

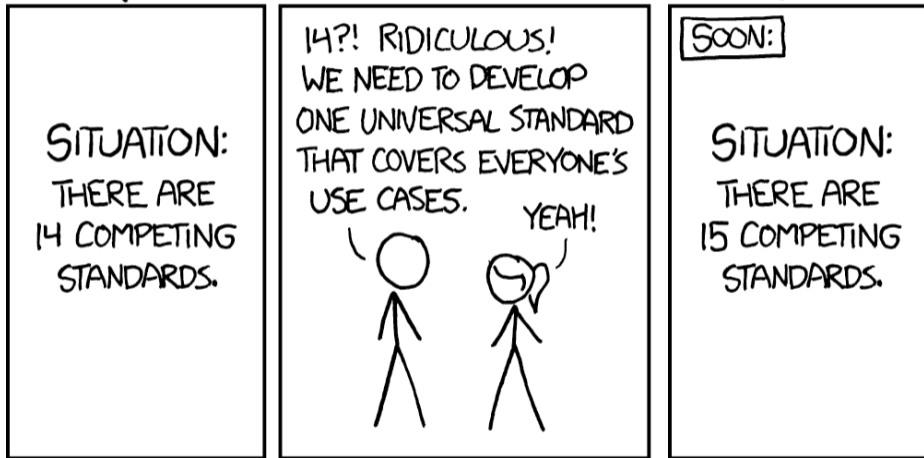
Yet another new SDR runtime?

Dr. Daniel Estévez

1 February 2025
FOSDEM, Brussels

HOW STANDARDS PROLIFERATE:

(SEE: A/C CHARGERS, CHARACTER ENCODINGS, INSTANT MESSAGING, ETC.)

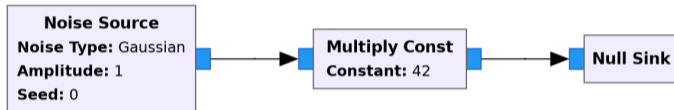


- Which is the *one true best* SDR runtime?

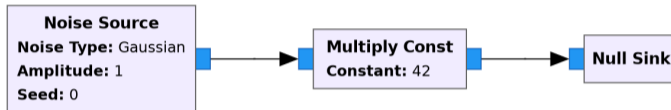
Asking the big questions

- Which is the *one true best* SDR runtime?
- What makes an SDR runtime be fast?
- Look at GNU Radio 3.10, GNU Radio 4.0, and FutureSDR

All the SDR runtimes look the same!

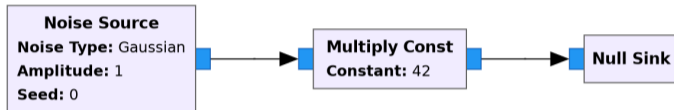


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- Connections in the flowgraph are a circular buffer shared between the connected blocks (single-producer multi-consumer)

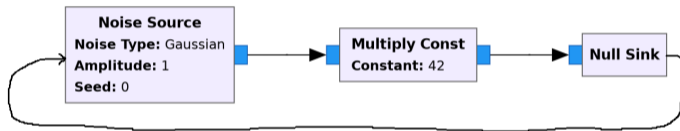
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- Connections in the flowgraph are a circular buffer shared between the connected blocks (single-producer multi-consumer)
- Different input and output buffers. Bad for cache.
- High (and difficult to control) latency on TX flowgraphs

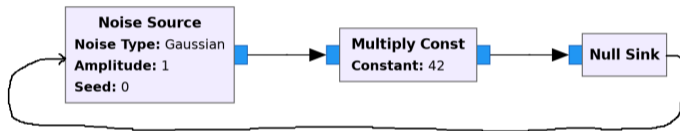
An alternative: closed circuits of “packets”

- Samples are sent in “packets” rather than as a continuous stream
- The flowgraph is divided into closed circuits, in which a fixed number of “packets” always exist
- “Packets” are recycled by sending them from a sink back to a source



An alternative: closed circuits of “packets”

