

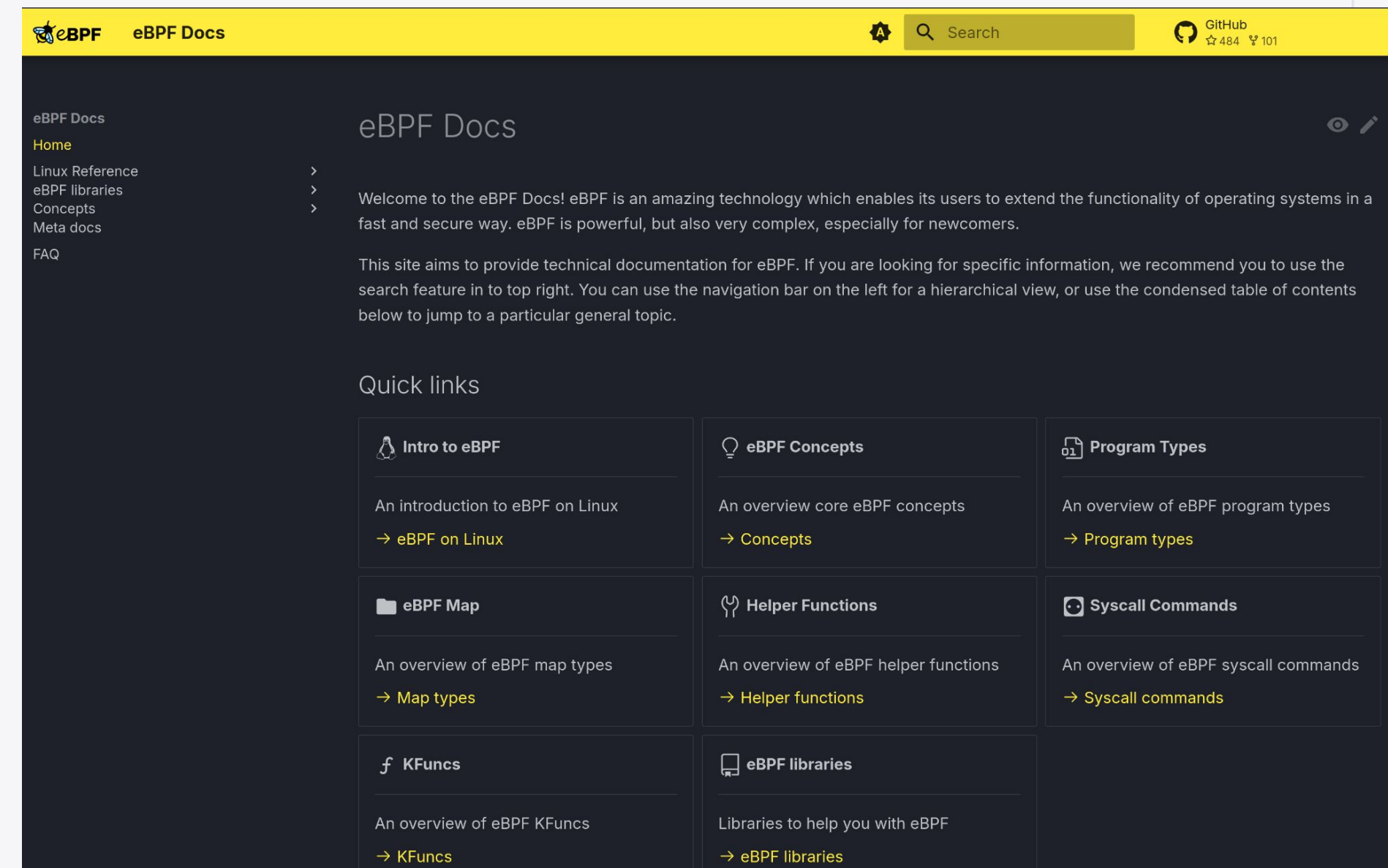
# eBPF Reachability Analysis



Speaker: **Dylan Reimerink**

# @me

- Dylan Reimerink
- Software Engineer - Isovalent @ Cisco
- Cilium committer
- Maintainer of [ebpf-go](https://github.com/cilium/ebpf) (cilium/ebpf)
- Maintainer of [docs.ebpf.io](https://docs.ebpf.io)



**ISOVALENT**  
now part of **cisco**

# Acknowledgement

- Timo Beckers
  - Cilium committer
  - ebpf-go maintainer
  - co-author of reachability analysis
- Robin Gögge
  - Cilium reviewer + org member
  - ebpf-go reviewer
  - Primary reachability analysis reviewer

