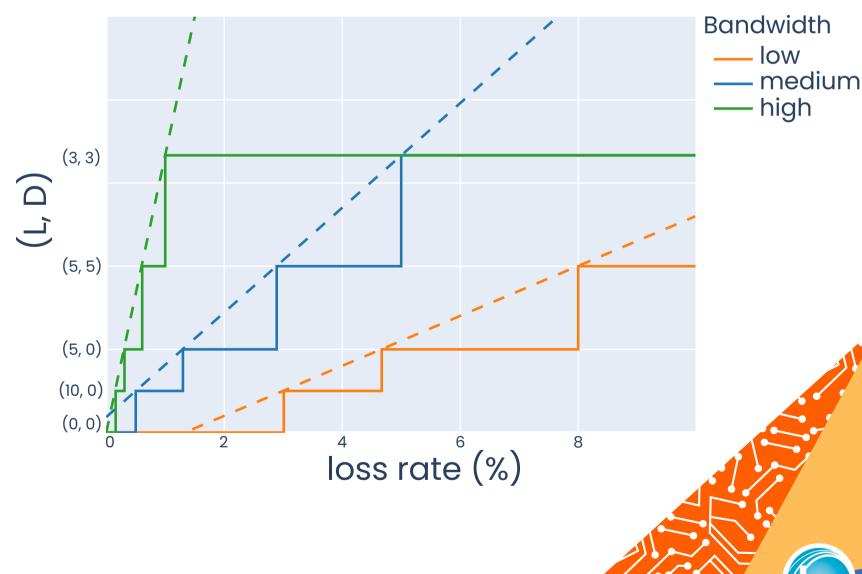
Best possible use of the bandwidth

- Optimal video setting with flexFEC
 - o definition, bitrate, frame rate
 - o no freeze

packet protection > high encoding setting

- Adaptatiblity: periodic control
 - o available bandwidth
 - loss rate
 - bandwidth dedicated to FEC
- Congestions: disabling FEC without delay