- Samples are sent in “packets” rather than as a continuous stream
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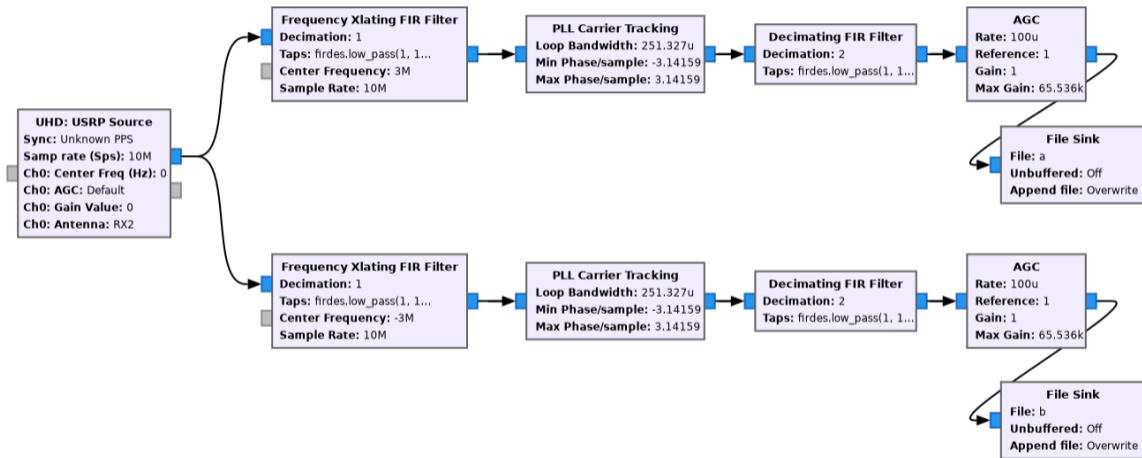


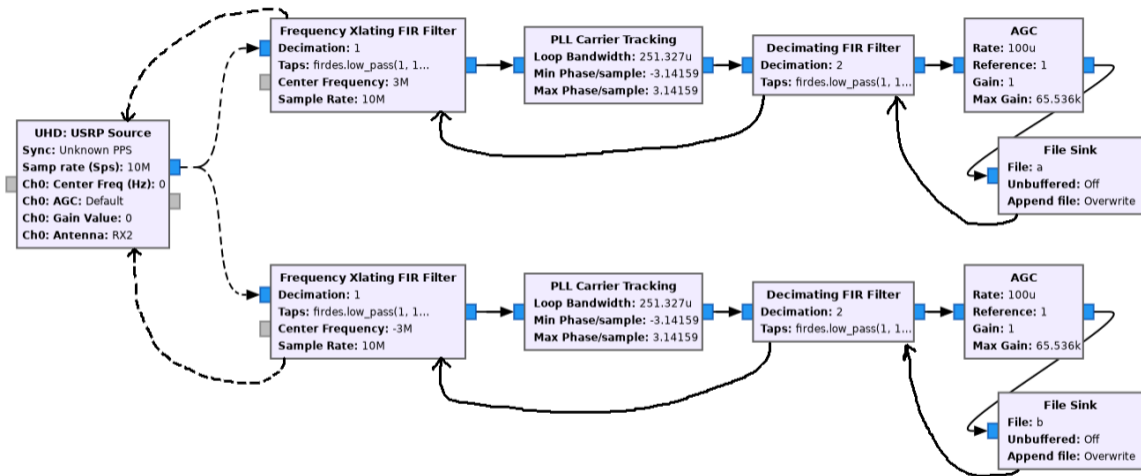
- Many blocks can work in-place on a “packet”
- Latency is determined by the number of “packets” in a circuit
- “Packets” can mark natural sections in the data (aligned to RF frame sections, etc.). Potentially less need for tags.
- More similar to a hand-crafted implementation where functions are called in sequence on the same buffer

The idea of a quantum

- Let's call this “packet” a **quantum**, because it is cool and because *packet* and *frame* are heavily overloaded
- A quantum contains:
 - A buffer with flexible *margins*. The margins allow inserting/adding a prefix/suffix in-place. Example use cases: CRC, synchronization word, cyclic prefix.
 - Tags, referred to sample indices within the packet
 - Perhaps other metadata that the user might need?