# Modernizing Cilium

- Cilium is over 10 years old, one of the first eBPF users
- Cilium compiles its eBPF programs at runtime (for now)
- Our goal is to pre-compile all eBPF programs
  - Faster program loading
  - Smaller docker images (not shipping clang toolchain and dependencies)
  - Less attack surface
  - Less complexity, so hopefully less bugs
- Working towards this since 2023

# Load time configuration

- Cilium has a significant amount of optional features and settings
  - Cilium has 94 datapath settings (at the moment)
- Disabled features were conditionally compiled with pre-processor macros
- Compile time config needs to become load time config
  - v5.1: [bpf: dead code elimination](#)
  - v5.5: [Track read-only map contents as known scalars in BPF verifiers](#)

```
volatile const bool enable_feature_a;

int SEC("tc") entrypoint(struct __sk_buff *ctx) {
    if (enable_feature_a) {
        // ...
    }
    return TC_ACT_OK;
}
```



# The “problem”

- Cilium core value: Only pay for what you use
- So how do I prevent paying for a map I don't use?

```
#ifdef ENABLE_FEATURE_A
struct {
    __uint(type, BPF_MAP_TYPE_HASH);
    __type(key, struct some_key);
    __type(value, struct some_value);
    __uint(max_entries, 1000000);
} feature_a_map SEC(".maps");
#endif

int SEC("tc") entrypoint(struct __sk_buff *ctx) {
#ifdef ENABLE_FEATURE_A
    struct some_key key = { /* ... */ };
    struct some_value *val = bpf_map_lookup_elem(&feature_a_map, &key);
    // ...
#endif
    return TC_ACT_OK;
}
```

# The naive fix?

```
struct {
    uint(type, BPF MAP TYPE HASH);
    __type(key, struct some_key);
    type(value, struct some_value);
    uint(max_entries, 1000000);
} feature_a_map SEC(".maps");

volatile const bool enable_feature_a;

int SEC("tc") entrypoint(struct __sk_buff *ctx) {
    if (enable_feature_a) {
        struct some_key key = { /* ... */ };
        struct some_value *val = bpf_map_lookup_elem(&feature_a_map, &key);
        // ...
    }
    return TC_ACT_OK;
}
```

# The naive fix?

```
$ sudo bpftool prog dump xlated id 2390
int entrypoint(struct sk_buff * ctx):
; int SEC("tc") entrypoint(struct __sk_buff *ctx)
    0: (b7) r0 = 0
; if (enable feature a)
    1: (18) r1 = map[id:281][0]+0
    3: (71) r1 = *(u8 *) (r1 +0)
; if (enable feature a)
    4: (15) if r1 == 0x0 goto pc+14
    5: (b7) r1 = 0
; struct some_key key = { /*...*/ };
    6: (63) *(u32 *) (r10 -8) = r1
    7: (bf) r2 = r10
    8: (07) r2 += -8
    9: (18) r1 = map[id:279]
   11: (85) call __htab_map_lookup_elem#294448
   12: (15) if r0 == 0x0 goto pc+1
   13: (07) r0 += 56
   14: (bf) r1 = r0
   15: (b7) r0 = 1
   16: (55) if r1 != 0x0 goto pc+1
   17: (b7) r0 = 0
   18: (67) r0 <=<= 1
; }
   19: (95) exit
```

```
$ sudo bpftool prog
2390: sched cls  name entrypoint  tag 61f771b5f10a00fc
      loaded at 2026-01-28T15:32:14+0100  uid 0
      xlated 160B  jited 97B  memlock 4096B  map_ids 281,279
      btf_id 310
```



# The naive fix? - Dead code maintains refcount

```
$ sudo bpftool prog dump xlated id 2392
int entrypoint(struct sk_buff * ctx):
; int SEC("tc") entrypoint(struct __sk_buff *ctx)
    0: (b7) r0 = 0
; if (enable feature a)
    1: (18) r1 = map[id:286][0]+0
    3: (71) r1 = *(u8 *) (r1 +0)
; }
    4: (95) exit
```

```
$ sudo bpftool prog
2392: sched cls  name entrypoint  tag 61f771b5f10a00fc
      loaded at 2026-01-28T15:37:22+0100  uid 0
      xlated 40B  jited 31B  memlock 4096B  map_ids 286,284
      btf_id 319
```