FEC protection adaptation to available bandwidth





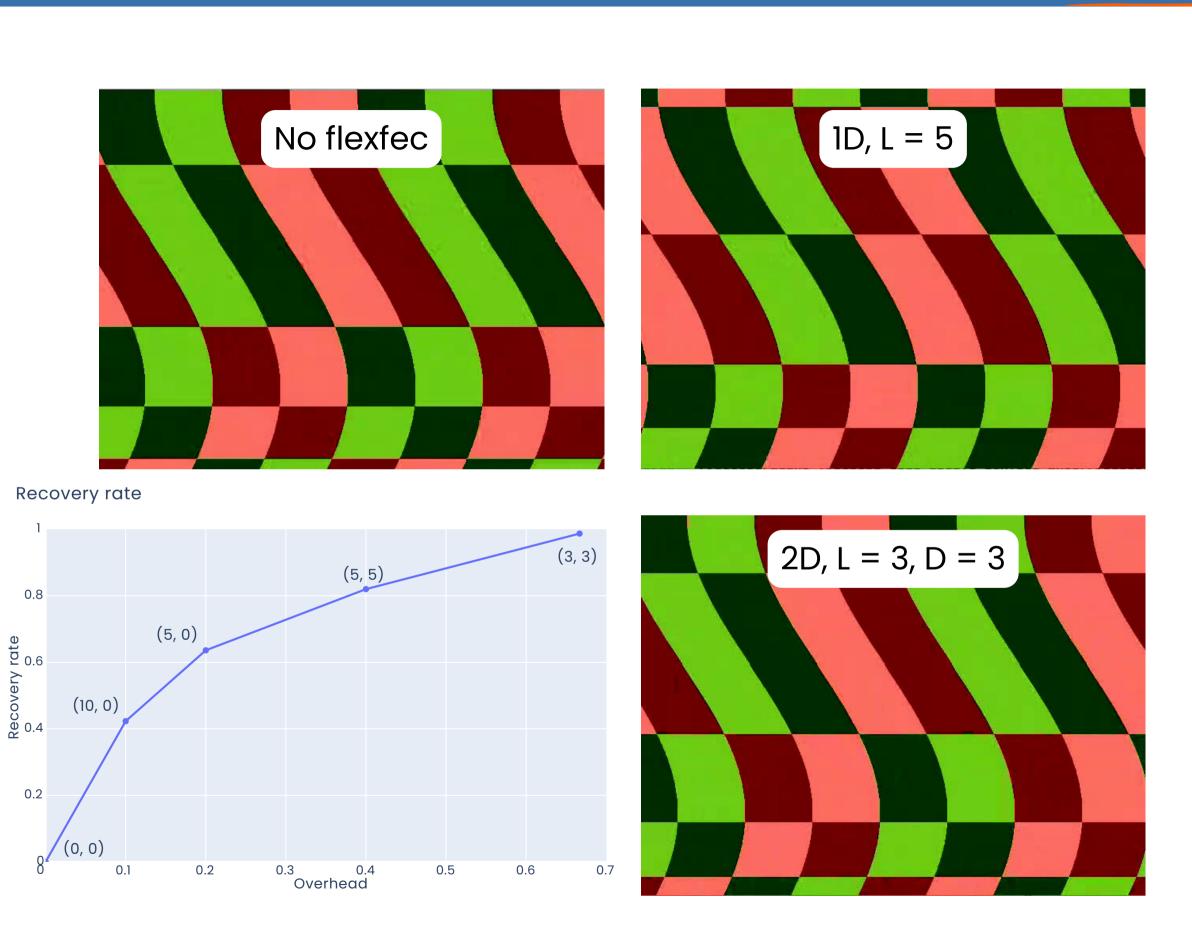
V. RESULTS





6% LOSS

AV1, 640x480 400kbits/s 30fps



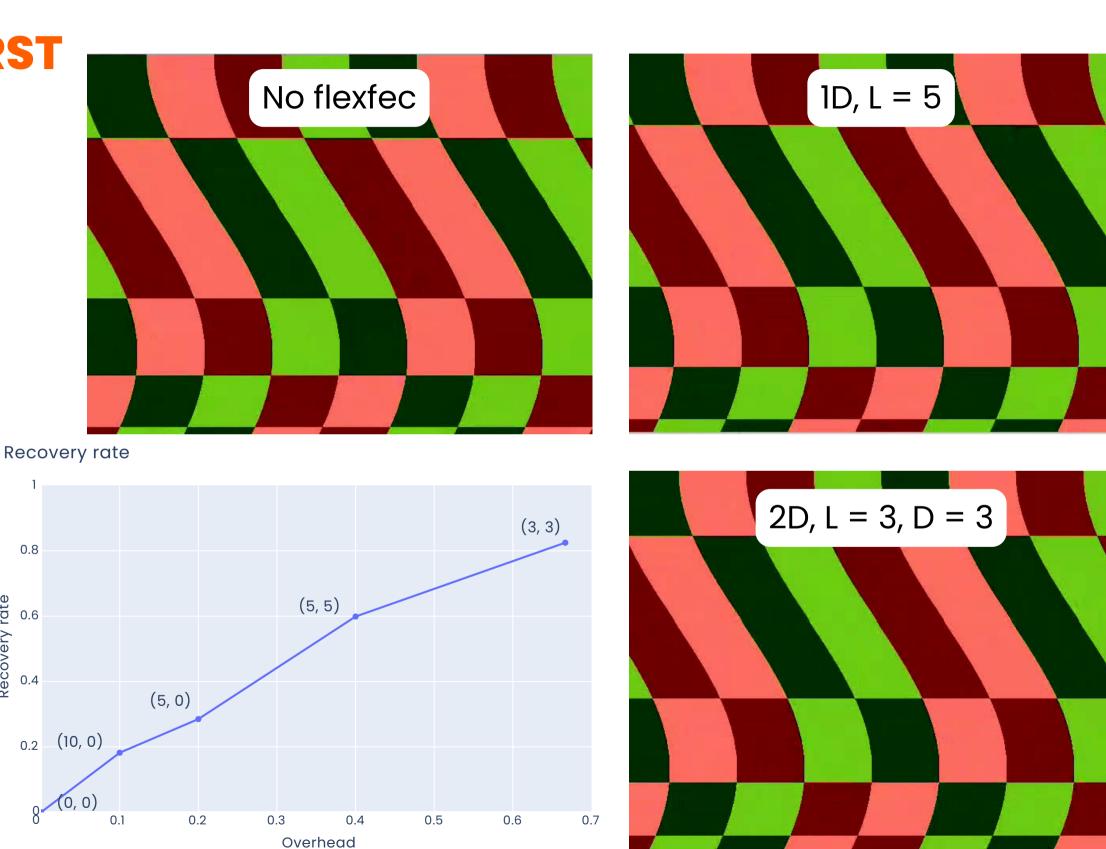


6% LOSS, BURST

0.8

Recovery rate 9.0 9.0

AVI, 640x480 400kbits/s 30fps





CONCLUSIONS

- Flexible FEC
 - Simple and efficient way to improve resiliency against packet loss
 - send redondant information
 - adaptable
 - short delay
 - Significant bandwidth needed ()
 - Complete description in RFC 8627, backward compatible
 - Real improvement in video quality
- FlexFEC coming soon in Linphone!
 - first for simple video call in 2024
 - o then for video conference, audio...





ANY QUESTIONS?





open source VOIP project



REFERENCES

- Linphone: https://www.linphone.org
- RFC 8627, RTP Payload Format for Flexible Forward Error Correction (FEC), 2019, https://datatracker.ietf.org/doc/html/rfc8627
- RFC 3550, RTP: A Transport Protocol for Real-Time Applications, 2003, https://datatracker.ietf.org/doc/html/rfc3550
- RFC 8854, WebRTC Forward Error Correction Requirements: https://datatracker.ietf.org/doc/html/rfc8854
- Source code linphone SDK: https://gitlab.linphone.org/BC/public/linphone-sdk