A more complex example flowgraph





- <https://github.com/daniestevez/qsdr>
- An implementation of these ideas in Rust using `async`
- Some benchmarks comparing with GNU Radio 3.10, GNU Radio 4.0 and FutureSDR
- Still an experimental work-in-progress

- GNU Radio 4.0 and FutureSDR support custom schedulers, but they are quite specific and not easy to write
- qsdr schedulers are based on Rust streams, so any code that can run streams can be a qsdr scheduler

- A stream is an object that can produce a sequence of values asynchronously

```
let mut stream = stream::iter(1..=3);  
  
assert_eq!(stream.next().await, Some(1));  
assert_eq!(stream.next().await, Some(2));  
assert_eq!(stream.next().await, Some(3));  
assert_eq!(stream.next().await, None);
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- A qcdr block has a work function which is `async` (it awaits to get its inputs), processes one quantum per input/output, and returns `Run`, `Done` or an error
- A qcdr block is converted to a stream. Each `next()` call on this stream calls the work function once, and returns either an error or nothing.

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- A qcdr block has a work function which is `async` (it awaits to get its inputs), processes one quantum per input/output, and returns `Run`, `Done` or an error
- A qcdr block is converted to a stream. Each `next()` call on this stream calls the work function once, and returns either an error or nothing.
- A qcdr scheduler is just some code that performs calls to multiple streams on one or several threads until all of them are done or there is an error

- Stream combinators are useful to build schedulers. For instance, `sequence2()` takes 2 streams (which produce items of type `Result<(), E>`) and produces a stream (with the same item type) that calls each of the 2 streams in sequence
- The helper function `run()` takes a stream and produces a future (an async result) that calls the stream until it is done or there is an error. This is usually the top-level element of the scheduler on each thread

```

type B = CacheAlignedBuffer<f32>;
let buffer_size = 4096;
let num_buffers = 4;

let head_elements = 100_000_000;

let buffers = std::iter::repeat_with(|| Quantum::new(B::new(buffer_size))).take(num_buffers);

let mut fg = Flowgraph::new();

let source = fg.add_block(NullSource::<Quantum<B>>::new());
let head = fg.add_block(Head::<Quantum<B>, Spsc, SpscRef>::new(head_elements));
let sink = fg.add_block(NullSink::<Quantum<B>>::new());

let mut circ = fg.new_circuit(buffers);
fg.connect(&mut circ, source.output(), head.input())?;
fg.connect_with_return(&mut circ, head.output(), sink.input(), source.input())?;

let mut fg = fg.validate()?;

let source = fg.extract_block(source)?;
let head = fg.extract_block(head)?;
let sink = fg.extract_block(sink)?;

println!("running flowgraph...");
block_on(run(sequence3(
    source.into_stream(),
    head.into_stream(),
    sink.into_stream(),
)))?;
println!("flowgraph finished");

```

```

#[derive(Block, Debug)]
#[qsdr_crate = "crate"]
#[work(WorkInPlace)]
pub struct Head<T, Cin = Spsc, Cout = Spsc>
where
    Cin: Channel,
    Cin::Receiver<T>: Receiver<T>,
    Cout: Channel,
{
    #[port]
    input: PortIn<T, Cin>,
    #[port]
    output: PortOut<T, Cout>,
    remaining: u64,
}

impl<T, Cin, Cout> Head<T, Cin, Cout>
where
    Cin: Channel,
    Cin::Receiver<T>: Receiver<T>,
    Cout: Channel,
{
    pub fn new(count: u64) -> Self {
        Self {
            input: Default::default(),
            output: Default::default(),
            remaining: count,
        }
    }
}

```