# The naive fix? - Dead code maintains refcount

- [resolve\\_pseudo\\_ldimm64](#) runs early
  - Does a flat scan, converting FDs to pointers
  - Adds map to ``env→used_maps``
  - Increments map refcount ``bpf_map_inc(map);``
- Majority of verification happens...
- Dead code elimination happens
  - [opt\\_hard\\_wire\\_dead\\_code\\_branches](#)
  - [opt\\_remove\\_dead\\_code](#)
  - [opt\\_remove\\_nops](#)
- But map refcounts are never released after pointers are eliminated 😞

# First workaround - The double load

- Load our program once, maps set to max\_entries 1
- Read back the jitted code, see which maps remain
- Change LDIMM64 of unused maps FDs to load imm 0xDEADC0DE
- Load with modified instructions and only used maps with full max\_entries
- Works (sort of), but...
  - Loading takes long, a lot of tailcall programs, and large programs
  - We still have to create maps, not free, even small ones
  - We risk hitting the 64 map limit on the first run
  - We cannot gate unsupported map types behind features (arena maps)
  - Reading back instructions impossible when ``kernel.kptr_restrict`` + ``net.core.bpf_jit_harden`` are set.

# Second workaround - matching userspace logic

- Write userspace logic that matches the logic in eBPF to disable maps
- We did not go this route
  - If the compiler re-orders code, the compiled program may not match userspace logic, even though sources do
  - Chances for human error are significant
  - Burdon on maintainers undesirable

# Final workaround - reachability analysis

- What if we knew before loading which instructions in our program are reachable under the given load time configuration?
- A reachability analysis if you will
- Applicable to both maps, tailcalls to hardcoded slots and bpf-to-bpf functions
- Addresses most concerns, but did cost some engineering effort

# Reachability analysis - basic block

```
0: r0 = 0
1: r1 = map[id:281][0]+0
3: r1 = *(u8 *) (r1 +0)
4: if r1 == 0x0 goto pc+14
5: r1 = 0
6: *(u32 *) (r10 -8) = r1
7: r2 = r10
8: r2 += -8
9: r1 = map[id:279]
11: call    htab map lookup_elem#294448
12: if r0 == 0x0 goto pc+1
13: r0 += 56
14: r1 = r0
15: r0 = 1
16: if r1 != 0x0 goto pc+1
17: r0 = 0
18: r0 <<= 1
19: exit
```



# Reachability analysis - basic block

```
0: r0 = 0
1: r1 = map[id:281][0]+0
3: r1 = *(u8 *) (r1 +0)
4: if r1 == 0x0 goto pc+14
---
5: r1 = 0
6: *(u32 *) (r10 -8) = r1
7: r2 = r10
8: r2 += -8
9: r1 = map[id:279]
11: call __htab_map_lookup_elem#294448
12: if r0 == 0x0 goto pc+1
---
13: r0 += 56
14: r1 = r0
15: r0 = 1
16: if r1 != 0x0 goto pc+1
---
17: r0 = 0
18: r0 <=<= 1
19: exit
```

# Reachability analysis - basic block

```
0: r0 = 0
1: r1 = map[id:281][0]+0
3: r1 = *(u8 *) (r1 +0)
4: if r1 == 0x0 goto pc+14
---
5: r1 = 0
6: *(u32 *) (r10 -8) = r1
7: r2 = r10
8: r2 += -8
9: r1 = map[id:279]
11: call __htab_map_lookup_elem#294448
12: if r0 == 0x0 goto pc+1
---
13: r0 += 56
14: r1 = r0
15: r0 = 1
16: if r1 != 0x0 goto pc+1
---
17: r0 = 0
18: r0 <<= 1
19: exit
```

```
graph TD
    L4[4: if r1 == 0x0 goto pc+14] -- red arrow --> L13[13: r0 += 56]
    L13 -- red arrow --> L14[14: r1 = r0]
    L14 -- red arrow --> L15[15: r0 = 1]
    L15 -- red arrow --> L16[16: if r1 != 0x0 goto pc+1]
    L16 -- red arrow --> L18[18: r0 <<= 1]
    L18 -- red arrow --> L19[19: exit]
    L12[12: if r0 == 0x0 goto pc+1] -- red arrow --> L13
```

