# **CONSTANT TIME BIG INTEGERS**

in Ada and SPARK

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https://github.com/AldanTanneo/bigints

- Avoid side channels in cryptography code
- Provide strong primitives for the Ada ecosystem



#### What I wrote

- Big integer library in Ada and SPARK, for the Ada ecosystem
- Design and algorithms heavily inspired by Rust's crypto-bigint [1]
- All functions in constant time unless explicitly named \*\_Vartime

Disclaimer: unaudited code written by a single person, **do not** use for security critical applications 😠

# "Constant time" algorithms:

- no branching depending on secret values
- no memory accesses depending on secret values

This includes hidden variable-time constructs like the div instruction on x86...

# Side channels from non-constant time implementations

- timing / power consumption attacks
- cache access analysis

*for ex:* binary exponentiation leaks secret RSA exponent (power analysis) or number of 1 bits in exponent (timing analysis)



#### **But:**

- Code less readable, harder to audit
- Code is (usually) slower
- Optimizing compilers might recognize your constant-time algorithm and replace it with a faster, variable-time one!

#### **Binary exponentiation**

Variable time

```
function Pow (A : Uint; N : Natural)
return Uint is
 Y : Uint := ONE;
 X : Uint := A;
begin
  for I in 0 .. BITS - 1 loop
    if Bit (N, I) then
    Y := Y * X;
    end if;
   X := X * X;
  end loop;
  return Y;
end Pow;
```

#### Constant time

```
function Pow (A : Uint; N : Natural)
return Uint is
 Y : Uint := ONE;
 X : Uint := A;
 B : Boolean;
begin
 for I in 0 .. BITS - 1 loop
    B := Bit (N, I);
   Y := Cond_Select (Y, Y * X, B);
   X := X * X;
 end loop;
  return Y;
end Pow;
```

The Cond\_Select primitive is implemented with xor-ing and masking:

$$f(a, b, \text{mask}) = a \oplus (\text{mask} \cdot (a \oplus b))$$

Some care must be taken when transforming a boolean condition into a mask, so as to not have a hidden branch from compiler optimization.

# Formal proof of contracts in SPARK

- Guarantee no runtime errors
- Proof of bitwise operations used in CT algorithms

function Cond\_Select (A, B : Uint; C : Choice) return Uint with
 Post => Cond\_Select'Result = (if To\_Bool (C) then B else A);
procedure CSwap (A, B : in out Uint; C : Choice) with
 Post => (if To\_Bool (C) then B = A'Old and then A = B'Old
 else A = A'Old and then B = B'Old);

4 We cannot prove that they effectively are in constant time

# In fact, all constant-time code that is not written in raw assembly relies on best effort implementations

- Obfuscate intent from the compiler
- Add optimization barriers
   *for ex:* System.Machine\_Code.Asm ("", Volatile => True)

### To convince ourselves that our code is constant time

- Audit assembly output
- Measure execution time w. statistics on cycles, cache, power draw...

#### **High level proofs**

- Higher level proofs in the library? using Ada's Big\_Integer package to model own representation
- Current GNATProve Alire package cannot prove these yet, so proof code is in a separate branch (addition-proof)

- Test suite includes small elliptic curve implementation using the library's modular integer facilities [2], [3]
- Define operations on a prime field, by instantiating a generic package:

```
package U256s is new Bigints.Uints (256);
subtype U256 is U256s.Uint;
use U256s;
P : constant U256 := Shl (ONE, 255) - From_U64 (19);
-- Prime P = 2^255 - 19
package GF_P is new Bigints.Modular (U256s, P);
-- Prime field GF(P), used in Curve25519 operations ⊜
```

• Available in the Alire index!

# REFERENCES

- [1] "RustCrypto: Cryptographic Big Integers." [Online]. Available: https://github.com/ RustCrypto/crypto-bigint
- [2] "Constant time big integers in SPARK." [Online]. Available: https://github.com/ AldanTanneo/bigints
- [3] "Elliptic Curves for Security." [Online]. Available: https://www.rfc-editor.org/rfc/ rfc7748