```

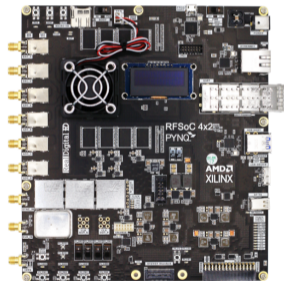
impl<T, Cin, Cout> WorkInPlace<T> for Head<T, Cin, Cout>
where
    Cin: Channel,
    Cin::Receiver<T>: Receiver<T>,
    Cout: Channel,
{
    async fn work_in_place(&mut self, _: &mut T) -> Result<WorkStatus> {
        assert!(self.remaining > 0);
        self.remaining -= 1;
        if self.remaining == 0 {
            Ok(DoneWithOutput)
        } else {
            Ok(Run)
        }
    }
}

```

- Choose a family of simple flowgraphs
- Write by hand an implementation that performs as best as possible
- Write implementations in qsdr, GNU Radio 3.10, GNU Radio 4.0 and FutureSDR
- Measure the rate (samples/second) at which each implementation can run

Benchmarking platform

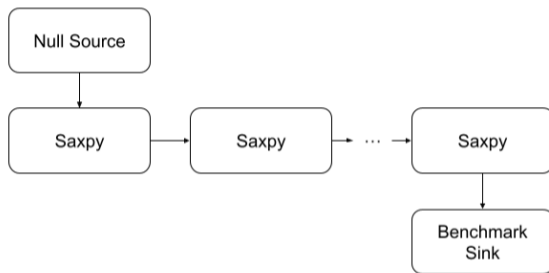
- Quad-core Cortex-A53 in AMD (Xilinx) MPSoC (1.33 GHz clock)
- Found in many AMD FPGA SoC platforms, in which a high-performance SDR runtime could be an alternative to an FPGA implementation for many signal processing problems



- Using a Kria KV260 board for development (\$249 MSRP)

Benchmarking flowgraph

- Saxpy kernel: $y[n] = ax[n] + b$ (note: usual saxpy has a vector in place of b)



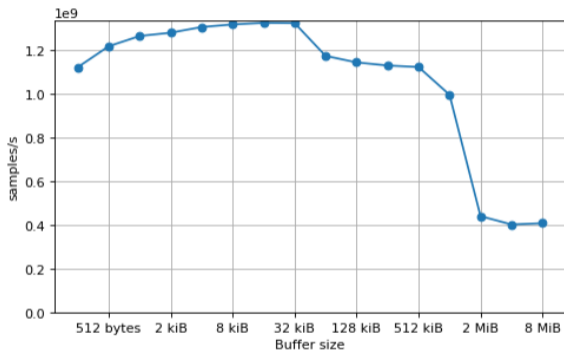
- Null Source does not do anything. Not even `memset()` the output to zero.
- Benchmark Sink counts samples and measures the sample rate

Saxpy kernel implementation

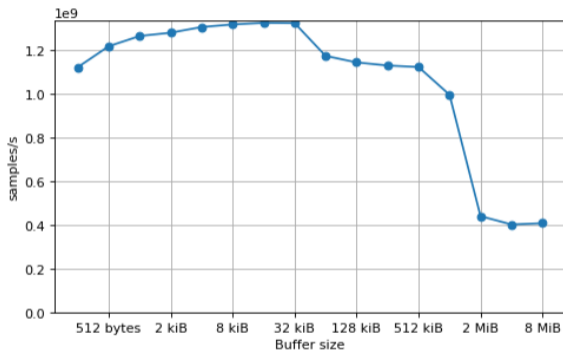
- Hand-written assembly using NEON SIMD
- Throughput of 1 float / clock cycle (2 FLOPs / clock cycle), which is the maximum according to the Cortex-A53 hardware limitations:
 - 2 NEON units each capable of operating on a 64-bit vector
 - 64-bit load path, 128-bit store path
 - Load/store to SIMD register uses corresponding NEON unit
 - One AGU for load/store
- For comparison, an optimal `memcpy()` is 1.33 floats / clock cycle (5.33 bytes / clock cycle)