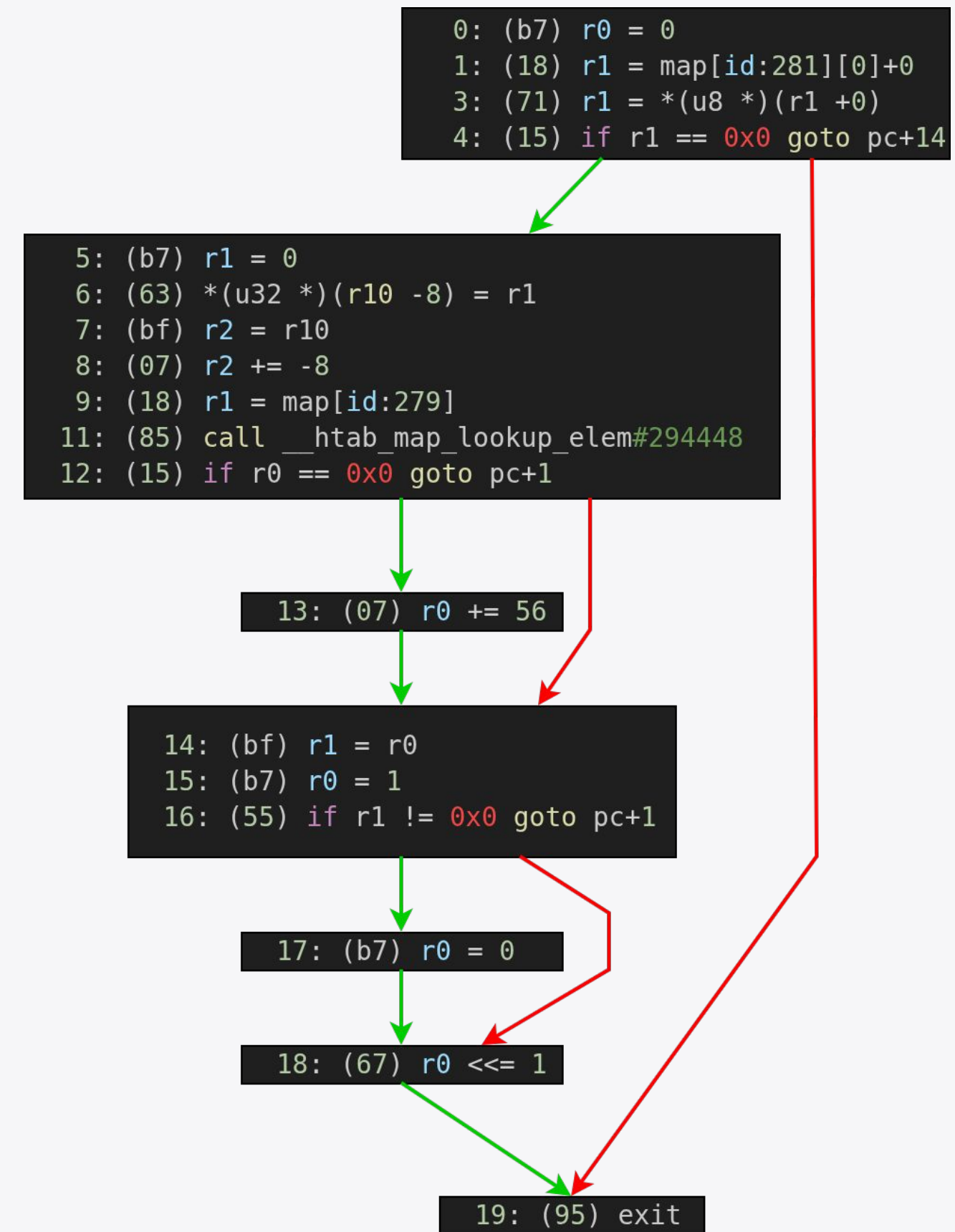
# Reachability analysis - basic block

```
0: r0 = 0
1: r1 = map[id:281][0]+0
3: r1 = *(u8 *) (r1 +0)
4: if r1 == 0x0 goto pc+14
---
5: r1 = 0
6: *(u32 *) (r10 -8) = r1
7: r2 = r10
8: r2 += -8
9: r1 = map[id:279]
11: call __htab_map_lookup_elem#294448
12: if r0 == 0x0 goto pc+1
---
13: r0 += 56
---
14: r1 = r0
---
15: r0 = 1
16: if r1 != 0x0 goto pc+1
---
17: r0 = 0
---
18: r0 <<= 1
---
19: exit
```

The diagram illustrates a control flow graph for a basic block. It starts at line 4 with a conditional jump: `if r1 == 0x0 goto pc+14`. A red line branches from this instruction to line 14 (`r1 = r0`) and continues down to line 19 (`exit`). Another red line branches from line 12 (`if r0 == 0x0 goto pc+1`) to line 14. Line 14 is the entry point for the loop body, which includes lines 15, 16, 17, and 18. A red line branches from line 16 (`if r1 != 0x0 goto pc+1`) back to line 14, forming a loop. The block concludes at line 19 (`exit`).