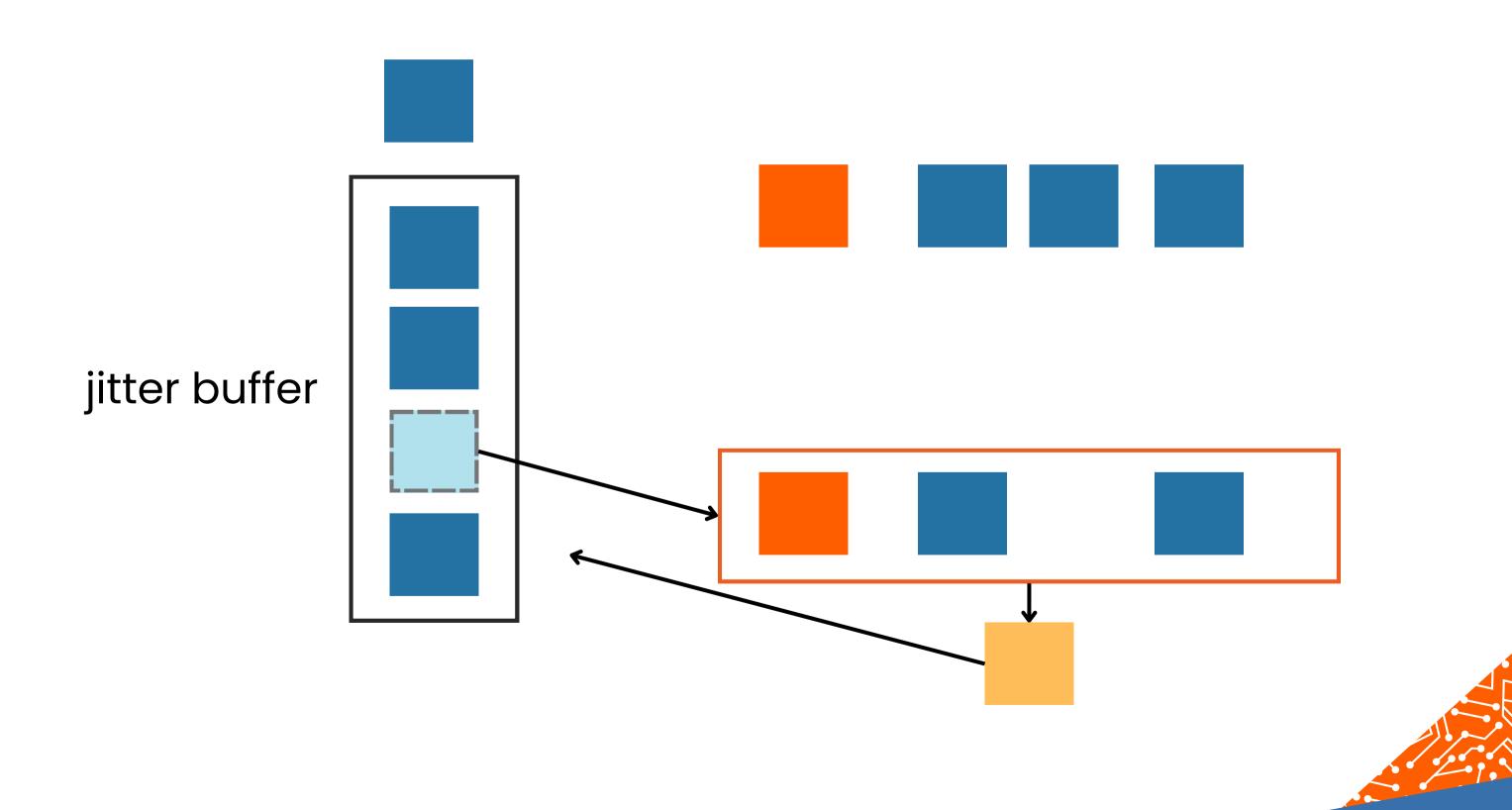


APPENDIX





RECEIVING FEC STREAM IN LINPHONE





FEC HEADER

• fixed L columns and D rows

0	1	Р	X	CC	М	Payload type recovery	Length recovery		
Timestamp recovery									
Sequence number base				Sequence r	num	ber base	L	D	
next SN base and L/D for CSRC in CSRC list									
Repair payload									

L and D can vary for different FEC blocks



ITERATIVE DECODING ALGORITHM (RFC 8627)

- 1. Set num_recovered_until_this_iteration to zero
- 2. Set num_recovered_so_far to zero
- 3. Recover as many source packets as possible by using the non-interleaved FEC repair packets and increase the value of num_recovered_so_far by the number of recovered source packets.
- 4. Recover as many source packets as possible by using the interleaved FEC repair packets and increase the value of num_recovered_so_far by the number of recovered source packets.
- 5. If num_recovered_so_far > num_recovered_until_this_iteration num_recovered_until_this_iteration = num_recovered_so_far Go to step 3.