- 2x as fast as the typical NEON code generated by gcc, clang and rustc, which is a 0.5 floats / clock cycle naïve implementation
- Heavily uses tricks related to the partial dual-issue capability of the Cortex-A53

14068:	4c402944	ldr	{v4.4s-v7.4s}, [x10]	1410c:	4e181da6	mov	v6.d[1], x13
1406c:	4f919084	fmul	v4.4s, v4.4s, v17.s[0]	14110:	4e081de7	mov	v7.d[0], x15
14070:	f9804140	prfm	pldl1keep, [x10, #128]	14114:	4e181d64	mov	v4.d[1], x11
14074:	4f9190a5	fmul	v5.4s, v5.4s, v17.s[0]	14118:	4e30d442	fadd	v2.4s, v2.4s, v16.4s
14078:	f940254b	ldr	x11, [x10, #72]	1411c:	f9806100	prfm	pldl1keep, [x8, #192]
1407c:	4f9190c6	fmul	v6.4s, v6.4s, v17.s[0]	14120:	4e30d463	fadd	v3.4s, v3.4s, v16.4s
14080:	f9402d4c	ldr	x12, [x10, #88]	14124:	f9808100	prfm	pldl1keep, [x8, #256]
14084:	4f9190e7	fmul	v7.4s, v7.4s, v17.s[0]	14128:	4f919084	fmul	v4.4s, v4.4s, v17.s[0]
14088:	f940354d	ldr	x13, [x10, #104]	1412c:	f940450b	ldr	x11, [x8, #136]
1408c:	4e30d484	fadd	v4.4s, v4.4s, v16.4s	14130:	4f9190a5	fmul	v5.4s, v5.4s, v17.s[0]
14090:	f9403d4e	ldr	x14, [x10, #120]	14134:	4f9190c6	fmul	v6.4s, v6.4s, v17.s[0]
14094:	4e30d4a5	fadd	v5.4s, v5.4s, v16.4s	14138:	f9404d0c	ldr	x12, [x8, #152]
14098:	f940394f	ldr	x15, [x10, #112]	1413c:	4f9190e7	fmul	v7.4s, v7.4s, v17.s[0]
1409c:	4e30d4c6	fadd	v6.4s, v6.4s, v16.4s	14140:	f940550d	ldr	x13, [x8, #168]
140a0:	f9804100	prfm	pldl1keep, [x8, #128]	14144:	4e30d484	fadd	v4.4s, v4.4s, v16.4s
140a4:	fd402140	ldr	d0, [x10, #64]	14148:	f9405d0e	ldr	x14, [x8, #184]
140a8:	4e181dc3	mov	v3.d[1], x14	1414c:	4e30d4a5	fadd	v5.4s, v5.4s, v16.4s
140ac:	fd402941	ldr	d1, [x10, #80]	14150:	f940590f	ldr	x15, [x8, #176]
140b0:	4e181d60	mov	v0.d[1], x11	14154:	4c002900	stl	{v0.4s-v3.4s}, [x8]
140b4:	fd403142	ldr	d2, [x10, #96]	14158:	fd404901	ldr	d1, [x8, #144]
140b8:	4e181d81	mov	v1.d[1], x12	1415c:	4e181dc3	mov	v3.d[1], x14
140bc:	4e081de3	mov	v3.d[0], x15	14160:	fd405102	ldr	d2, [x8, #160]
140c0:	4e181da2	mov	v2.d[1], x13	14164:	4e181d81	mov	v1.d[1], x12