# Reachability analysis - basic block

```
--- #0 p:[], b:#6, f:#1
0: r0 = 0
1: r1 = map[id:281][0]+0
3: r1 = *(u8 *) (r1 +0)
4: if r1 == 0x0 goto pc+14
--- #1 p:[#0], b:#3, f:#2
5: r1 = 0
6: *(u32 *) (r10 -8) = r1
7: r2 = r10
8: r2 += -8
9: r1 = map[id:279]
11: call htab map lookup_elem#294448
12: if r0 == 0x0 goto pc+1
--- #2 p:[#1], f:#3
13: r0 += 56
--- #3 p:[#1,#2], b:#5 f:#4
14: r1 = r0
15: r0 = 1
16: if r1 != 0x0 goto pc+1
--- #4 p:[#3], f:#5
17: r0 = 0
--- #5 p:[#3,#4], f:#6
18: r0 <=<= 1
--- #6 p:[#0,#5]
19: exit
```



# Reachability analysis - load time config

```
0: r0 = 0
; Get pointer to .rodata map
1: r1 = map[id:281][0]+0
; Deref variable at offset
3: r1 = *(u8 *) (r1 +0)
; Compare to some constant
4: if r1 == 0x0 goto pc+14
```

# Reachability analysis - backtracking

```
--- #0 p:[], b:#55, f:#1
  0: r0 = 0
  1: r1 = map[id:281][0]+0
  3: r1 = *(u8 *) (r1 +0)
  4: r3 = 0
  5: r2 = *(u64 *) (r10 -16)
  6: if r2 == 0x123 goto pc+200
--- #1 p:#0, b:#3, f:#2
  7: r3 = 1
  8: if r1 == 0x0 goto pc+14
```



# Reachability analysis - sign extension

```
0: r0 = 0
; Get pointer to .rodata map
1: r1 = map[id:281][0]+0
; Deref variable at offset
3: r1 = *(u16 *) (r1 +0)
; Cast s16 to 64-bit
4: r1 <=<= 48
5: r1 s>>= 48
; Compare to some constant
6: if r1 s> 0x0A goto pc+14
```

```
0: r0 = 0
; Get pointer to .rodata map
1: r1 = map[id:281][0]+0
; Deref variable at offset
3: r1 = *(u16 *) (r1 +0)
; Cast s16 to 32-bit
4: w1 <=<= 24 ; new in ISAv3
5: w1 s>>= 24
; Compare to some constant
6: if w1 s> 0x0A goto pc+14
```

# Reachability analysis - masks

```
0: r0 = 0
; Get pointer to .rodata map
1: r1 = map[id:281][0]+0
; Deref variable at offset
3: r1 = *(u16 *) (r1 +0)
4: r1 &= 0x01
; Compare to some constant
5: if r1 == 0x00 goto pc+14
```

```
0: r0 = 0
; Get pointer to .rodata map
1: r1 = map[id:281][0]+0
; Deref variable at offset
3: r1 = *(u16 *) (r1 +0)
; Mask and shift bitfield
4: r1 &= 0x04
5: r1 >>= 2
; Compare to some constant
6: if r1 == 0x00 goto pc+14
```

# Reachability analysis - 64 bit constants

```
0: r0 = 0
; Get pointer to .rodata map
1: r1 = map[id:281][0]+0
; Deref variable at offset
3: r1 = *(u16 *) (r1 +0)
; LD64IMM, branching instructions have a 32-bit imm
4: r2 = 0xFFFFFFFFFFFFFFFF
; Compare register to register
6: if r1 == r2 goto pc+14
```

# Reachability analysis - edge cases



```
0: r0 = 0
1: r1 = map[id:281][0]+0
3: r1 = *(u16 *) (r1 +0)
4: *(u16 *) (r10 -8) = r1
...
80: r8 = *(u16 *) (r10 -8)
81: if r8 == 0x00 goto pc+14
```

```
#define CONFIG(name) \
(*({ \
    void *out; \
    asm volatile("%0 = " stringify(name) " 11" \
                : "=r"(out)); \
    (typeof(name) *)out; \
}))

if (CONFIG(enable_feature_a)) {
    // ...
}
```

# Conclusions / final notes

- Reachability analysis seems like a good tool for optimization
  - Reducing map creation and un-releasable maps
  - Reducing load time by pruning unused tail calls and global functions
- The verifier could be improved with regards to map refcounting
  - But even then this likely has a place
- This system has false negatives but no false positives
- Is this a Cilium specific use case? Or might this be useful elsewhere?
- Can we make signing work with this? (at some point in the future)

ISOVALENT  
now part of cisco

# Questions?



**ISOVALENT**  
now part of **cisco**

**Thank  
you!**