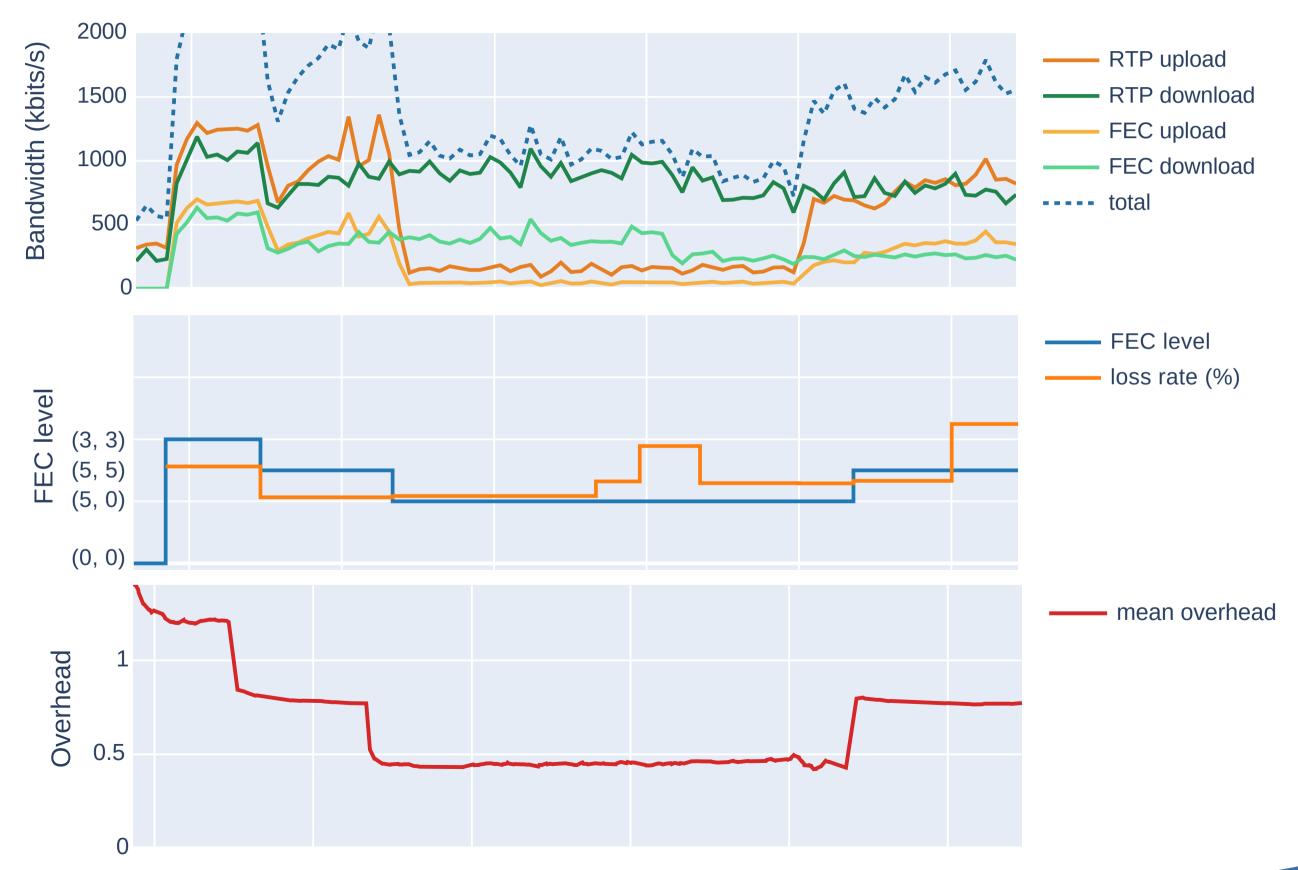
Else

Terminate











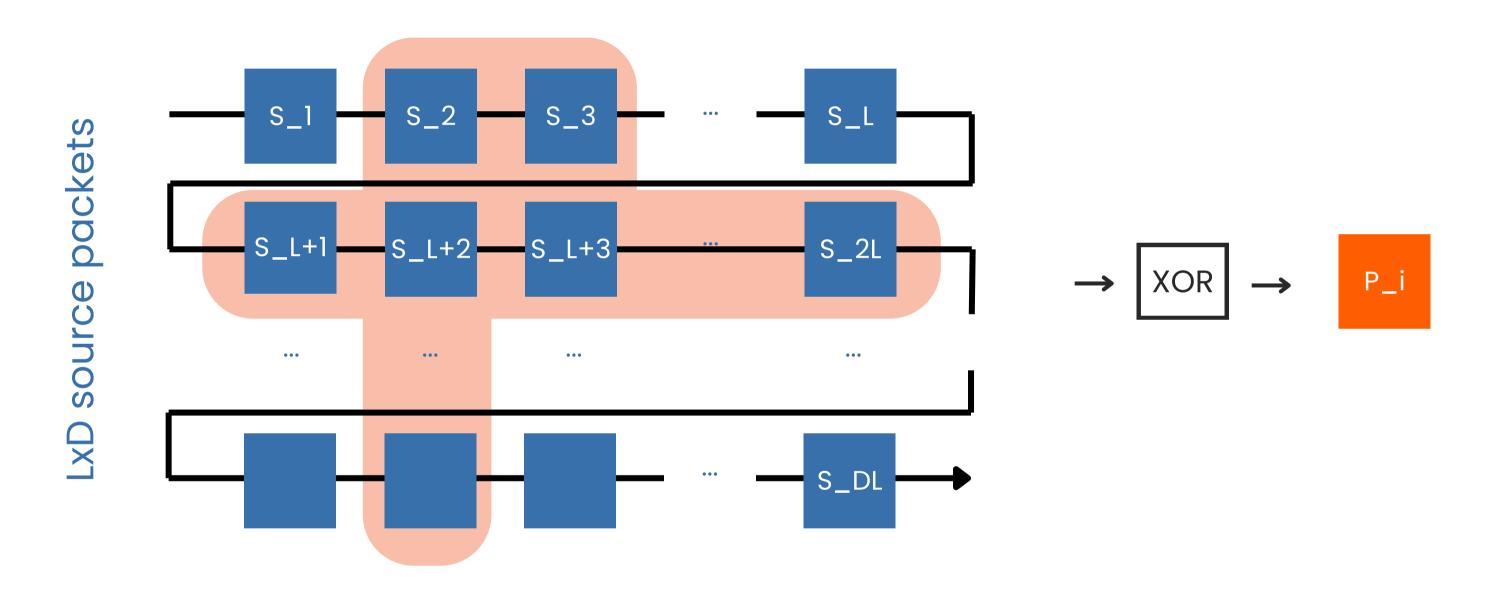


- Repair window: impacts the buffering
 - sender proposes a value in microseconds
 - rejected by receiver if it can't support the value
- Any unknown option can be ignored in the answer and FEC deleted

```
V=0
o=marie%20laroueverte_6jr06r~ 153 2733 IN IP6
s=Talk
c=IN IP6 2a01:e0a:27e:8210:fc05:3486:429e:714a
t=0 0
a=rtcp-xr:rcvr-rtt=all:10000 stat-summary=loss
a=group:BUNDLE as vs
m=audio 35787 RTP/AVP 0 101
a=rtpmap:101 telephone-event/8000
a=mid:as
a=extmap:1 urn:ietf:params:rtp-hdrext:sdes:mid
a=rtcp:58774
a=rtcp-fb:* trr-int 5000