140c4:	4e30d4e7	fadd	v7.4s, v7.4s, v16.4s	14168:	fc480d00	ldr	d0, [x8, #128]!
140c8:	4f919000	fmul	v0.4s, v0.4s, v17.s[0]	1416c:	4e181da2	mov	v2.d[1], x13
140cc:	f940454b	ldr	x11, [x10, #136]	14170:	4e081de3	mov	v3.d[0], x15
140d0:	4f919021	fmul	v1.4s, v1.4s, v17.s[0]	14174:	4e181d60	mov	v0.d[1], x11
140d4:	4f919042	fmul	v2.4s, v2.4s, v17.s[0]	14178:	4e30d4c6	fadd	v6.4s, v6.4s, v16.4s
140d8:	f9404d4c	ldr	x12, [x10, #152]	1417c:	eb09015f	cmp	x10, x9
140dc:	4f919063	fmul	v3.4s, v3.4s, v17.s[0]	14180:	4e30d4e7	fadd	v7.4s, v7.4s, v16.4s
140e0:	f940554d	ldr	x13, [x10, #168]	14184:	54fffa21	b.ne	140c8
140e4:	4e30d400	fadd	v0.4s, v0.4s, v16.4s	14188:	4f919000	fmul	v0.4s, v0.4s, v17.s[0]
140e8:	f9405d4e	ldr	x14, [x10, #184]	1418c:	4f919021	fmul	v1.4s, v1.4s, v17.s[0]
140ec:	4e30d421	fadd	v1.4s, v1.4s, v16.4s	14190:	4f919042	fmul	v2.4s, v2.4s, v17.s[0]
140f0:	f940594f	ldr	x15, [x10, #176]	14194:	4f919063	fmul	v3.4s, v3.4s, v17.s[0]
140f4:	4c002944	stl	{v4.4s-v7.4s}, [x10]	14198:	4c002944	stl	{v4.4s-v7.4s}, [x10]
140f8:	fd404945	ldr	d5, [x10, #144]	1419c:	4e30d400	fadd	v0.4s, v0.4s, v16.4s
140fc:	4e181dc7	mov	v7.d[1], x14	141a0:	4e30d421	fadd	v1.4s, v1.4s, v16.4s
14100:	fd405146	ldr	d6, [x10, #160]	141a4:	4e30d442	fadd	v2.4s, v2.4s, v16.4s
14104:	4e181d85	mov	v5.d[1], x12	141a8:	4e30d463	fadd	v3.4s, v3.4s, v16.4s
14108:	fc480d44	ldr	d4, [x10, #128]!	141ac:	4c002900	stl	{v0.4s-v3.4s}, [x8]

Saxpy kernel benchmark depending on buffer size



Saxpy kernel benchmark depending on buffer size



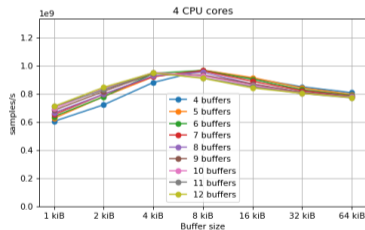
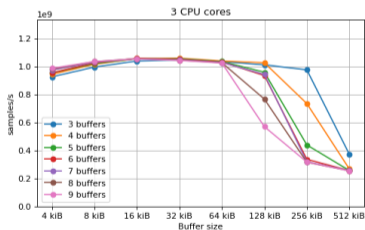
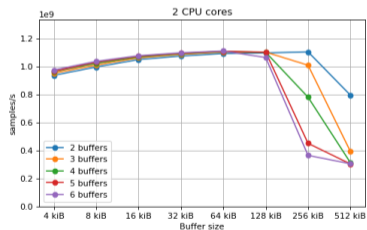
- 32 KiB L1 cache (per core), 1 MiB L2 cache (shared by all cores)

Some comments about SDR runtime performance

- We want most of the CPU time to be spent on running our `work()` functions, rather than something else
- SDR runtime overhead depends on number of `work()` calls per second