a=rtcp-fb:* ccm tmmbr
m=video 51437 RTP/AVP 96 97
a=rtpmap:96 AV1/90000
a=rtpmap:97 flexfec/90000
a=fmtp:97 repair-window=200000
```



FEC PROTECTION WITH FLEXIBLE MASK



- must send a bitmap with the mask
- dynamic and flexible protection





REPAIR PACKET CONSTRUCTION

• Fixed L columns and D rows

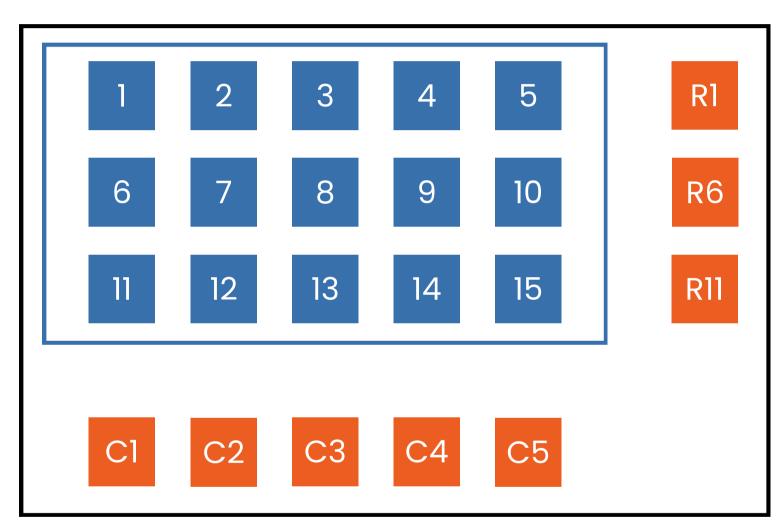
- \circ L > 0, D = 0: row FEC, source packets: SN, SN+1, SN+(L-1)
- L > 0, D = 1: row FEC followed by column, all columns following all rows repair packets
- L > 0, D > 1: column FEC, source packets: SN, SN+L, ..., SN+(D-1)*L
- L = 1, D = 0: retransmission of a single source packet







source block



FEC block