- Using large buffers for each `work()` call is not an option. Ideally we want to stay in L1 cache (and definitely not go to DDR).
- Fast simple kernels which can process an L1 cache worth of data quickly are the worst case, since SDR runtime overhead can be significant
- Some numbers for Saxpy kernel: 16 KiB buffer \rightarrow 4096 floats. ~ 1 float/cycle at 1.33 GHz $\rightarrow 3 \mu\text{s}$ per `work()` call.

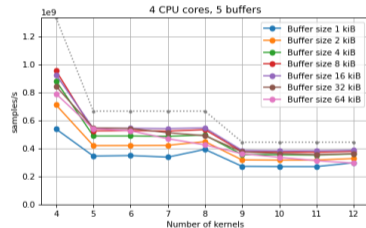
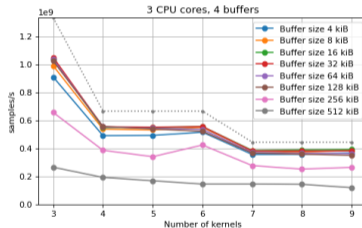
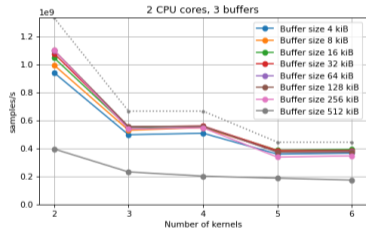
Saxpy kernel benchmark in multiple CPU cores

- Rust channels used to send buffers between threads. Custom high-performance channels implemented in qsdr.
- One thread pinned to each CPU core
- Fixed number of buffers passed around in a circuit formed by the threads



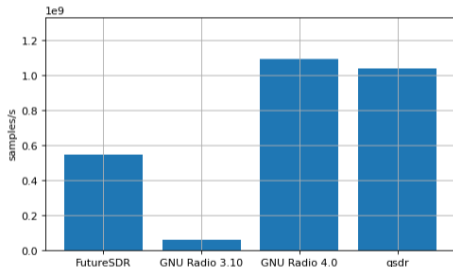
Saxpy kernel benchmark of multiple kernels

- M kernels and N CPU cores. Kernels statically pinned to CPU cores. Worst case, $\text{ceil}(M/N)$ kernels sharing a core.

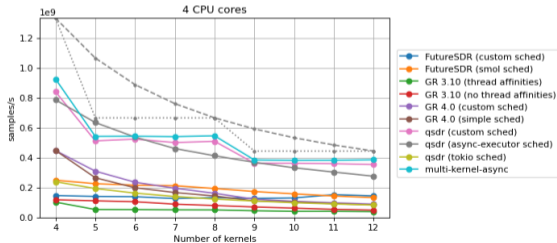
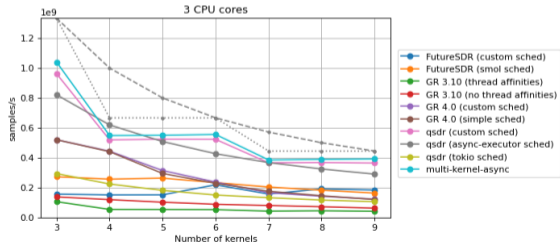
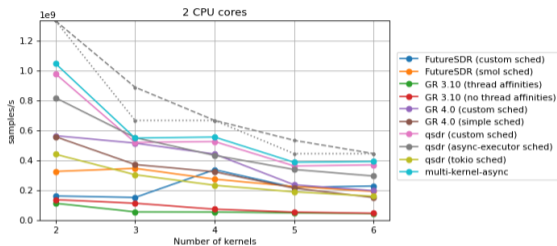
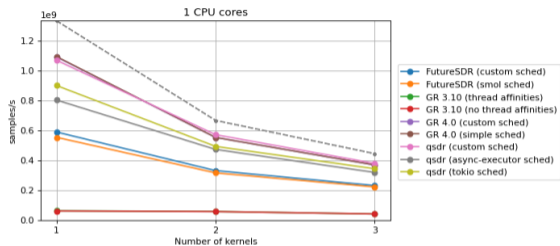


SDR runtime benchmark with a single Saxpy kernel and single core

- Flowgraph: null source → saxpy → benchmark sink
- Uses simplest scheduler to run all the blocks in the same CPU



SDR runtime benchmark of multiple Saxpy kernels



- There is a lot of room for improvement in SDR runtime performance
- What is the future of qsdr?
 - Hard to say at this point. Might develop further or remain as an experiment
 - Currently more intended as a source of ideas and to compare with other SDR runtimes
 - If it develops further, the goal would be a lower-level runtime than GNU Radio. A middle ground between having many things done/chosen for you versus having to write everything from scratch.